Winning Ways for Your Mathematical Plays, Volume 4

Winning Ways for Your Mathematical Plays



Elwyn R. Berlekamp, John H. Conway, Richard K. Guy



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To Martin Gardner

who has brought more mathematics to more millions than anyone else



Elwyn Berlekamp was born in Dover, Ohio, on September 6, 1940. He has been Professor of Mathematics and of Electrical Engineering/ Computer Science at UC Berkeley since 1971. He has also been active in several technology business ventures. In addition to writing many journal articles and several books, Berlekamp also has 12 patented inventions, mostly dealing with algorithms for synchronization and error correction.

He is a member of the National Academy of Sciences, the National Academy of Engineering, and the American Academy of Arts and Sciences. From 1994 to 1998, he was chairman of the board of trustees of the Mathematical Sciences Research Institute (MSRI).



John H. Conway was born in Liverpool, England, on December 26, 1937. He is one of the preeminent theorists in the study of finite groups and the mathematical study of knots, and has written over 10 books and more than 140 journal articles.

Before joining Princeton University in 1986 as the John von Neumann Distinguished Professor of Mathematics, Conway served as professor of mathematics at Cambridge University, and remains an honorary fellow of Caius College. The recipient of many prizes in research and exposition, Conway is also widely known as the inventor of the Game of Life, a computer simulation of simple cellular "life," governed by remarkably simple rules.



Richard Guy was born in Nuneaton, England, on September 30, 1916. He has taught mathematics at many levels and in many places— England, Singapore, India, and Canada. Since 1965 he has been Professor of Mathematics at the University of Calgary, and is now Faculty Professor and Emeritus Professor. The university awarded him an Honorary Degree in 1991. He was Noyce Professor at Grinnell College in 2000.

He continues to climb mountains with his wife, Louise, and they have been patrons of the Association of Canadian Mountain Guides' Ball and recipients of the A. O. Wheeler award for Service to the Alpine Club of Canada.

	Preface to the Second Edition	xiii
	Preface to the Original Edition	xiv
	Solitaire Diamonds!	803
23	Purging Pegs Properly	803
	Central Solitaire	804
	Dudeney, Bergholt and Beasley	805
	Packages and Purges	
	Packages Provide Perfect Panacea	
	The Rule of Two and the Rule of Three	
	Some Pegs Are More Equal Than Others	
	Reiss's 16 Solitaire Position Classes	
	The Continental Board	
	Playing Backwards and Forwards	
	Pagoda Functons	
	The Solitaire Army	
	Managing Your Resources	
	Unproductivity	
	The Prodigal Son's Opening	
	Deficit Accounting and the G.N.P.	
	Accounting for Two-Peg Reversal Problems	826
	Forgetting the Order Can Be Useful	827
	Beasley's Exit Theorems	829
	A Stolid Survivor Problem	829
	Another Hard Problem	831
	The Spinner	833

 \diamond

	Our Fine Finalist834Doing the Splits834All Soluble One-Peg Problems on the Continental Board833The Last Two Moves833A 20-Man Solitaire Army833Fool's Solitaire, Etc.833Beasley Proves Bergholt Is Best833The Classical Problems833References and Further Reading844	455579
24	Pursuing Puzzles Purposefully 843	3
	Soma	3
	Blocks-in-a-Box	4
	Hidden Secrets	4
	The Hidden Secrets of Soma	5
	Hoffman's Arithmetico-Geometric Puzzle	7
	Coloring Three-by-Threee-by-Three by Three, Bar Three	8
	Wire and String Puzzles	9
	The Magic Mirror Method	9
	The Barmy Braid	3
	The Artful Arrow	4
	The Magic Movie Method	4
	Party Tricks and Chinese Rings	6
	Chinese Rings and the Gray Code	
	The Tower of Hanoï	1
	A Solitaire-Like Puzzle and Some Coin-Sliding Problems	
	The Fifteen Puzzle and the Lucky Seven Puzzle	
	All Other Courses for Point-to-Point	
	Rubik's Hungarian Cube — Bűvös Kocka	
	Just How Chaotic Can the Cube Get?	
	Chief Colors and Chief Faces	
	Curing the Cube	
	A: Aloft, Around (Adjust) and About	
	B: Bottom Layer Corner Cubelets	
	C: Central Layer Edge Cubelets	
	D: Domiciling the Top Edge Cubelets	
	E: Exchanging Pairs of Top Corners	
	F: Finishing Flips and Fiddles	
	Explanations	_
	Improvements	
	Elena's Elements	
	Are You Partial to Partial Puzzles?	
	Other "Hungarian" Objects	
	A Trio of Sliding Block Puzzles	7

xi

	Tactics for Solving Such Puzzles	878
	Counting Your Moves	885
	Paradoxical Pennies	
	Paradoxical Dice	886
	More on Magic Squares	886
	The Magic Tesseract	891
	Adams's Amazing Magic Hexagon	
	Strip-Jack-Naked, or Beggar-My-Neighbour **1	
	The Great Tantalizer	
	Polyominoes, Polyiamonds and Searching Policy	
	Alan Schoen's Cyclotome	
	MacMahon's Superdominoes	
	Quintominal Dodecahedra	
	The Doomsday Rule	
	and Easter Easily	
	How Old is the Moon?	
	Jewish New Year (Rosh Hashana)	
	Blocks-in-a-Box	
	The Somap	
	Solutions to the Arithmetico-Geometric Puzzle	
	and One for "Three" Too!	
	Hares and Tortoises	
	The Lucky Seven Puzzle	
	Top Face Alterations for the Hungarian Cube	
	The Century Puzzle	
	Adams's Amazing Magic Hexagon	
	Flags of the Allies Solution	
	All Hexiamond Solutions Found	
	The Three Quintominal Dodecahedra	
	Answer to Exercise for Experts	
	Where Do the Black Edges of MacMahon Squares Go?	
	A Doomsday Answer	
	References and Further Reading	923
\λ/h;	at is Life?	927
***	Still Life	
	Life Cycles	
	The Glider and Other Space Ships	
	The Unpredictability of Life	
	Gardens of Eden	
	Life's Problems are Hard!	
	Making a Life Computer	
	When Glider Meets Glider	
	How to Make a NOT Gate	
	now to make a not Gate	942

25

	The Eater							 				943
	Gliders Can Build Their Own Guns! .							 				947
	The Kickback Reaction							 				947
	Thinning a Glider Stream							 				947
	Building Blocks for Our Computer							 				948
	Auxiliary Storage							 				951
	How We Move Blocks							 				952
	A Little Difficulty							 				954
	Mission Completed—Will Self-Destruct							 				955
	Life is Still Being Lived!							 				958
	Life Computers Can Reproduce!							 				958
	Genetic Engineering							 				959
	Whither Life?							 				959
	References and Further Reading							 				960
Glos	sary											963

Index to Volumes 1–4

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Preface to the Second Edition

This is the volume that has seen the fewest changes in the new edition, if we ignore typographical corrections and minor improvements of wording. Most of them take the form of additional references to the vast amount of work that has been done in our subject since the previous edition appeared.

Some more substantial corrections and additions have been made to the discussion of Rubik's Cube in Chapter 24, based on the ideas of several authors. We have also added a small section to that chapter to report on Marc Paulhus's solution to the question raised by one of us many years ago: "Can a game of Strip-Jack-Naked continue for ever?

The Game of Life still has many adherents more than 30 years after its invention, and some of their discoveries and comments have been incorporated in Chapter 25.

We thank everyone who has written to us with suggestions for improving the book, even when we have not taken their advice. We also thank everyone at A K Peters who has worked on the book, in particular Jon Peters, who has fixed some of the illustrations. Additional thanks go to A and K themselves for having undertaken its republication without knowing what they were in for!

Elwyn Berlekamp, University of California, Berkeley John Conway, Princeton University Richard Guy, The University of Calgary, Canada

February 7, 2004

Preface to the Original Edition

Does a book need a Preface? What more, after fifteen years of toil, do three talented authors have to add.

We can reassure the bookstore browser, "Yes, this is just the book you want!"

We can direct you, if you want to know quickly what's in the book, to page xvi. This in turn directs you to volumes 1,2,3 and 4.

We can supply the reviewer, faced with the task of ploughing through nearly a thousand information-packed pages, with some pithy criticisms by indicating the horns of the polylemma the book finds itself on. It is not an encyclopedia. It is encyclopedic, but there are still too many games missing for it to claim to be complete. It is not a book on recreational mathematics because there's too much serious mathematics in it. On the other hand, for us, as for our predecessors Rouse Ball, Dudeney, Martin Gardner, Kraitchik, Sam Loyd, Lucas, Tom O'Beirne and Fred. Schuh, mathematics itself is a recreation. It is not an undergraduate text, since the exercises are not set out in an orderly fashion, with the easy ones at the beginning. They are there though, and with the hundred and sixty-three mistakes we've left in, provide plenty of opportunity for reader participation. So don't just stand back and admire it, work of art though it is. It is not a graduate text, since it's too expensive and contains far more than any graduate student can be expected to learn. But it does carry you to the frontiers of research in combinatorial game theory and the many unsolved problems will stimulate further discoveries.

We thank Patrick Browne for our title. This exercised us for quite a time. One morning, while walking to the university, John and Richard came up with "Whose game?" but realized they couldn't spell it (there are three tooze in English) so it became a one-line joke on line one of the text. There isn't room to explain all the jokes, not even the fifty-nine private ones (each of our birthdays appears more than once in the book).

Omar started as a joke, but soon materialized as Kimberly King. Louise Guy also helped with proof-reading, but her greater contribution was the hospitality which enabled the three of us to work together on several occasions. Louise also did technical typing after many drafts had been made by Karen McDermid and Betty Teare.

Our thanks for many contributions to content may be measured by the number of names in the index. To do real justice would take too much space. Here's an abridged list of helpers: Richard Austin, Clive Bach, John Beasley, Aviezri Fraenkel, David Fremlin, Solomon Golomb, Steve Grantham, Mike Guy, Dean Hickerson, Hendrik Lenstra, Richard Nowakowski, Anne Scott, David Seal, John Selfridge, Cedric Smith and Steve Tschantz.

Preface

No small part of the reason for the assured success of the book is owed to the well-informed and sympathetic guidance of Len Cegielka and the willingness of the staff of Academic Press and of Page Bros. to adapt to the idiosyncrasies of the authors, who grasped every opportunity to modify grammar, strain semantics, pervert punctuation, alter orthography, tamper with traditional typography and commit outrageous puns and inside jokes.

Thanks also the Isaak Walton Killam Foundation for Richard's Resident Fellowship at The University of Calgary during the compilation of a critical draft, and to the National (Science & Engineering) Research Council of Canada for a grant which enabled Elwyn and John to visit him more frequently than our widely scattered habitats would normally allow.

And thank you, Simon!

University of California, Berkeley, CA 94720 Elwyn Berlekamp University of Cambridge, England, CB2 1SB John Conway University of Calgary, Canada, T2N 1N4 Richard Guy

PRACTICE There are a number of other connexions between various chapters of the book: page 803 page 277 page 461 turn to the little notes about our four main themes: If you want to know roughly what's elsewhere, page 1 Doing It Yourself Adding Games Bending the Rules Case Studies **ΤΗΕΟRΥ** You are now here



without reference to anything earlier, except perhaps the basic ideas at the start of the book. However, you should be able to pick any chapter and read almost all of it



Solitaire Diamonds!

Twinkle, twinkle little star, How I wonder what you are! Up above the world so high Like a diamond in the sky! Jane Taylor, The Star.

We are all in the dumps, For diamonds are trumps; The kittens are gone to St. Paul's. the babies are bit, the Moon's in a fit, And the houses are built without walls. *Nursery Rhyme*

If you've followed everyting in *Winning Ways* so far, you're probably finding it hard to get people to play with you, so you will need something to do on your own. Here are our favorite solitaire diamonds:

The classical game of Peg Solitaire, treated by old and new methods in Chapter 23. A host of puzzles, pastimes and other party tricks in Chapter 24. And finally, every automaton will enjoy playing the notorious game of Life (Chapter 25).

-23-

Purging Pegs Properly

We can merely mention bean-bags, peg-boards, size and form boards, as some of the apparatus found useful for the purpose of amusing and instructing the weak-minded.

Allbutt's Systematic Medicine, 1899, VIII, 246.



Figure 1. The English Solitaire Board.

Figure 1 shows the English Board on which the game of Peg Solitaire is usually played. It's easier to refill the board if you use marbles, but pegs are steadier when it comes to analysis.

The game is played (by one person of course) as shown in Fig. 2. If in some row or column two adjacent pegs are next to an empty space as in Fig. 2(a), then we may jump the peg p over r into the space s (Fig. 2(b)). The peg r that has been jumped over is then removed (Fig. 2(c)). Jumps are like captures in Draughts or Checkers, but they *never* take place diagonally, but only in the East, South, West or North directions.



Figure 2. Making a Solitaire Jump.

Central Solitaire

The standard problem is to start as in Fig. 1, with a peg in every hole except the centre, and then aim, by making a series of these jumping moves, to reduce the situation to a single peg in the central hole (Fig. 3).



Figure 3. Success!

Like many card solitaire ("Patience") games, Solitaire is probably called a *game* rather than a puzzle because one often feels one is playing against an invisible opponent. Many people not normally interested in puzzles will recall some period of their lives when they have struggled with this opponent for days at a time; yet it seems that most of those who can readily solve simple Solitaire problems have been taught the trick by someone else as a child. It is rare indeed to find someone who has acquired the knack single-handed, and surely Peg Solitaire (nowadays selling in many parts of the world under the trade name of Hi-Q) must be the hardest game of its kind to have gamed substantial popularity. It is an ideal game to while away hours of enforced idleness during illness or long journeys, and perhaps we should believe those old books which tell us that the game was invented by a French nobleman who first played it on the stone tiles of his prison cell. If you haven't played this game before, put down this book, go out right now, buy a board, and try to solve the Central Solitaire Game. Those of you who are left will have plenty of time to read the chapter before the novices come back in a week or so—why not learn a particularly elegant solution to impress them all?



Figure 4. A Move of Five Jumps.

Dudeney, Bergholt and Beasley

Since you must already know how to solve the problem, you'll want to do it quickly, so let's agree to count any number of consecutive jumps with a single peg as just one **move**. Figure 4 shows such a move—the five shaded jumped-over pegs are to be taken off as part of the move.

		a	b	с		
	у	d	e	f	z	
g	h	i	j	k	1	m
n	0	р	x	Р	0	Ν
Μ	L	К	J	Ι	Н	G
· · · ·	Ζ	F	E	D	Y	
		C	В	A		

Figure 5. Labelling the Places.

























Figure 6. Dudeney's 19-move Solution for Central Solitaire.

Packages and Purges

In order to describe a solution concisely, we label the places as in Fig. 5, and write S_t for a jump from S to t and shorten this to S when we don't need to indicate the direction. The 5-jump move of Fig. 4 is L_{JHljh} which we will abbreviate to L_5 when it is unambiguous (we can't do that here since L_5 could also mean L_{hjJHl} and various other things). In this notation, Dudeney's elegant 19-move solution of his Central Solitaire problem is

$$eJO_2 fmh_2ap FMH_2AP c_2g_2C_2G_2 p_6o$$
,

and this is set out in Fig. 6.

Dudency thought that the number 19 could not be improved, but, in *The Queen* four years later, Ernest Bergholt gave an 18-move solution, unfortunately not quite as symmetrical as Dudency's:

 $elcPDGJm_2igL_5CpA_2M_2a_3d_5o.$

Here the notation L_5 is ambiguous, but the intended 5-jump move is the one depicted in Fig. 4. The move d_5 is also ambiguous, but either interpretation leads to the same result.

The whole truth emerged only 52 years later, in 1964, when John Beasley used the methods described in this chapter to prove that a solution in fewer than 18 moves is impossible. With Beasley's kind permission we publish his proof for the first time in the Extras to this chapter. It is very condensed, so the reader who wishes to follow it should first study the chapter diligently!

Packages and Purges

It's nice to be able to know the effect of a whole collection of moves before you make them, so let us sell you some of our instant **packages**. When a package is used to clear all the pegs from a region, we call it a **purge**.



Figure 7. Purging Three Pegs.

Figure 7 shows the handy little **3-purge**, our most popular package. When three pegs—the *tail*, the *body* and the *head*—are adjacent in line, this will remove them all, provided the head has an additional peg on one side of it, and an empty space on the other, as in the figure. Move 1 of the package jumps the additional peg *over* the head; move 2 jumps tail over body *into* head; and move 3 jumps back *over* the head to its original position. Since the peg and the

space on either side of the head are essential to the package, but are restored to their original state, we call them **the catalyst**.

In Figs. 8(a) to 8(h), \bigcirc indicates a peg to be purged, \bigcirc a space to be filled, and $\times\times$ indicate catalyst places of which one must be full and the other empty. In most of the purges there are two catalyst moves in opposite directions over the same position (which may initially be either a peg or an empty space) and the remaining moves form one or two packages which deliver pegs to that place. For the 3-purge (8(a)), one peg was already in place and the second is delivered by a single jump which we might call the "2-package" (8(b)). The 6-purge is usually accomplished (8(c)) using a 2-package to deliver the first peg and a 4-package (8(d)) for the second.



Figure 8. A Parcel of Packages.

The L-purge (8(e)) and L-package (8(f)) are very useful indeed. The first peg for the Lpurge is already in place and an L-package supplies the second. The first two moves of the L-package form a 2-purge (8(g)) which can also be used in other situations. The catalyst for the 2-purge is restored in a rather unorthodox way, as is the alternative catalyst for the 6-purge shown in Fig. 8(h).

Packages Provide Perfect Panacea

Plenty of problems are performed with panache by people who purchase our packages.

In Fig. 9 we can see at a glance a solution for Central Solitaire, consisting of two 3-purges (1 and 2) followed by three 6-purges (3, 4 and 5) and an L-purge, leaving only the final jump to be made. You should check that every purge has the catalyst it needs.



Figure 9. Central Solitaire Painlessly Packaged.

Instead of Central Solitaire we can consider other one-peg reversal problems: start with only one empty space and finish with only one peg in the same place. Figures 9 and 10 show that most such problems can be solved by purely purgatory methods, but in Fig. 10(e) we start with a 4-package indicated by the arrow (1), and the notorious problem (b) needs more complicated methods.

To clarify our notation we explain our solution for (b) in detail. For the first jump we have no choice but to jump from the place marked 1 in the figure. Our second jump, from the place marked 2, clears a space which enables us to make the L-package, indicated by the bent arrow (3). We now have a catalyst for the L-purge (4) which is followed by a single jump from the place marked 5. We are now on the home run with purges 6, 7 and 8 followed by a single jump from place 9. If the reader plays this through she will find that we have set up a spectacular 5-jump move from the place marked 10_5 .

The reader might like to try her hand at some *two*-peg reversal problems—start with just two spaces on the board and end with just two pegs in those places.



Figure 10. The Other Six One-Peg Reversals.

The Rule of Two and the Rule of Three

Here is another type of problem (Fig. 11). We start with just one empty space and declare that some particular peg is to be the **finalist** (last on the board). In the example the initial hole is at position d and we want the finalist to be the peg that starts at b. Where must it end?



Figure 11. Find Where the Finalist Finishes!

There is an obvious **Rule of Two**—the peg can only jump an even number of places in either direction, as indicated by the arrows in Fig. 11(a). But there is a much more interesting **Rule of Three**. One of the consequences of this is that if we start with a single space on the English board and end with a single peg, then we can move in steps of three from the initial space to that of the finalist, as in Fig. 11(b).

The Rule of Two and the Rule of Three, taken together, can lead to surprises. See how they point to the unique finishing place H in Fig. 11(a) and (b). Now that we know that H is the only place permitted by both the Rule of Two and the Rule of Three, the problem is a lot easier than it might have been. Figure 12 shows a neatly packaged solution; how did we find it?



Figure 12. (a) The Position after the First Two Moves. (b) The Position before the Last Two Moves.

What we did was plan the 3-jump move 9_3 which puts the finalist in his place, and our second jump was to clear a space for this. But after we made this second jump most of the pegs parcelled themselves up naturally. The one apparent exception was the peg starting just right of the finalist, and the best way of clearing this seemed to be to use it as in move 8_2 to provide the final jump.

For other problems, gentle reader, we recommend a similar procedure. Plan the last few moves of your solution and let the first few be used to smooth the way for these and leave the remaining pegs in tidy packages. Remember that the catalyst for the very last purge must be among the pegs in your planned finale.

Here's a nice finalist problem for you. Let the initial hole be in position B and the finalist be the peg which starts at J. Can you end with only this peg?

Some Pegs Are More Equal Than Others

How do we explain the Rule of Three? The best way is to introduce "multiplication" for Solitaire positions. In Fig. 13(a) the two adjacent pegs s and t can obviously be replaced by a single peg at r, so we write

s t x (a) t х s u (b)6 u t t x u = s (c) 0 s t u s

Figure 13. Multiplying Pegs.

$$st = r$$
,

but Fig. 13(b) shows that we can also write

st = u.

Now Euclid tells us that things that are equal to the same thing are equal, so we must agree that r = u.

> Places three apart in any line are considered equal.

Let's see what other rules of algebra tell us. Combining Figs. 13(b) and 13(c), we have

$$\begin{array}{l} st=u, \ tu=s,\\ st^2u=us, \end{array}$$

or, cancelling,

which seems to tell us that

Remember how catalysts do precisely this—they remove two pegs which are delivered to the same place by the other moves of a purge. In fact it follows from our algebra that

For example, in Fig. 13(c), tu = s, so

$$stu = ss = 1.$$

Three adjacent pegs in line cancel,

(3-purge)

 $t^2 = 1,$

two pegs in the same place cancel.

> any set of pegs that can be purged cancel.

and since r = u,



But in the algebra there are less obvious equalities: for since $s^2 = 1 = rst$, we find

s = rt. **r t r t s Figure 14.** s = rt.

Reiss's 16 Solitaire Position Classes

We've now said enough to see how our algebra cuts the Solitaire board down to size, for since places three apart are algebraically equal, every place is equal to one of the nine in the middle of the board (Fig. 15); for example a = p. Now we can use our most recent rule to express each of these nine in terms of the four corner ones, i, k, I, K:

$$\begin{array}{ll} j = ik & P = Ik \\ p = iK & J = IK \\ x = jJ = ikIK. \end{array}$$



Figure 15. Stripping Down to Essentials.

Since equal pegs cancel,

every Solitaire position is algebraically equal to one of the 16 combinations of the places i, k, I, K.

THE SIXTEEN REISS CLASSES





Figure 16 (b). Reduced to Size.

The position of Fig. 16(a) was found unattended in a Jugoslav railway train. Those filmy packages and letters weren't there—but just came into our mind's eye when we pondered the possibility of reducing the position to a single peg. Where must this single peg be?

Our rules allow us to cancel those four packages of three and then the four pairs RR, SS, TT, UU, so that the position is algebraically equal to the four shaded pegs. We then can move two of these three spaces and cancel another 3-package as in Fig. 16(b) to see that the position equals a single peg at I. So the Rule of Three says that the finalist must be at I, L or f. For which of these places can you find solutions?

How do we know that Reiss's sixteen classes are really different? Might not our algebraic rules imply perhaps that i = kK? No! For consider the numbers ± 1 shown in the places of Fig. 17(i). Whenever three of these numbers

r, s, t

are adjacent in line, we really do have

$$rs = t$$
,

and from this we can see that all our algebraic rules hold for these numbers. But in this system we have

$$i = -1, \ k = K = +1,$$

so we can't prove i = kK! In fact Figs. 17(i, k, I, K) show that all 16 combinations of the pegs i, k, I, K are algebraically distinct: for example the value on Fig. 17(i) is -1 just if i is involved in the combination. Making a Solitaire move or applying any of our algebraic rules will never change the value in any of the four Figures.



Figure 17. "Answers" to the Algebra.

In algebraic language, the first thing we told you about the Rule of Three may be restated a position with just one empty space is algebraically equal to the **complementary** position in which *only* that place is full. More generally,

> any position on the English board is algebraically equal to the complementary position which has empty spaces replacing pegs and pegs replacing empty spaces.

For our rules allow us to complement any line of three adjacent places, and the whole board can be parcelled into such threes.

This property fails for the Continental Board.

The Continental Board

The Continental Board which has the four extra holes at y, z, Y, Z in Fig. 5. So no reversal problems are possible on this board. Which of the problems which start with a single hole and end with a single peg are solvable on this board? See the Extras.

Playing Backwards and Forwards

"The game called Solitaire pleases me much. I take it in reverse order. That is to say that instead of making a configuration according to the rules of the game, which is to jump to an empty place and remove the piece over which one has jumped, I thought it was better to reconstruct what had been demolished, by filling an empty hole over which one has leaped."

The famous philosopher plainly thought that playing Solitaire backwards was different from playing it forwards, but really it's exactly the same game! For let's see what happens when he makes one of his backward moves from Figs. 18(a) to 18(c). Leibniz regards this as jumping *piece t* into *hole r* and *filling* the empty *hole s* over which he has leaped, but Fig. 18(b) shows that we can regard him as jumping the *hole* at r over the *hole* at s into the *piece* at t and removing the *hole* over which he has jumped. (Of course to *remove* a hole he *inserts* a piece!)



Backwards Solitaire is just forwards Solitaire with the notions "empty" and "full" interchanged.

TIME-REVERSAL = ANTI-MATTER?



Leibniz.

This can be useful as well as interesting. A quite spectacular Solitaire finale happens in Beasley's remarkable 16-move solution of the *i*-reversal problem:

$apc_2F_2gdM_2IAP_{\downarrow}f_{\downarrow}C_3Gm\ldots$

After 14 of the 16 moves the board still seems quite full (Fig. 19(a)) but can be cleared to a single peg in just 2 moves. (Can you find them?)



Figure 19. Two Moves to Go! ... And How to Get Back?

How would you find the moves leading up to this positron? The time-reversal trick should make it easy. Instead of reducing a position with only one space, at i, to Fig. 19(a), try to reduce the complementary position (Fig. 19(b)) to just one peg at i. If you've been doing your homework and practising diligently, you won't find this too hard. You too can astonish your friends with grand finales to other Solitaire problems set up by the time-reversal trick.

Pagoda Functons

Reiss's algebraic theory (known to many!) applies even when we allow you to make moves backward in time (like Leibniz) as well as the ordinary forward ones. Of course this lets you take back any of your bad moves, but you may also "undo" moves you haven't even made! If two positions are in different Reiss classes, then we can never get from one to the other by normal moves, by Leibniz's backward moves nor by any mixture of the two.

Unfortunately this means, of course, that the Reiss theory can never tell you when you've made a bad move, because the Reiss class never changes. You need something like the **Pagoda Functions** (known to few!) we are about to show you, that can *change* when you make a move, albeit in a restricted way. Mike Boardman was one of those who helped us to develop these.

Those friends of yours should now be back from the store with their Solitaire boards, so why not present them with a couple of innocent-looking problems? Since these are *reversal* problems, your friends won't be able to prove them impossible even if they've got as far as the last section.

The two problems are shown in Figs. 21(a) and (b) where circles show the only places which are initially empty and which must also be the only places which are finally full. Figures 21(c)and (d) show two pagoda functions which prove the problems impossible. In general, if pag is any such function and X any Solitaire position, we shall write pag X

for the sum of the numbers that pag assigns to the pegs which are present in X. If X is partitioned into smaller positions Y and Z, then, in our algebraic notation we have

$$X = YZ,$$

and

$$pag X = pag Y + pag Z$$

so that pagoda functions behave like logarithms.



Figure 20. The Golden Pagoda.



Figure 21. Two Impossible Reversal Problems.

The essential property which defines pagoda functions is that no move may increase the value. To check this condition you must make sure that

THE PAGODA FUNCTION CONDITION

$$pag r + pag s \ge pag t$$

holds for every conceivable Solitaire jump r over s into t.

CHECK THIS CONDITION NOW IN FIGURES 21(c) AND (d)!

When you've done that you will see the impossibility of our two problems, since the pag (Fig. 21(c)) of the initial position in 21(a), namely 4, can't be increased to 6, the pag of the final position; nor can 8 be increased to 10 (Figs. 21(b) and (c)).

In Fig. 22 we show the pagoda functions you're most likely to find useful; so you'd better check the Pagoda Function Condition for each of them! The values in the blank spaces are zero and you can make any of the indicated swaps. Figures 22(c, d, h, and v) are obvious pagoda functions since they just indicate all the places that a given peg can go to. The 12 places in 22(c) are called **corners** and the 5 in 22(d) are the **dodos**, because one of the easiest mistakes you can make is to let your dodos become extinct when you need one in your final position. Those extra minus ones often make 22(a) and (b) more useful than (h) and (v).


Figure 22. Some Useful Pagoda Functions.

The Solitaire Army

A number of Solitaire men stand initially on one side of a straight line beyond which is an infinite empty desert (Fig. 23). How many men do we need to send a scout just 0, 1, 2, 3, 4, or 5 paces out into the desert?



Figure 23. How About Sending a Scout Out?

It's not hard to see that the answers for 0, 1, 2 and 3 paces are 1, 2, 4 and 8 men, so you might guess that the next two answers are 16 and 32. But in fact no less than 20 men are needed to get 4 paces out. Can you find the two possible configurations of 20 men? (See the Extras.)

For 5 paces the answer is even more surprising—it is *impossible* to send a scout five paces into the desert , no matter how large an army we hire! The pagoda function which proves this is shown in Fig. 24. It was the shape of the graph of this function (Fig. 20) which first suggested the name "pagoda". The number σ is determined by the golden ratio:

$\sigma = \frac{1}{2}(\sqrt{5}-1) = 0 \cdot 618 \dots,$ $\sigma^2 + \sigma = 1.$												
					1							
					σ							
					σ^2							
				σ^4	σ^3	σ^4						
			σ^6	σ^5	σ^4	σ^5	σ^6					
σ^{10}	σ9	σ^8	σ^7	σ^6	σ^5	σ^{6}	σ^7	σ^8	σ^9	σ^{10}		
σ^{11}	σ^{10}	σ^9	σ^8	σ^7	σ^6	σ^7	σ^8	σ^9	σ^{10}	σ^{11}		
	σ^{11}	σ^{10}	σ^9	σ^8	σ^7	σ^8	σ9	σ^{10}	σ^{11}			
		σ^{11}	σ^{10}	σ^9	σ^8	σ^9	σ^{10}	σ^{11}				
		•••	σ^{11}	σ^{10}	σ^9	σ^{10}	σ^{11}					
		•										
σ^{n+5}	σ^{n+4}	σ^{n+3}	σ^{n+2}	σ^{n+1}	σ"	σ^{n+1}	σ^{n+2}	σ^{n+3}	σ^{n+4}	σ^{n+5}		

Figure 24. Pagoda Function for the Solitaire Army.

By some easy mathematics we have

$$\sigma^n + \sigma^{n+1} + \sigma^{n+2} + \ldots = \frac{\sigma^n}{1 - \sigma} = \sigma^{n-2},$$

so that the total score of the line whose middle element is σ^n is

$$\sigma^{n-2} + \sigma^{(n+1)-2} = \sigma^{n-3},$$

and the total score of this line and all lower lines is

$$\frac{\sigma^{n-3}}{1-\sigma} = \sigma^{n-5}.$$

In particular, the sum of *all* the men on or below the σ^5 line is *exactly* 1, so no finite number of these men will suffice to send a scout to the place whose score is 1. But infinitely many men are *almost* enough, because we once showed that if any man of our army is allowed to carry a comrade on his shoulders at the start, then no matter how far away the extra man is, the problem can now be solved.

Managing Your Resources

Your score on a pagoda function is in some sense a measure of your resources, which you should not consume too rapidly. But mere worldly goods are not enough: they must be capably managed to preserve a balance between your commitments in various directions.

The **Balance Sheet** of Fig. 25 has been cunningly devised to do just this. The subtlety of the English board is that you are often forced to consume assets in order to maintain the

		b-1	β	$b^{-1}\beta$	β ² =	
		b	аβ	bβ	ρ- =	= 1
a 1	а	1	а	1	а	<i>a</i> ⁻¹
α	bα	b	<i>c</i> αβ	Ьβ	<i>bαβ</i>	x
$a^{-1}\alpha$	aα	1	аα	1	αα	$a^{-1}\alpha$
α ² =	_ 1	b	ахβ	bβ		
ο.	- 1	<i>b</i> ⁻¹	β	$b^{-1}\beta$		

Figure 25. The Balance Sheet.

balance, as measured by the greek letters α and β , of your position in the North-South and East-West directions. The latin letters a, b and c measure the **assets** on a number of pagoda functions simultaneously (a and b for Figs. 22(a) and 22(b) and abc^2 for Fig. 21(c)).

To estimate the overall capacity of a position, find the product of the resources of all its pegs in Fig. 25, using the relations

$$\alpha^2 = \beta^2 = 1.$$

A problem has two such products, the **raw product** (for its initial position), which must be taken to the **finished product** (for the final position) while consuming the **available resources**:

$$\frac{\text{raw product}}{\text{finished product}} = \text{available resources.}$$

In Fig. 26, all the jumps that change the product are shown to do so in units of sizes

so your available resources will only be **productive** if they can be made up of such units.

Central Solitaire, for example, has raw product a^4b^4 and finished product $c\alpha\beta$ so that its available resources are

$$\frac{a^4b^4}{c\alpha\beta} = a^4b^4c^{-1}\alpha\beta.$$



Figure 26. Using Resources in Various Units.

Unproductivity

In Dudeney's solution only the opening and closing few moves actually use any of these:

Unproductivity

Many problems are impossible for the simple reason that

$$b^2 \alpha$$
 and $a^2 \beta$ are
unproductive!

Why is this? In the case $b^2 \alpha$, for example, we are hamstrung for lack of *a*'s, so the α forces us to make a jump $c\alpha$, leaving only $b^2 c^{-1}$ for the remaining moves, in which c^{-1} demands a move $b^2 c^{-1} \beta = B$, and we then have no assets with which to adjust the remaining β .

The Prodigal Son's Opening

Jump into centre; jump over centre; jump into centre; jump back over centre;

is the only way Central Solitaire can go wrong in as few as four moves. What's so bad about these moves? The prodigality lies in the second and fourth moves which both use $c\alpha$ or both use $c\beta$ and therefore leave only

$$a^4b^4c^{-!}\alpha\beta/c^2 = a^4b^4c^{-3}\alpha\beta$$

for the remaining moves. But

$$a^4b^4c^{-3}lphaeta$$

is unproductive!

For the only way to cope with c^{-3} without overspending either a or b is to use the units

A, A, B or A, B, B

which leave only the unproductive products

$$b^2 \alpha$$
 or $a^2 \beta$.

Of course the same argument shows that *no* two moves in *any* solution of Central Solitaire can have product c^2 .

Can you find the only way (**Fool's Solitaire**) of getting absolutely pegbound (unable to move) in six jumps? And can you **Succour the Sucker** by solving the position reached after five of these moves? And can you flag yourself down to *another* pegbound position in as few as ten jumps from the start? See the Extras.

Deficit Accounting and the G.N.P.

The **deficit** of a problem is the amount by which its initial position lacks the resources of the entire board, combined with the total resources of the final position and the costs of any moves you intend to make. Since the resources of the entire board are

 $a^4b^4c\alpha\beta$ (the (English) Gross National Product)

we have

remaining resources =
$$\frac{a^4 b^4 c \alpha \beta}{\text{deficit}}$$

The deficit is found very easily by multiplying the initial hole values by the final peg values. For Central Solitaire, the basic deficit is

$$c\alpha\beta.c\alpha\beta = c^2$$

which the Prodigal Son's bad moves extravagantly enlarged to c^4 . He clearly didn't know the **Deficit Rule**:

If deficit/ c^4 IS productive, your remaining resources AREN'T!

This is because (G.N.P.)/ c^4 is our unproductive product $\alpha^4 \beta^4 c^{-3} \alpha \beta$.

Accounting for Two-Peg Reversal Problems

We know that all the one-peg reversal problems are possible, but there are just four different impossible two-peg reversal problems. The first of these is **Hamlet's Memorable Problem** (to be or not to be):

Get to only b, e present	(to be)
from only b, e absent	(not be).

- Deficit account for Hamlet.	s Peoblem
Io: <u>Deficit account for Hamlet</u> . Initial holes @ b & c :	<i>в.ар</i> = 2
Final pegs @ b& e:	$\beta.a\beta = a$
First & last jumps into e:	$\cos \cos = c^2$
Jump into b:	a²
<u>Deficit :</u>	$\underline{a^4c^2}$

Since

$$\frac{a^4c^2}{c^4} = a^4c^{-2} = A^2$$

is productive, Hamlet's Problem succumbs to the Deficit Rule. The other three impossible two-peg reversals are the **Dodo Problems**, for which the two places are two of the five dodo pegs (Fig. 22(d)). Deficit accounts for the typical problems eo, ex and eE are:

Dodo Problem	eo	ex	eE
Initial holes and final pegs	$(\alpha\beta.b\alpha)^2$	$(a\beta.clphaeta)^2$	$(a\beta.alphaeta)^2$
Required moves	$c\alpha.c\beta$	$c\alpha$	$c\alpha.c\alpha$
Deficit	$a^2b^2c^2lphaeta$	$a^2c^3\alpha$	$a^{4}c^{2}$
$\operatorname{Deficit}/c^4$	A.B	A	A.A

The reader who has been paying attention will have no difficulty in finding solutions to any other two-peg reversal problem.

John Conway, Mike Guy and Bob Hutchings have shown that the only impossible *three*-peg reversals are typified by

- 1. The **Bumble-bee Problems** (b, e and any third place other than g, m, M, G),
- 2. The **Deader Dodo Problems** (two dodos and any third place other than an outside corner *acgmMGCA*),
- 3. The **Three B'ars Problems** (any three of the unlucky 13 places in the three rows *def*, *nopxPON*, *FED*).

These can be shown to be impossible by deficit accounting . In fact in any reversal problem, an additional peg other than an outside corner merely aggravates the deficit.

Forgetting the Order Can Be Useful

If you allow yourself to have 2 or more, or -1 or less, pegs in a hole, you can make your moves in any order! It's a good idea to alter a hard problem in this way, and when you've solved the altered problem, go back and find a sensible order for the original one.

We'll do out loud for you the tricky 3-peg reversal:

start with 0 pegs in b, N, n; 1 peg everywhere else; end with 1 peg in b, N, n; 0 pegs anywhere else.

In the altered problem it's easier to

start with -1 peg in b, N, n; 1 peg everywhere else; end with 0 pegs anywhere.



Figure 27. Solving a Tricky 3-peg Reversal Problem.

From the starting position, Fig. 27(a), we'll need, at some time, to make three jumps to fill those -1's, so we make these three jumps *now*, reaching Fig. 27(b). The only plausible way to deal with the six isolated corner pegs in this is to jump them inwards, and after the indicated upwards jump over the centre we reach Fig. 27(c). The remaining pegs in this can be cleared by a 3-purge, a 6-purge and four double jumps over the inner corners.

If you follow Fig. 27(d) in the order A to L you'll find yourself making all the above moves in a legal way. The double jumps have been incorporated into the 3-purges B, C, I and K.

Beasley's Exit Theorems

Sometimes you can work out exactly what moves to make in a problem, but find it hard to get them into the right order. The following remarks can help you get your moves in order, or prove that it can't be done.

> A region of at least three squares that starts full *or* ends empty needs at least one *exit move*. A region of at least three squares that starts full *and* ends empty needs at least two exit moves.

BEASLEY'S FIRST AND SECOND EXIT THEOREMS

An exit move for a region is a jump that empties some square in the region and fills some square outside the region. To justify Beasley's Second Exit Theorem, note that the first and last moves affecting a region must both be exits. We'll illustrate with a stolid survisivor problem.

A Stolid Survivor Problem

Suppose we want to do an *a*-reversal, with the added condition that peg K is the **stolid survivor**, i.e. that the first move of K is also the final move from K to a. Can the grand finale be a 6-chain?

The ideas of the first part of our discussion are often useful in long chain problems. Then we'll try to put the moves we've found into order, using Beasley's Exit Theorems.

How do we use the 16 side pegs h^8v^8 of Figs. 22(h) and (v)? Each of the outer corner pegs must at some time be jumped into the central 3×3 square, and those at C and M must sidestep first to avoid the stolid survivor at K. So the jumps mentioned use up side pegs as follows:

leaving h^3v^4 for the remaining jumps. Since the first move uses one side peg and the final chain six, we have accounted for all the side pegs and no *other* move can destroy one.

This forces us to make the first move c_a , since the alternative, i_a , would move an inside corner peg to the outside and make us use another side peg to bring it back later. Next k_c would use another side peg, so the second move is j_b and this peg must stay at b until the grand finale because a move refilling e would use yet another side peg. We now know that the final 6-chain uses h^2v^4 and involves a horizontal jump over b, so it must be as in Fig. 28(a).

Since K doesn't move till the end, L can't be jumped over and can only be cleared by the upward jump L_h . We need to make two jumps over B, once to get corner peg C out,



Figure 28. Can the Stolid Survivor Make a Grand 6-chain Finale?

and once in the finale, so we must deliver an extra peg there by a downward jump J_B . For similar reasons *two* extra pegs are needed at D, so we must make two downward jumps P_D . For the second of these, and for the finale, we need two more pegs delivered at P; these must come from N and p. We've now found 23 (Fig. 28(b)) of the 31 jumps. If we make these we arrive at Fig. 28(c). The two pegs on each of I and k must be cleared by pairs of vertical or horizontal to and fro jumps, and the four on i by two such pairs.

To find the right order in which to make these moves we use Beasley's Second Exit Theorem. Consider the region of Fig. 29(a). The moves we've copied from Fig. 28(b) incorporate just one exit from the region; the vertical jump across P. To make sure there's another we must remove the two pegs on k by a *vertical* pair of to and fro jumps as in Fig. 29(b). But the region of *that* figure can now have only one exit, the vertical jump across f. So our problem's impossible!



Figure 29. Using Beasley's Second Exit Theorem.

Another Hard Problem



Figure 30. The Reversal Problem *abceg*.

We'll try the 5-peg reversal problem *abceg*, i.e. start with the board full except for spaces at a, b, c, e, g, and finish with pegs in just those places. An equivalent problem is to clear the board of Fig. 30(a), which has negpegs (or **negs**) in each of the places a, b, c, e, g and pegs in the other 28 places.

For the original problem the pagoda function of Fig. 30(b) changes from 20 to 16, or, in the form of Fig. 30(a), from 4 to 0. This pag kills any possible move across b, which would lose 8, so the jumps i_a , k_c shown in Fig. 30(c) are forced, as is the jump j_b to fill b. In order to make this last jump a peg must be delivered to e, and e must also be full by the end, so two jumps x_e are also needed. If we make these five jumps, using negs where needed, we arrive at



Figure 31. We Make Some Progress.

Fig. 31(a), whose resources are $a^2b^4c^{-1}\alpha\beta$, showing a deficit of a^2c^2 . The Deficit Rule tells us we can't make another move of value a^2 , since $a^4c^2/c^4 = A^2$, so the jump M_g (Fig. 31(b)) is forced. Now use the pagoda function of Fig. 31(c). Its value for Fig. 31(a) is 2, so the peg at N can't jump inwards, nor can we jump over it upwards, since these moves lose 4 on this peg. So the jump m_G (Fig. 31(b)) is forced. The two pegs at G must now both jump to I, and a peg must be delivered to H for the second of these. This can't come from l, as this loses 4 on the 1st pag, so the jump J_H is also forced. Moreover, as l can't jump downwards, and can't be jumped over (this would lose a) it must make the jump l_j over k. This needs delivery of a peg at k, which can't come from i (loses 6 on the last pag) so the jump I_k is forced. If we make all these moves, which have been collected in Fig. 31(b), we arrive at Fig. 32(a), for which the resources are now

$$a^{2}b^{3}c^{-1}\alpha = Ab^{3}$$
 or $A.B.b.c\beta$ or $B.a.a\alpha.b\beta$

so there is no jump $c\alpha$.



Figure 32. A Cul-de-Sac!

There are two ways in which we might remove the peg at O: by j_{lHJ} or by O_x . The former (after the necessary delivery I_k) leads to Fig. 32(b) which is impossible to clear, as the pag of Fig. 32(c) shows. So O_x is forced (Fig. 33(a)) and this requires the delivery D_P (horizontal delivery is prohibited by the pag of Fig. 31(c)). These two jumps lead to a position whose resources $a^2 \alpha \beta$ are uniquely productive: $(a^2 c^{-1} \alpha)(c\beta)$ and the jump E_x is forced. The L-package of Fig. 33(a) will deliver a second peg to K and the board is cleared by L_J and h_{IJj} .



Figure 33. The Problem Solved.

The 23 jumps are shown in Fig. 33(b). How do we do them in practice? In what order? The answer isn't unique, but one possibility is given in Fig. 33(c). It involves two L-packages, 10 and 11, and two chain moves, 13_3 and 14_2 .





Figure 34. The Spinner.

If you start with empty spaces at b, B and marked pegs at g, M, G, m, can you finish with just the four marked pegs on the board in the respective positions M, G, m, g?

Extras

Our Fine Finalist

The Rule of Two and the Rule of Three together tells us that if the initial hole is at B, then a finalist that starts at J must end in either B or b. Here's a solution for b:



The letters indicate the order of the successive moves, except for the finalizing flourish. The bent arrow we've used for move I is our notation for an L-package, as distinct from an L-purge.

However, it's impossible for the finalist to finish at B. This is because there are forced moves $J_B \quad x_E \quad x_E \quad J_B$

which consume

 a^2 $c\alpha$ $c\alpha$ a^2

on the Balance Sheet, giving a deficit of a^4c^2 . Since

$$\operatorname{deficit}/c^4 = a^4 c^{-2}$$

is productive, your remaining resources aren't.

Doing the Splits

If you start from Fig. 16(a) and make the moves A to I indicated in Fig. 35, where the pairs of circles C, F, G indicate 2-purges, you'll reach a 5-peg configuration which can easily be reduced to I, L or f. (We found this solution by the ordering process after subtracting this 5-peg configuration from the starting position.)



Figure 35. The Train was going to IvanicGrad, Ljubljana or foća.

All Soluble One-Peg Problems on the Continental Board

All soluble one-peg problems on the Continental board were found by Reiss using his theory. In our language, the Continental board is algebraically equal to its centre, and so for a one-peg problem to be soluble we must have

(initial hole)
$$\times$$
 (final peg) = centre,

in our algebraic sense. You can easily check that the initial hole and final peg are at opposite ends of an arrow in

$$(a \, p \, C \, O) \leftrightarrow (A \, P \, c \, o) \quad (e \, G \, J \, M) \leftrightarrow (E \, g \, j \, m).$$

There is a 41-hole board for which Lucas gives all the soluble problems; but see the appendix to his *second* edition, because he first conjectured that most of the problems were insoluble!

The Last Two Moves

The last two moves in Fig. 19(a) are n_9G_3 ,

A 20-Man Solitaire Army

A 20-man solitaire army can get a scout 4 places out by arranging itself as shown in Fig. 23. The two men with guns can be moved to the shaded places so as to obtain the only other arrangement.

Fool's Solitaire, Etc.

If each of your moves is confined to the middle row or column you'll reach a position like Fig. 36(a) after six jumps. The next pegbound position is the Hammer and Sickle position of Fig. 36(b), reached after ten jumps.



Figure 36. (a) Sickle and Sickle, (b) Hammer and Sickle.



Figure 37. Succouring the Sucker.

(d)

(c)

To succour the sucker who's made five of the six moves leading to Fig. 36(a) it's best to try to clear Fig. 37(a) to zero. If you set this up on the board (use an upside down peg for the -1!) you should see how the moves of Fig. 37(b) suggest themselves in order, leading to the easily cleared position, Fig. 37(c). The L-moves in Fig. 37(b) are L-packages, not L-purges . You then have the tricky little problem of arranging the moves in order, one solution of which is given in Fig. 37(d), in which A and F are L-purges, but C is an L-package.

Beasley Proves Bergholt Is Best

Suppose there were a 17-move solution to Central Solitaire. Then Beasley first uses the scoring function of Fig. 38(a) ("score" refers to this function—which is *not* a pagoda function—throughout the proof) and his First Exit Theorem, to show that no move begins or ends on b, n, B or N. The initial and final scores are 20 and 0. Moves which begin or end on b, n, B or N increase the score by at least 1. Others decrease it by at most 2 (the careful reader will make a table of score changes for each type of move).



Figure 38. Scoring Function and Regions Used in Beasley's Proof.

Any solution to Central Solitaire contains 11 reserved moves: the first, which we'll take to be e_x , the last, a single jump into the centre, the penultimate one, taking a peg to j, p, J or P, and eight moves bringing the outside corner pegs to inside corner squares so they may be captured.

The first and last moves each increase the score by 2, the penultimate one decreases it by at most 1 and each of the other eight decreases it by at most 2. So the other six (loose) moves must decrease the score by at least 7.

The second move is a loose move, either J_j or h_j say. The move J_j doesn't change the score and it leaves the region of Fig. 38(b) full. The first exit from this region is a loose move, either of type b_j or of type (ending with) h_j . The former increases the score by 2 and the other four loose moves would have to decrease it by at least 9, which is impossible. The latter decreases the score by 1, and our four loose moves have to reduce it further by at least 6. If any of these increased the score, the others could not then reduce it to zero, so moves starting or ending on b, n, B or N are again impossible. Such a move *might* occur as the penultimate one, but

 \diamond



Figure 39. What is the First Exit?

the six loose moves would then have to reduce the score by 10, and the same argument shows this to be impossible.

The second move h_j reduces the score by 1 and is a first exit from the region of Fig. 38(b). The other five loose moves must reduce the score by at least 6. What is the first exit from the region of Fig. 38(c)? There are several possibilities, all of them loose moves, which we'll leave the reader to pursue. In some cases he'll want to ask a further question about one of the regions of Fig. 39, whichever is still full. From now on we'll assume that no move begins or ends on b, n, B or N.

How do we clear a, b and c? We've proved that b can't jump out, so there must be a jump over it, say c_a . The two pegs at a now force two jumps a_i and a jump into d, which we shall call a **side delivery**. The four jumps

$$c_a \quad a_i \quad a_i \quad ?_d$$

are parts of at least three moves

 $a_i \dots ? \dots d c_{ai} \dots$ (the normal case), or $a_i \dots ? \dots d ? \dots kcai$ (a U-turn).

However a U-turn demands a previous clearance of c and an extra side delivery to f.

Since the same argument applies at n, B and N, we shall need at least four side deliveries, none of which can be among our 11 reserved moves, and none of which can be the first exit from Fig. 38(b). This accounts for 16 moves; call the other the **spare**. Moreover, if a U-turn is involved we have a further side delivery, and so no spare. Note that after *eha*, p is a side delivery, but j doesn't count as one while g is still occupied, because we'll still need one to clear gnM.

The final stage of Beasley's proof just enumerates all the variations. In the list below the spare move is in **bold**. In the first two variations the first exit from Fig. 38(b) is L, and in all the others it's h. Each variation ends with \ddagger , \S or a colon and a number.

- ‡ means that the next move can't be a corner move or a side delivery, but the spare has already been used,
- § means that there aren't enough moves left to reduce the score to zero, and
- :9 refers to variation number 9, for example.

This list of variations covers the cases where no U-turns are used. If there is a U-turn then is *no* spare move so we have only the variation $ehapc_2$ [‡](cf. 56).

The Classical Problems

1	$eJLCpA_{2}^{\dagger}$	17	$ehKMJg_2$ ‡	33	eha j gpL§	48	ehap F c3 ⁺
2	$D\S$	18	\bar{CD} §	34	c_2L §	49	M_2
		19	$P_{2}g_{2}^{+}$	35	$J_2 M_{2*}^{+}$	50	$c_2 M J g_2$
3	eh x af ‡	20	A_2	36	$l_2 M_{2*}^{*}$	51	$Mc_{2}:50$
4	pc_2	21	$d_2g_2_{\ddagger}_{\ddagger}$	37	L§ – –	52	$\bar{Jc}_{3^{\ddagger}}$
5	Lap§	22	A_2	38	$J_2 M_2 p_d c_2^+$	53	g_{3}^{+}
6	<i>gj</i> §	23	<i>j</i> ₂ §	39	$l_2 M_2 p_d c_2 \ddagger$	54	c2:50
7	$kcPa_2$	24	$CD\S$			55	$g_{2}c_{2}$;
8	mH§	25	P_2A_2	40	ehapc:30	56	$c_2 P_{2*}^{+}$
9	J_2a_2	26	MJ:19	41	k _2§	57	F,MJg,\ddagger
10	$G_{2^{+}}$	27	<i>j</i> ₂ §	42	<i>x</i> :4	58	x:4
11	L_{3} §	28	d_2A_2 ‡	43	L:5	59	L§
12	mH§	29	MJ:21	44	k mH§	60	jgL§
13	$J_{l}G_{2}^{*}$			45	$J_l G_2$	61	$J_{2}M_{2}^{+}$
14	cP:9	30	ehacpa‡	46	L_3 §	62	$l_2 M_{2*}^{\dagger}$
15	L_3G_2ap §	31	x :3	47	Pc_2 ‡	63	P ‡
16	c_2 §	32	L:5			64	FMj_2 §
						65	Jg_2 ‡

The Classical Problems

These are: start with one empty space, finish with a single peg. They include the reversals, for which Bergholt's results were:

bdeijreversal ax - $_{in}$ 1618 16191616 18moves.

We've just seen that his x-reversal is best possible, but Harry 0. Davis has given a 15-move solution of the i-reversal:

 $kmh_2cPKCD_{PF}A_3MG_2H_4a_2d_5g_3.$

And here are his solutions, which equal Bergholt's, for the *b*- and *j*-reversals:

 $jhapc_2xl_hIf_PA_2GJm_2gL_{Hh}M_2CB_5,$ $hKCd_2MJkmH_{Jl}G_3cA_2D_{Fd}g_2ab_7.$

Hermary identified the 21 distinct problems, one place empty to one place full (see Lucas) and Davis has made a table of best known solutions (see Martin Gardner, "The Unexpected Hanging and Other Mathematical Diversions"). The numbers of moves are:

aa	ap	aO	aC	bb	bn	bx	bB	dd	dK	dH	ee	eM	eJ	ii	il	$_{jj}$	jg	jE	xx	xb
16	16	17	16	18	17	18	18	16	15	16	19	17	17	15	16	16	16	17	18	17

For this information we thank Wade E. Philpott, who has copies of the solutions. Omar will want to find better ones, or prove them best possible.

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-24-

Pursuing Puzzles Purposefully

The chapter of accidents is the longest chapter in the book. John Wilkes

I shall proceed to such Recreations as adorn the Mind; of which those of the Mathematicks are inferior to none. William Leybourne; *Pleasure with Profit.*

We know you want to use your winning ways mostly when playing with other people, but there are quite a lot of puzzles that are so interesting that you really feel you're playing a game against some invisible opponent—perhaps the puzzle's designer—maybe a malevolent deity. In this chapter we'll discuss a few cases where some kind of strategic thinking simplifies the problem. But because we don't want to spoil your fun we'll try to arrange not always to give the *whole* game away.

Soma

This elegant little puzzle was devised by Piet Hein. Figure 1 shows the seven non-convex shapes that can be made by sticking 4 or fewer $1 \times 1 \times 1$ cubes together. Piet Hein's puzzle is to assemble these as a $3 \times 3 \times 3$ cube.



Figure 1. The Seven Pieces of Soma.

We advise you to use seven different colors for your pieces as in the figure. Many people solve this puzzle in under ten minutes, so it can't be terribly hard. But we've got a distinct feeling that it's much harder than it ought to be. Is this just because the pieces have such awkwardly wriggly shapes?

Blocks-in-a-Box

Here is another puzzle invented by one of us some years ago, in which all the pieces are rectangular cuboids but it still seems undeservedly hard to fit them together. We are asked to pack one $2 \times 2 \times 2$ cube, one $2 \times 2 \times 1$ square, three $3 \times 1 \times 1$ rods and thirteen $4 \times 2 \times 1$ planks into a $5 \times 5 \times 5$ box (Fig. 2). It's quite easy to get all but one of the blocks into the box, but somehow one piece always seems to stick out somewhere. A friend of ours once spent many evenings without ever finding a solution. Why is it so much harder than it seems to be?



Figure 2. The Eighteen Pieces for Blocks-in-a-Box.

Hidden Secrets

In our view the good puzzles are those with simple pieces but difficult solutions. Anyone can make a hard puzzle with lots of complicated pieces but how can you possibly make a hard puzzle out of a few easy pieces?

When a seemingly simple puzzle is unexpectedly difficult, it's usually because, as well as the obvious problem, there are some hidden ones to be attended to. Both Soma and Blocksin-a-Box have such hidden secrets, but let's look at a much simpler puzzle, to fit six $2 \times 2 \times 1$ squares into a $3 \times 3 \times 3$ box, leaving three of the $1 \times 1 \times 1$ cells empty—the *holes* (Fig. 3). This now seems fairly trivial, but even so there's a hidden secret which sometimes makes people take more than 5 minutes over it. This hidden problem comes from the fact



Figure 3. A Much Simpler Puzzle.

The Hidden Secrets of Soma

that the square pieces can only occupy an even number of the cells in each horizontal layer. So since 9 is odd each horizontal layer must have a hole and there are only just enough holes to go round. Of course these holes must also manage to meet each of the three layers in each of two vertical directions—you can't afford to have two holes in any layer, because some other layers would have to go without.

So the problem wasn't really to fit the *pieces* in but rather the *holes*. Only when you've realized this do you see why the unique solution (Fig. 4) has to be so awkward looking, with the holes strung out in a line between opposite corners rather than neatly arranged at the top of the box.

Perhaps you'd like to try the big Blocks-in-a-Box problem now, before looking at the extra hints in the Extras.



Figure 4. Six Squares in a $3 \times 3 \times 3$ Box.

The Hidden Secrets of Soma

It's because the Soma puzzle pieces have to satisfy some hidden constraints as well as the obvious ones, that it causes most people more trouble than it should. Let's see why.

The $3 \times 3 \times 3$ cube has 8 vertex cells, 12 edge cells, 6 face cells and 1 central cell as in Fig. 5.



Figure 5. The Vertex, Edge, Face and Central (invisible) Cells.

Now the respective pieces can occupy at most

of the vertex cells, so just one piece, the **deficient** one, must occupy just one less vertex-cell than it might. The green piece can't be deficient without being doubly so, and therefore:

the Green piece has its spine along an edge of the cube.

Now let's color the 27 cells of the cube in two alternating colors,

Flame for the 14 FaVored cells, F and V, Emerald for the 13 ExCeeded ones, E and C.

Then in *one* solution that we know, the respective pieces occupy

but the Yellow, Orange, bLue and Red pieces, and we now know also the Green piece, *must* occupy these numbers in *every* solution, and therefore so must the White and the Black, since an interchange of colors in either or both of these would alter the totals.

The White piece occupies 2 FV cells, 1 EC cell.

The Black piece occupies 1 FV cell and 3 EC ones.

For the placing of a single piece within the box, these considerations leave only the positions of Fig. 6 (which all arise). You'll see that up to symmetries of the cube, the placement of any single piece is determined by whether or not it is deficient and whether or not it occupies the central cell.

The hidden secrets of Soma make it quite likely that one of the first few pieces you put in may already be wrong, when of course you'll spend a lot of time assembling more pieces before such a mistake shows its effect. This would happen for instance if you started by putting the

846



Figure 6. All Possible Positions for the Seven Soma Pieces.

corner of the White piece into a corner of the cube. But if you only put the pieces into the allowed positions, you'll find a solution almost as soon as you start. The complete list of 240 Soma solutions was made by hand by J.H. Conway and M.J.T. Guy one particularly rainy afternoon in 1961. The SOMAP in the Extras enables you to get to 239 of them, when you've found one—and located it on the map!

Hoffman's Arithmetico-Geometric Puzzle

A well-known mathematical theorem is the inequality between the arithmetic and geometric means:

$$\sqrt{ab} \le \frac{a+b}{2}.$$

Figure 7 provides a neat proof of this in the form

$$4ab \le (a+b)^2$$

and the three variable version

$$27abc \le (a+b+c)^3$$

has prompted Dean Hoffman to enquire whether 27 $a \times b \times c$ blocks can always be fitted into a cube of side a + b + c. This turns out to be quite a hard puzzle if a, b, c are fairly close



Figure 7. Proof of the Arithmetico-Geometric Inequality.

together but not equal. A good practical problem is to fit

27 $4\times5\times6$ blocks into a $15\times15\times15$ box.

With these choices, as for any others with

$$\frac{1}{4}(a+b+c) < a < b < c,$$

it can be shown that each vertical stack of three blocks must contain just one of each height a, b, c, while there must be just three of each height in each horizontal layer. There must be the same unused area on each face (just 3 square units in the $4 \times 5 \times 6$ case).

It's almost impossible to solve the puzzle if you don't keep these hidden secrets constantly in mind because you'll make irretrievable mistakes like making a stack of three height 5 blocks, or leaving a 2×2 empty hole on some face. When you *do* keep them in mind, the puzzle becomes much easier, being only extremely difficult! You'll find some information about solutions to Hoffman's puzzle in the Extras.

Coloring Three-by-Three-by-Three by Three, Bar Three

In Hoffman's $3 \times 3 \times 3$ puzzle, the three lengths along any line of three had to be different. Can you color the cells of a $3 \times 3 \times 3$ tic-tac-toe board with

three different colors,

using all

three colors the same

number (9) of times, in such a way that *none* of the $\frac{1}{2}(5^3 - 3^3) = 49$ tic-tac-toe lines uses

three different colors,

nor has all its

three colors the same?

Wire and String Puzzles

Figure 9 shows a number of topological puzzles which can be made with wire and string. It's a pity that manufacturers don't seem to know about all of these.

You wouldn't expect to be able to say much about such varied looking objects, but in fact there's a quite general principle which helps you to solve a lot of them.

The Magic Mirror Method

We'll just take the one-knot version of the puzzle shown in Fig. 9(c) which has been commercially sold as The Loony Loop (Trolbourne Ltd., London). You're to take the string off the rigid wire frame in Fig. 8(a).

If only that rigid wire were a bit stretchable, the puzzle would be quite easy. After squashing the string up (Fig. 8(b)) so as not to get in the way, we could stretch the loops over the ends (Fig. 8(c)) and shrink them again (Fig. 8(d)). After this we can take the string right off (Fig. 8(e)) and then put the loops back as they were (Fig. 8(f)) so as not to upset the owner.

Now the change from Fig. 8(b) to Fig. 8(d) could be accomplished by continuously distorting space. Think of embedding the puzzle in a flexible jelly, if Mother has one made up. Now old-fashioned fairgrounds had special mirrors which seemed to distort space in very funny ways. Now let's imagine a magic mirror with the wonderful property that the distortion is just what's required to make Fig. 8(b) look like Fig. 8(d). Hold the wire frame absolutely still before the magic mirror (Fig. 10(a)) and bunch the string up until it's almost a single point on the axis. Because the space distortion was continuous, its image will also be almost a single point on the image axis.



Figure 8. Solving The Loony Loop.



Figure 9. Shifting Rings, Strings



..... and Other Things.



Figure 10. The Magic Mirror.

Now, very carefully, move the string in just such a way that its *image* in the magic mirror moves completely away from the wire and shrinks to a small point at some little distance from it. Once again, because the distortion was continuous, the real string must now be almost a point, some distance from the wire, and you've solved the puzzle. Easy, wasn't it?

In such cases it often helps to imagine an intermediate distortion. In Fig. 11 we show two stages in an intermediately distorted one-knot Loony Loop . Perhaps you're ready for the two knot version (Fig. 9(c))? Or the Double Treble Clef (Fig. 9(b)) (Pentangle Puzzles, Over Wallop, Hants, U.K.)?



Figure 11. A Less Distorting Mirror.

If a puzzle has got just one completely rigid piece and a number of completely flexible pieces then you can often use the magic mirror method to pretend that the rigid piece is also flexible. For instance, although it may seem impossible to make the braided piece of paper in



Figure 12. Can You Braid this Strip of Paper?

Fig. 12 without glue, it can be undone quite easily. This principle is quite familiar to craftsmen in leather. (To make it you should start braiding at one end and undo the tangle which forms at the other.)

The Barmy Braid

The Barmy Braid problem appears for the first time in this book. It's to take the string off the rigid wire frame in Fig. 13(a). You know you can do it, because in a suitable magic mirror it looks like Fig. 13(b).



Figure 13. Barmy Braid Meets Magic Mirror.

Chapter 24. Pursuing Puzzles Purposefully

The Artful Arrow

Figure 8(a) is our version of a puzzle that appears in many different forms. The basic framework is often a bar of wood with a drill hole in place of our hexagon. We have even seen a version in which the ends of our arrow are a giant's arms and the central hole his nostrils, but the solution is always the same! You can solve this puzzle, and some similar ones, by a modification of the Magic Mirror Method which we call

The Magic Movie Method

If the Artful Arrow had a much smaller ring, there'd be no difficulty about solving it; we'd just slide the ring along the string from the tail of the arrow to its head. Let's suppose we have a kinematic friend who takes a movie of this, but that through some accident with his filters, the string doesn't show up too well, so that what the movie shows is the rigid arrow framework and a little ring that wanders about in space. In fact the ring moves downward through the hexagon (1 to 2 to 3 in Fig. 14(a)), sweeps around (3 to 4 to 5) and then comes safely back up again (5 to 6 to 7).



Figure 14(a). The Magic Movie M_0 .

854



Figure 14(b). An Intermediate Half-Magic Movie.

What we want to do is to watch this movie in a sort of hyperspace magic mirror which distorts both space and time. Our friend can arrange this for us by taking the movie M_0 to the animation department where they can change the whole movie bit by bit, first to M_1 , in which the ring goes down through the hexagon and wanders about a bit less before it comes back up again, then to M_2 , in which it wanders hardly at all before coming up, then to M_3 in which it only takes a timid dip through the hexagon, then in M_4 not at all, while in M_5 , M_6 ... the size of the ring gradually increases until it is too big to go through the hexagon.

The trouble with all these movies is that we can't see the string! But since we intend the sequence of movies to realize a continuous distortion of space-time, we can ask the animation department to work overtime and fill in the position of the string as well. The final movie, M_{10} say, should satisfy the producer as representing a solution to the puzzle.

As usual, it helps if the whole process is only half-magic. What must actually happen in this sequence of movies is that the excursion of the ring through the hexagon is gradually replaced by a pulling up of the central loop of string (Mahomet coming to the mountain). In Fig. 14(b) we show an intermediate movie in which you can hardly tell whether this loop, as it passes through the ring in position 4, is above or below the hexagon. You can therefore solve this puzzle by passing the ring from 1 to 2 to 3 while the loop is *below* the hexagon, then lifting the loop a bit while you slide the ring from 3 to 4 to 5 and drop it again so that you can go from 5 to 6 to 7. Since all these movies can be made with a full-sized ring, this will solve the puzzle.

This argument allows us to extend the idea we noted when introducing the Barmy Braid. Suppose that a puzzle has *any number* of rigid pieces (like our arrow and ring) and some arbitrarily flexible ones (e.g., our string) and you could find a solution if the rigid pieces were made flexible. Then, if the motion of the rigid pieces in your solution can be continuously distorted into a rigidly permissible motion, you can use the Magic Movie Method to solve the original puzzle. In topologists' technical language we are using the *Isotopy Extension Principle*.

Party Tricks and Chinese Rings



Figure 15. Girl Meets Boy.

You must have met the party trick where the boy and the girl have to separate themselves without untying the knots in the string. Usually they have lots of fun stepping through one another's arms without effect before they find the real answer.

Let's look at one of those fists more closely (Fig. 16(a)). With a really magic mirror this looks like (Fig. 16(b)) and the solution is obvious, but as usual it's slightly easier to see what to do if your mirror is only half magic (Fig. 16(c)).



Figure 16. Boy Leaves Girl.


Figure 17(a). Pajamas on Hanger.

One of our wire and string puzzles is very like this. The pajama-shaped frame at the bottom of Fig. 17(a) is made of wire rather than string, but it happens to be just about the shape that a piece of string would need to get to while being taken off. In Fig. 8(d) you'll see there's a similar puzzle, but with an extra piece.

The magic mirror in Fig. 17(b) shows that this puzzle can certainly be solved if the wire pajama shape is replaced by a completely flexible string—once again this funny shape is sufficient to overcome its lack of flexibility.



Figure 17(b). Another Look in the Magic Mirror.

The Chinese rings are an indefinite extension of this principle. The magic mirror method shows that the string in Fig. 18(a) can be taken right off. In the course of doing so it reaches a position like that of the wire loop in Fig. 18(b), and removal of this is the usual Chinese Rings puzzle.



Figure 18. The Chinese (st)ring Puzzle.

Chinese Rings and the Gray Code

Figure 19(a) shows a certain position of a 7-ring Chinese Rings puzzle. We call this position

 $1 \quad 0 \quad 1 \quad 1 \quad 1 \quad 0 \quad 0$

because the rings we've numbered

 $64 \quad 32 \quad 16 \quad 8 \quad 4 \quad 2 \quad 1$

are respectively

on off on on off off

the loop. ("On" means that the ring's retaining wire passes through the loop.) Which positions neighbor this?

You hardly need a magic mirror to see how the state of the rightmost ring, number 1, can always be changed (Fig. 19(b)), showing that our position neighbors

1 0 1 1 1 0 1.

But it also neighbors

 $1 \quad 0 \quad 1 \quad 0 \quad 1 \quad 0 \quad 0 \\$

as well!

To see this, slip ring number 8 up over the end of the loop as suggested by the dotted arrow in Fig. 19(a) and then drop it down through the loop as hinted in Fig. 19(c).

In general the rightmost ring, number 1, can always be slipped on or off the loop, so that

 \dots ? ? ? 0 neighbors \dots ? ? ? 1.



Figure 19. Gray Code and Chinese Rings.

But also a ring can be slipped on or off provided that the ring just right of it is on and all ring right of that are off, so that

 \dots ? 1 1 0 0 0 neighbors \dots ? 0 1 0 0 0.

With these neighboring rules the entire set of 2^n positions in the *n*-ring puzzle form one continuous sequence, which for n = 4 is:

	ring $\#$	8	4	2	1	
state $\#$	8, i.e.	1	0	0	0	is 15 moves from being off,
state $\#$	9, i.e.	1	0	0	1	is 14 moves from being off,
state $\#$	11, i.e.	1	0	1	1	is 13 moves from being off,
state $\#$	10, i.e.	1	0	1	0	is 12 moves from being off,
state $\#$	14, i.e.	1	1	1	0	is 11 moves from being off,
state $\#$	15, i.e.	1	1	1	1	is 10 moves from being off,
state $\#$	13, i.e.	1	1	0	1	is 9 moves from being off,
state $\#$	12, i.e.	1	1	0	0	is 8 moves from being off,
state $\#$	4, i.e.	0	1	0	0	is 7 moves from being off,
state $\#$	5, i.e.	0	1	0	1	is 6 moves from being off,
state $\#$	7, i.e.	0	1	1	1	is 5 moves from being off,
state $\#$	6, i.e.	0	1	1	0	is 4 moves from being off,
state $\#$	2, i.e.	0	0	1	0	is 3 moves from being off,
state $\#$	3, i.e.	0	0	1	1	is 2 moves from being off,
state $\#$	1, i.e.	0	0	0	1	is 1 moves from being off,
and state $\#$	0, i.e.	0	0	0	0	is OFF!

How do we tell how many moves it takes to get all the rings off if we're given only the state number, i.e. the sum of the numbers of the rings that are on? The answer displays a remarkable connexion with nim-addition! When you're in state number n, it will take you exactly

$$n \stackrel{*}{+} \lfloor n/2 \rfloor \stackrel{*}{+} \lfloor n/4 \rfloor \stackrel{*}{+} \lfloor n/8 \rfloor \stackrel{*}{+} \ldots = m$$

moves to get off. For example in state 13 you're just

$$13 + 6 + 3 + 1 = 9$$

moves away. And if you're given a number m, then state number

$$m + |m/2| = n$$

is the one that's just m moves from off. For example

$$9 + 4 = 13.$$

The Tower of Hano"

Let's find the position that's 99 moves from off in the 7-ring puzzle. Because the binary expansions of 99 and $\lfloor 99/2 \rfloor$ are

	1	1	0	0	0	1	1
and		1	1	0	0	0	1
the answer is	1	0	1	0	0	1	0.

How many moves is state

 $1 \ 1 \ 0 \ 1 \ 1 \ 1 \ 1$

from off? The answer is found by the 7-term nim-sum

1	1	0	1	1	1	1
	1	1	0	1	1	1
		1	1	0	1	1
			1	1	0	1
				1	1	0
					1	1
						1
 1	0	0	1	0	1	0,

which is the binary expansion of 74.

In various kinds of control device it's important to code numbers in such a way that the codes from adjacent numbers differ in only one place and the code that appears above, known to engineers as the Gray code, has this useful property. It has also been used in transmitting television signals. However, its connexion with the Chinese Rings puzzle was known to Monsieur L. Gros, more than a century ago. Incidentally, the multiknot Loony Loop is connected with a ternary version of the Gray code.

The Chinese Rings have occasionally been used as a sort of combination lock. In recent years several mechanical and electronic puzzles, completely different in appearance, but employing the same mathematical structure, have appeared on the market.

The Tower of Hanoï

In happier times, Hanoï was mainly known to puzzlers as the fabled site of that temple where monks were ceaselessly engaged in transferring 64 gold discs from the first to the last of three pegs according to the conditions that

> only one disc may be moved at a time, and no disc may be placed above a smaller one.

Figure 20(a) shows the initial position in a smaller version of the puzzle and Fig. 20(b) shows the position 13 moves later.

In this puzzle it's possible to make mistakes, unlike in the Chinese Rings where the only mistake you can make is to start travelling in the wrong direction. However, you won't make too many mistakes if you use discs that are alternately gold and silver and

 \diamond



Figure 20. The Tower of Hanoï.



To find out where you should be after m moves, expand m in binary, and then, according as the total number of discs is

	even	\mathbf{or}	odd,
replace a 1 digit by the ternary number	1	\mathbf{or}	2,
replace a 2 digit by the ternary number	21	\mathbf{or}	12,
replace a 4 digit by the ternary number	122	\mathbf{or}	211,
replace an 8 digit by the ternary number	2111	\mathbf{or}	1222,
replace a 16 digit by the ternary number	12222	\mathbf{or}	21111,
replace a 32 digit by the ternary number	211111	\mathbf{or}	122222,
replace a 64 digit by the ternary number	1222222	\mathbf{or}	2111111,

These ternary numbers, when added mod 3 without carrying, show you what peg each disc should be on. For 13 moves and a 7-disc tower, since 7 is odd and

$\begin{pmatrix} 1\\ +4\\ +8 \end{pmatrix}$	we find the ternary numbers	$ \left\{\begin{array}{c} 2\\ 211\\ 1222 \end{array}\right. $
= 13		0001102,

showing that disc 1 should be on peg 2, discs 4 and 8 on peg 1, and the rest on peg 0 as in Fig. 20(b).

The Tower of Hanoï puzzle and the fable which usually accompanies it were invented by Messieurs Claus (Édouard Lucas) and De Parville in 1883 and 1884.



Figure 21. A Solitaire-Like Puzzle.

A Solitaire-Like Puzzle and Some Coin-Sliding Problems

A little puzzle we came across recently is played in a way very similar to the game of Peg Solitaire, except that the pegs are not removed after jumping. Starting from the position of Fig. 21(a), go to the "opposite" position of Fig. 21(b) jumping only in the N-S and W-E directions. The three special pegs are to move to the three special places.



Figure 22. Swap the Hares and Tortoises.

This is rather like various two-dimensional forms of the familiar Hares and Tortoises (or sheep and goats) puzzle (Fig. 22) in which the animals (you can use coins) have to change places and the permitted moves are as in the game of Toads and Frogs in Chapter 1. Other problems with the same coins are:

- 1. get from Fig. 23(a) to Fig. 23(b) with just 3 moves of 2 contiguous coins (the coins to be slid on the table, remaining in the same orientation and touching throughout);
- 2. the same, but reversing the orientation of each pair of coins as it is moved;
- 3. similar problems, but with more coins;
- 4. form the six coins of Fig. 24(a) into a ring (Fig. 24(b)) with just three moves. At each move one coin must be slid on the table, without disturbing any of the others, and positioned by touching it against just two coins. For example, you might try Fig. 24(c) for your first move, but then you wouldn't be able to slide the middle one out.



Figure 24. Ringing the Changes.

The Fifteen Puzzle and the Lucky Seven Puzzle



Figure 25. Sam Loyd's Fifteen Puzzle.

The most famous sliding puzzle is Sam Loyd's *Fifteen Puzzle* in which the home position is Fig. 25 and the move is to slide one square at a time into the empty space. You are required to get home from the random position you usually find the puzzle in. Nowadays the puzzle is usually sold with pieces so designed that it is impossible to remove them from the base.



Figure 26. The Lucky Seven Puzzle.

A more interesting puzzle is the **Lucky Seven Puzzle**, for which the home state is displayed in Fig. 26 and similar rules apply.

In such puzzles there are certain basic permutations of the pieces that bring the empty space back to its standard position. For the Seven Puzzle you can either move the four discs in the left pentagon in the order 1, 2, 3, 4, 1, leading to the position of Fig. 27(a) or treat the right pentagon similarly, moving 7, 6, 5, 4, 7, leading to the position of Fig. 27(b).



Figure 27. After a Few Moves.

In the first case we have effected the permutation α in which

and in the second case the permutation β in which

We can obviously combine these basic permutations to any extent. For instance, by performing the sequence



Figure 28. Crossing Bridges.

By combining any given permutations in all possible ways we get what mathematicians call a **group** of permutations. Is there an easy way to see which permutations belong to the group of the Lucky Seven Puzzle? Yes! The trick, as always in such cases, is to find some permutations which keep most of the objects fixed. In the case of the Seven Puzzle it seems best to regard the outer edges as forming a complete circle across which there is a single bridge between places 0 and 4 (Fig. 28). In this form the seven discs can be freely cycled round the outer circle (which we hardly count as a move) or else a single disc may be slid across the bridge (remember that in the actual form of the puzzle the bridge is too short for several discs to traverse it at once). It doesn't really matter whether the disc we slide across the bridge goes upwards or downwards, since this has the same effect on the cyclic order, so we'll always slide our discs *downward*.



Figure 29. Swapping Two and Four.

If we think of the puzzle in this way and, starting from the home position, slide discs 2, 4, 2, 4, 2 down the bridge, we reach the position of Fig. 29 in which discs 2 and 4 have been interchanged and all the others are in their original places. Obviously we can interchange any pair of discs which are two places apart round the circle in this way. It's not hard to see how *any* desired rearrangement can be reached by a succession of such interchanges. For instance if we wanted to get

disc	1	2	3	4	5	6	7
to place	$\overline{7}$	6	5	4	3	2	1

we might perform the interchanges of the following scheme

leading to a solution in which 45 discs have crossed the bridge. This method is not very efficient but it has the great advantage of providing an almost mechanical technique by which you can obtain any position. Can you find a shorter solution to the above problem?

All Other Courses for Point-to-Point

The history of the Fifteen Puzzle has been given too many times to bear further repetition here. Exactly half of the

$$15! = 1 \times 2 \times 3 \times \ldots \times 15 = 1\,307\,674\,368\,000$$

permutations (the so called *even* permutations) can be obtained. In technical language, the available permutations form the **alternating group**, A_{15} , whereas for the Lucky Seven Puzzle we have the full **symmetric group**, S_7 , of $7! = 7 \times 6 \times 5 \times 4 \times 3 \times 2 \times 1 = 5040$ permutations.

You can make a puzzle of this type by putting counters on all but one of the nodes of any connected graph and then sliding them, point to point, always along an edge into the currently empty node. We can afford to ignore the *degenerate* cases, when your graph is a cycle, or is made by putting two smaller graphs together at a single node, because then the puzzle is trivial, or degenerates into the two smaller puzzles corresponding to the two smaller graphs.

Rick Wilson has proved the remarkable theorem that for every non-degenerate case but one we get either the full symmetric group (if some circuit is odd) or the alternating group (otherwise). The single exception is the graph of the **Tricky Six Puzzle** (Fig. 30) for which the group consists of all possible Möbius transformations

$$x \to \frac{ax+b}{cx+d} \pmod{5}$$



Figure 30. Rick's Tricky Six Puzzle.

Rubik's Hungarian Cube -- Bűvös Kocka

The Hungarian words actually mean "magic cube". If you're crazy enough to get one of these you'll see that when it comes to you from the manufacturer it has just one color on each face (Fig. 31(a)) but your Hungarian cube is unlikely to stay in this beautiful state because you can rotate the nine little **cubelets** that make up any face (Fig. 31(b)) and so disturb the color scheme. For example, if you complete the turn started in Fig. 31(b), and then turn the top face clockwise you'll arrive at Fig. 31(c). After three more turns the colors are all over the place (Fig. 31(d)) and you'll find it very hard to recover the original arrangement; in other words to get each of the cubelets back into its own **cubicle**, and the right way round.



Figure 31. The Hungarian Magic Cube.

There are really two problems about this elegant little puzzle. The first is how its brilliant designer, Ernő Rubik, can possibly have managed to make all those motions feasible without all the cubelets falling apart. We'll leave that one to you! The other is, of course, to provide a method by which we can guarantee to get home from any position our friends have muddled the cube into.

868

Just How Chaotic Can the Cube Get?

At least there are six permanent landmarks: the cubelets at the centres of the faces always stay in their own cubicles although they may be rotated. We call these the **face cubelets** and have framed them in Fig. 31 (a). No matter how confused your cube looks, you can tell what the final color of each face should be, just by looking at the face cubelet at its centre. So, for instance, in Fig. 31 (d) we call the top face **white** even though only one third of it really is.

So you can work out the home cubicle for any cubelet by just looking at its colors and thinking which faces these belong to. For instance the LWO cubelet * in Fig. 31(d) should end up at \dagger (in our cube the colors opposite R, W, O are L, B, Y). We recommend the nervous novice always to hold the cube with its white face uppermost and then to take a careful note of the color of the bottom face, which we call the **ground** color.

Since the other 20 visible cubes are of two types,

8 **corner cubes**, which have 3 possible orientations in their cubicles, and 12 **edge cubes**, which have 2,

there are at most

 $3^8 \times 2^{12} \times 8! \times 12! = 519\,024\,039\,293\,878\,272\,000$

conceivable arrangements. However, Anne Scott proved that only one-twelfth of this number, namely

 $43\,252\,003\,274\,489\,856\,000$

are attainable.

Chief Colors and Chief Faces

These notions help us keep track of the orientations of cubelets, even when they're not in their home cubicles. We'll call the **chief face** of a *cubicle* the one in the top or bottom surface of the cube, if there is one, and otherwise the one in the right or left wall. The **chief color** of a *cubelet* is the color that should be in the chief place when the cubelet gets home. In other words White or the Ground color if possible, and otherwise the color that should end up in the left or right wall of the cube.

If a cubelet, no matter where it is, has its chief color in the chief face of its current cubicle we'll call it **sane** and otherwise **flipped** (if it's an edge cubelet) or **twisted** (if it's a corner one). There's only one way to make an edge-flip (e), but a corner may be twisted anticlockwise (a) or clockwise (c).

Now, as shown in Fig. 32, turning the top (or bottom) preserves the chiefness of every cubelet. Turning the front (or back) changes the chiefness at four corners and turning the left (or right) changes it at four corners and four edges. Since each turn flips an even number of edges, you can see that for attainable positions

the total number of edge-flips will always be even.





Figure 32. Changes in Chiefness.

And since each turn produces equal numbers of clockwise and anticlockwise twists

the total corner twisting will always be zero, mod 3.

In computing corner twists we count +1 for clockwise and -1 for anticlockwise — of course three clockwise twists of a cubelet produce no effect. Finally, for reasons as in the Fifteen Puzzle

the total permutation of all the 20 movable cubelets must be *even*.

An even permutation is one we might imagine making by an even number of interchanges.



Figure 33. The Six Stages and Our Notation for the Moves.

Curing the Cube

Benson, Conway and Seal have simplified Anne Scott's proof that you really can get home from any position for which

- (i) the total edge flipping is zero, mod 2,
- (ii) the total corner twisting is zero, mod 3, and
- (iii) the total permutation of all 20 movable cubes is even.

We have adapted our names for the moves (Fig. 33) so as to agree with David Singmaster's in the hope that a single notation will rapidly become universal. Note that the unprimed letters L,R,F,B,U,D, refer to *clockwise* turns, and the primed letters L',R',F',B',U',D' to *anticlockwise* ones. Our notation for the **slice moves** is illustrated in Fig. 34. Note that in these moves only the *middle* layer of the cube is turned. We shall also use the common notation in which, for example, X^2 means "do X twice" and X^{-1} means "undo X".



Figure 34. Slice Moves.

Our method has six stages which correspond roughly to the letters in Fig. 33(a).

- A: Aloft, Around (Adjust) and About.
- B: Bottom Layer Corner Cubelets.
- D: Domiciling the Top Edge Cubelets.
- E: Exchanging Pairs of Top Corners.

C: Central Layer Edge Cubelets.

F: Finishing Flips and Fiddles.

We've collected the figures for these stages in Fig. 35 for easy reference, so keep a finger on page 871.

Warning: Be very careful when applying this algorithm. Think of "tightening" or "loosening" a screw-cap, so that you never mistake a clockwise turn for an anticlockwise one, even from behind. Be aware at all times which way up you are holding the cube, and don't stop to think in the middle of a sequence of moves. Remember that if you make a tiny mistake you'll probably have to go all the way back to Stage A.

THE CUBE SELDOM FORGIVES!

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A: Aloft, Around (Adjust) and About

Our first stage (Fig. 35A) gets the bottom edge cubelets (A in Fig. 33(a)) into their correct cubicles, the right way round. You bring the ground (= chief) color of such a cubelet into the topmost surface (Aloft) then turn the top layer Around to put this cubelet into the correct side wall which can be turned About to home the cubelet. Sometimes this disturbs a bottom edge cubelet that's already home, but this can be Adjusted by turning the appropriate side wall just before the About step.

B: Bottom Layer Corner Cubelets

Now, without disturbing the bottom layer edge cubelets, you must get the bottom layer corner cubelets home.

If the cubelet that's to stand on the shaded square of Fig. 35B is in the top layer, turn the top layer until this cubelet's ground color is in one of the three numbered positions. Then do the appropriate one of

```
B1:F'U'F\ B2:RUR'\ B3:F'UF.RU^2R'
```

If the cubelet is already *in* the bottom layer, but wrongly placed, use one of these to put any corner cubelet from the top layer into its current position, thereby evicting it into the top layer. Then work as above to put it into the proper place. Repeat this procedure for the other three bottom layer corner cubelets.

C: Central Layer Edge Cubelets

This stage corrects the central layer edge cubelets without affecting the bottom layer.

If the cubelet destined for the shaded cubicle of Fig. 35C is in the top layer, turn the top layer until you want to move this cubelet in one of the two ways of Fig. 35C (its side face will then be just above the face cubelet of the same color). Then do the appropriate one of

 $C1: URU'R'.U'F'UF\ C2: U'F'UF.URU'R'$

If the cubelet is already *in* the central layer, but wrongly placed, use one of these to evict it into the top layer. Then work as above. Repeat the procedure for the other three central layer edge cubelets.

D: Domiciling the Top Edge Cubelets

i.e. putting the top layer edge cubelets into their own home cubicles without as yet worrying about their orientations.

You can do this by a sequence of swaps of adjacent edge cubelets as in Fig. 35D for which the moves are

UF.RUR'U'.F'

Of course you can first turn the top layer to reduce the number of swaps needed.



Aloft-



Around



(Adjust)?











BI: FUF B2: RUR'

CI:URUR UFUF C2:UFUF.URUR

B3:FUF.RU²R'







Figure 35. Six Simple Stages Cure Chaotic Cubes.

E: Exchanging Pairs of Top Corners

1

Now you must get the top layer corner cubelets into their own cubicles by moves that, when they are finally completed, won't have affected the bottom two layers or moved the top layer edge cubelets. Usually you can do this in just two swaps of adjacent corners, but sometimes four will be needed.

Correct performance requires some care. Work out a pair of successive swaps of adjacent corner cubelets that will improve things. Then turn the cube until the first required swap is as in Fig. 35E and do our

$$\mathbf{monoswap}, \ \mathrm{M_s} = \mathrm{FD}.\mathrm{F}^2\mathrm{D}^2\mathrm{F}^2.\mathrm{D}'\mathrm{F}'$$

Then turn THE TOP LAYER *ONLY* to bring the second desired swap into the position of Fig. 35E, do another monoswap, and then return the top layer to its original position.

Since the bottom two layers are disordered by a single monoswap, but restored by a second one, it's important not to move these layers (by turning the cube, say) between the two mono-swaps of each pair.

F: Finishing Flips and Fiddles

Since every cubelet should now be in its own cubicle, the only remaining problems can be solved by edge-flips and corner twists in the top layer. To tackle any particular top layer cubelet, turn THE TOP LAYER *ONLY* to bring that cubelet into one of the two shaded cubicles of Fig. 35F and then, according as its white face is in position

1, 2 or

anticlockwise monotwist,	clockwise monotwist	or	edge monoflip
$\rm M_a = (F'RFR')^2$	$\rm M_c = (RF'R'F)^2$		$M_e = (\varepsilon R)^4$

where ε is a slice move (Fig. 34).

Once again it's important not to move the bottom two layers by turning the cube between operations, since individual monotwists and monoflips affect these layers. However, the entire set of operations needed to correct the top layer will automatically correct the bottom two layers as well.

Explanations

do our

Stage E works because our monoswap operation M_S leaves the top layer unchanged except for the desired swap of the two near corner cubelets, while two copies of the monoswap cancel $(M_S^2 = 1)$.

So a sequence such as

monoswap, turn top clockwise, monoswap, turn top back

doesn't really disturb the bottom two layers, which "feel" only the two cancelling monoswaps. The top layer, however, effectively undergoes a swap of the two near corners followed by a swap of two right corners, which are brought into position by the first top turn and returned by the second.

3

Improvements

Stage F works similarly because M_a , M_c and M_c have exactly the desired effects on the top layer, and enjoy the properties

$$\mathbf{M}_e^2 = \mathbf{M}_c^3 = \mathbf{1}, \quad \mathbf{M}_c \mathbf{M}_e = \mathbf{M}_e \mathbf{M}_c, \quad \mathbf{M}_a = \mathbf{M}_c^{-1}.$$

So Anne Scott's laws ensure that the bottom two layers feel a cancelling combination of operations, while the top layer undergoes the desired flips and twists.

Improvements

Our method is easy to explain, perform and remember, but usually takes more moves than an expert would. If you're prepared to take more trouble and have a rather larger memory, you can often shorten it considerably. For instance, the original monoflips and monotwists (due to David Seal and David Goto) are shorter:

 $\mathbf{m}_{\mathbf{e}} = \mathbf{R} \varepsilon \mathbf{R}^2 \varepsilon^2 \mathbf{R} \qquad \mathbf{m}_{\mathbf{e}}^{-1} = \mathbf{R}' \omega^2 \mathbf{R}^2 \omega \mathbf{R}' \qquad \mathbf{m}_{\mathbf{c}} = \mathbf{R}' \mathbf{D} \mathbf{R} \mathbf{F} \mathbf{D} \mathbf{F}' \qquad \mathbf{m}_{\mathbf{a}} = \mathbf{m}_{\mathbf{c}}^{-1} = \mathbf{F} \mathbf{D}' \mathbf{F}' \mathbf{R}' \mathbf{D}' \mathbf{R}$

but with these you must always be careful to follow a mono-operation by the corresponding inverse one.

Explore the effects of the following moves, which many people have found useful. The first few only affect the top layer. Here and elsewhere we've credited moves to those who first told them to us. We expect that many facts about the cube were found by clever Hungarians long before we learnt of them. For the Greek letter slice moves, see Fig. 34.

David Benson's "special" RUR2. FRF².UFU² David Singmaster's "Sigma" FURU'R'F' Margaret Bumby's top edge-tricycle $\beta U^{\pm 1}\alpha$. $U^2.\beta U^{\pm 1}\alpha$ Two more top edge-tricycles $U^2F.\alpha U\beta.U^2.\alpha U\beta.FU^2$; FUF'UFU²F'U² Top corner tricycle RU'L'UR'U'LU Clive Bach's cross-swap $(\alpha^2 U^2 \alpha^2 U)^2$ Kati Fried's edge-tricycle $\beta F^2 \alpha F^2$ Tamas Varga's corner tricycle $((FR'F'R)^3 U^2)^2$ Two double edge-swaps $(R^2 U^2)^3$; $((\alpha^2 U^2)^2$ Andrew Taylor's Stage C moves $F^2(RF)^2(R'F')^3$; $(FR)^3(F'R')^2F^2$ Other Stage C moves FUFUF. U'F'U'F'U'; R'U'R'U'R'. URURU

In the Extras you'll find lists of the shortest known words (improvements welcome!) to achieve any rearrangement, or any reorientation of the top layer. These are quoted from an algorithm due to Benson, Conway and Seal which guarantees to cure the cube in at most 85 moves (a half turn still counts as one move, but a slice counts as two). More Thistlethwaite has recently constructed an impressive algorithm which never takes more than 52 moves.

Because there are 18 choices for the first move, but only 15 (non-cancelling) choices for subsequent ones, the number of positions after 16 moves is at most

 $18 \times 15^{15} = 7\,882\,090\,026\,855\,468\,750 < 43\,252\,003\,274\,489\,856\,000$

proving that there are many positions that need 17 or more moves to cure. We can improve this to 18 moves by using the estimates $u_1 = 18$, $u_2 = 27 + 12u_1 = 243$, $u_{n+2} \le 18u_n + 12u_{n+1}$, which take into account relations like LR = RL.

Elena's Elements

Elena Conway likes making her cube into pretty patterns. Here are some ways she does this:

"4 Windows" $\alpha \gamma^2 \beta \delta^2$	"6 Windows" $\alpha \gamma \beta \delta$	"Chequers" $\alpha^2 \gamma^2 \varepsilon^2$	"Harlequin" $\alpha\gamma\beta\delta\alpha^2\gamma^2\varepsilon^2$
"Stripey" $(L^2F^2R^2)^2.LR'$	$^{ m `'Zigzag"}$ $({ m LRFB})^3$	"4 Crosses" (LRFB) ³ (FBLR) ³ or $(\gamma^2 L' \gamma^2 R)^3$	"6 Crosses" $(\gamma^{2}L'\gamma^{2}R)^{3}(\alpha^{2}B'\alpha^{2}F)^{3}$

And try following "6 Crosses" with any of the earlier ones.

Are You Partial to Partial Puzzles?

It's interesting to see what you can do using only *some* of the available moves. You might restrict yourself to just a specified selection effaces, to half-turns, to slice moves, or to the **helislice moves** like LR. Mathematically these correspond to subgroups we call the 2-, 3-, 4- and 5-face groups, the square group, the slice group and the helislice group.

Beginners are recommended to stay in the slice group because they cannot get lost. From any position you can cure the edge-cubes in 3 slices, getting to "4 Windows" or "6 Windows" and so home in 4 more slices. Frank O'Hara has shown that in fact at most 5 slices are needed in all. The slice group has order $4^3.4!/2 = 768$ and the helislice group has order $2^{11}.3 = 6144$.

The 2-face group has been intensively studied by Morwen Thistlethwaite. It's interesting to notice that it involves both the lucky Seven Puzzle (on the edge cubelets that move) and Rick Wilson's Tricky Six Puzzle (on the corners).

Roger Penrose first proved that everything can be done using just 5 faces. David Benson has a simple proof:

$$RL'F^2B^2RL'.U.RL'F^2B^2RL' = D.$$

Other "Hungarian" Objects

A $2 \times 2 \times 2$ cube and $2 \times 3 \times 3$ "domino" have also been manufactured. Their design seems even more mysterious, although as puzzles they're much easier. One can *imagine* Hungarian tetrahedra, octahedra, dodecahedra, icosahedra, etc. Although, as far as we know, these have not all been manufactured nor completely solved, Andrew Taylor has found a neat proof that (for any choice of chief faces and colors)

the total permutation on edges and corners is even, the number of edge-flips is even, and

the total corner twisting is zero, modulo the corner valence.

Despondent Domino dabblers should need but three little words (with effects):

$$\begin{array}{ccc} X = EhEhEh & Y = EcEhNcE & Z = cYcYc \\ (28) & (13)(26)(1'3')(2'6') & (13)(26) \end{array}$$

(c is a clockwise $\frac{1}{4}$ -turn of the top; h,E,N $\frac{1}{2}$ -turns of top, East, North).

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876





Figure 36. Dad's Puzzler.

Dad's Puzzler (Fig. 36) is unfortunately almost the only sliding block puzzle that's generally available from toy stores, although it goes under many different names. The problem is to slide the pieces without lifting any out of the tray, until the 2×2 square arrives in the lower left hand corner. Fifty years ago the puzzle represented Dad's furniture-removing difficulties, and the 2×2 block was the piano; at other times it has been depicted as a pennant, a car, a mountain, or space capsule but the puzzle has remained unchanged, probably for a hundred years. Some more enterprising manufacturer should sell a set containing one 2×2 , four 1×1 and six 2×1 pieces which can be used either for Dad's Puzzler or for the following more interesting puzzles.

In the **Donkey** puzzle the initial arrangement is as in Fig. 37(a) and the problem is to move the 2×2 square to the middle of the bottom row. The name arises from the picture of a red donkey which adorned the 2×2 square in the original French version (L'Âne Rouge, which probably goes back to the last century) but we think that our choice of starting position already looks quite like a donkey's face.



Figure 37. The Donkey and The Century (and a Half).

The **Century Puzzle**, published for the first time in *Winning Ways*, was discovered by one of us several years ago as a result of a systematic search for the hardest puzzle of this size. Start from Fig. 37(b) and, as in the Donkey, end with the 2×2 block in the middle of the bottom row. Or, if you're a real expert, you might try the **Century-and-a-Half Puzzle** in which you're to end in the position got by turning Fig. 37(b) upside-down.

Tactics for Solving Such Puzzles

As in our previous sliding puzzles the basic idea is to see what can be done while quite a lot of the pieces are kept fixed. In all three of these examples one occasionally sees one of the configurations of Fig. 38 somewhere, and any of these can be exchanged for any other, moving



Figure 38. A Micropuzzle.

only the pieces in the area shown. They form a kind of micro-puzzle within the larger one. Figure 40 is a complete "map" of Dad's Puzzler showing how it consists of a dozen of these micro-puzzles joined by various paths of moves that are more or less forced. Using this map, you'll find it easy to get from anywhere to anywhere else.



Figure 39. Micro- and Mini-puzzles Found in Donkey and Century.

P

b

A

E

H

START

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A ->-

D

E

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11-

F



Figure 40. Map of Dad's Puzzler.



Figure 41. Map of the Donkey.

In the Donkey and Century puzzles there are several micro- and mini-puzzles: see what moves you can make inside the regions shown in Fig. 40. The Century and Donkey puzzles will never become easy but it will help if you become an adept at these minipuzzles. Figure 41 is our map of the Donkey. The positions are classified according to the location of the 2×2 square and in most cases we have only drawn one of a left-right mirror-image pair. Some unimportant culs-de-sac will be found in the directions indicated by the signs $-|||_{\bullet}$, and the rectangle containing (?) represents many positions connected to the left and right shoulders. The arrows indicate other connexions to the shoulders. Left-right symmetric positions are boldly bordered.

The Century puzzle is very much larger, and we need more abbreviations to draw its map within a reasonable compass. The positions are best classified by the position of the large square together with information about which of the two horizontal pieces should be counted as "above" or "below" the square. We remark that in Fig. 42 both horizontal pieces should be counted as *below* the square despite their appearance, because the only way to move these pieces takes the horizontals *down* and the square *up*.



Figure 42. The Two Horizontal Pieces are Below the Square!

The key to the puzzle is to find one of the two possible **narrow bridges** in the map at which the first horizontal piece changes from *below* the square to *above* it. In fact it's best to think out the possible configurations in which this can happen and then work the puzzle backwards and forwards from one of these. Very few people have ever solved the puzzle by starting at the initial configuration and moving steadily towards its end. A much abbreviated map appears as Fig. 43.

Our maps were prepared with much help from some computer calculations made by David Fremlin at the University of Essex, who found incidentally that the Donkey pieces may be placed in the tray in 65880 positions and the Century pieces in 109260 ways. Although the Century puzzle can be inverted (this is our Century-and-a-Half problem) Fremlin's computer found that the Donkey cannot. It would be nice to have a more perspicuous proof of this.



Figure 43. Map



The starting position is heavily outlined; see centre column of opposite page, near top

Positions are classified thus:

overleaf: Square "between" horizontals. opposite: Two verticals left, one right. this page: Three verticals left.

They are further classified according to the location of the square:



The "narrow bridges" are the two thick connexions between the top of the opposite page and the left of overleaf.

..... of the



Counting Your Moves

It's customary to follow Martin Gardner and declare that any kind of motion involving just one piece counts as a single move. It takes 58 moves to solve Dad's Puzzler and 83 to solve the Donkey. How many do you need to solve the Century puzzle? And how many for the Century-and-a-Half?

Paradoxical Pennies

You tell me your favorite sequence of three Heads or Tails and then I'll tell you mine. We then spin a penny until the first time either of our sequences appears as the result of three consecutive throws. I bet you 2 to 1 it's mine!

The graph



shows the sequence I'll choose for each possible sequence of yours, together with the odds that I win. You'll see that it's always at least 2 to 1 in my favor.

Here's a rule for computing the odds. Given two Head-Tail sequences a and b of the same length, n, we compute the **leading number**, aLb, by scoring 2^{k-1} for every positive k for which the last k letters of a coincide with the first k of b. Then we can show that the odds, that b beats a in Paradoxical Pennies, are exactly

$$aLa - aLb$$
 to $bLb - bLa$.

Leo Guibas and Andy Odlyzko have proved that, given a, the best choice for b is one of the two sequences obtained by dropping the last digit of a and prefixing a new first digit. Notice the paradoxical fact that in the length 3 game:

THH beats HHT beats HTT beats TTH beats THH.

Paradoxical Dice

You can make three dice, A, B, C, with a similar paradoxical property, using the magic square:

	D	Е	F
Α	6	1	8
В	7	5	3
С	2	9	4

Each die has the numbers of one row of the square on its faces (opposite faces bearing the same number). For these dice

A beats B beats C beats A,

all by 5 to 4 odds! Similarly for the three dice, D, E, F, obtained from the columns. The only other paradoxical triples of dice using the same numbers are those obtained from A, B, C by interchanging 3 with 4 and/or 6 with 7. These interchanges improve the odds.

It's possible to put positive integers on the faces of two dice in a unique non-standard way that gives the same probability for each total as the standard one. Algebraically, the problem reduces to factorizing

$$x^{2} + 2x^{3} + 3x^{4} + 4x^{5} + 5x^{6} + 6x^{7} + 5x^{8} + 4x^{9} + 3x^{10} + 2x^{11} + x^{12}$$

into the form f(x)g(x) with f(0) = g(0) = 0 and f(1) = g(1) = 6. The two factorizations are

$$(x + x^2 + x^3 + x^4 + x^5 + x^6)^2$$
 and
 $(x + 2x^2 + 2x^3 + x^4)(x + x^3 + x^4 + x^5 + x^6 + x^8)$

so the new pair of dice have the numbers

$$1, 2, 2, 3, 3, 4$$
 and $1, 3, 4, 5, 6, 8$.

More on Magic Squares

It's an old puzzle to arrange the numbers from 1 to n^2 in an array so that all the rows and columns and both the diagonals have the same sum, which turns out to be $1/2n(n^2 + 1)$. The only 3×3 magic square (see the last section), often called the Lo-Shu, was discovered several dynasties ago by the Chinese. We also used it in Chapter 22. In 1693 Frenicle de Bessy had worked out the 880 magic squares of order 4. In this section we'll show you how to find all these.

It's handy to subtract 1 from all the numbers, because the numbers 0 to 15 are closed under nim-addition. With this convention the magic sum is 30. We shall call a square **perfect** if we can nim-add *any* number from 0 to 15 to its entries and still obtain a magic square; if only 1/2 of these additions are possible we'll call it 1/2-**perfect**, and so on. Since nim-adding 15 is the same as complementing in 15, it always preserves the magic property, showing that *every* square is at least 1/8-perfect. We shall also classify the squares by the disposition of complementary pairs as in Fig. 44.



Figure 44. Classifying Squares by Complementing Pairs.

There are essentially just three ways to write the numbers from 0 to 15 as an addition table:

0	1	2	3	0	1	4	5	0	2	4	6
4	5	6	7	2	3	6	7	1	3	5	7
8	9	10	11	8	9	12	13	8	10	12	14
12	13	14	15	10	11	14	15	9	11	13	15

but you can then freely permute the rows and columns in any of these. Take any table obtained in this way, say

15	11	14	10
13	9	12	8
7	3	6	2
5	1	4	0

Apply the interchanges indicated by our Quaquaversal Quadrimagifier



and you get the magic square on the left in:

15	2	1	12	16	3	2	13
4	9	10	7	5	10	11	8
8	5	6	11	9	6	7	12
3	14	13	0	4	15	14	1
3	14	13	0	4	15		14

Adding I to this particular example we obtain the right hand square which features in Albrecht Dürer's famous self-portrait, *Melencolia I*, in which the boxed figures indicate the date of the work. In this case complementary numbers appear according to the scheme called CD in Fig. 44, and so this square is called Central Diagonal.

By applying the Quaquaversal Quadrimagifier to the other forms of addition table we can get 432 essentially different perfect magic squares. The complementary pairs enable us to classify these as:

- 48 Adjacent Diagonal (AD),
- 48 Broken Diagonal (BD),
- 48 Central Diagonal (CD),
- 96 Adjacent Horizontal (AH) or Adjacent Vertical (AV),
- 96 Broken Horizontal (BH) or Broken Vertical (BV), and
- 96 Central Horizontal (CH) or Central Vertical (CV).

Because we don't count squares as different when they are related merely by a reflexion or a rotation of the diagram, we must regard Adjacent-Horizontal and Adjacent-Vertical squares as the same type. You can find out what type your square will be by looking at the position occupied by the complement of the addition table's leading entry before Quadrimagification:

*	\bigcirc	\bigcirc	0
\bigcirc	AD	\mathbf{CH}	$_{\rm BV}$
0	\mathbf{CV}	BD	AH
0	$_{\rm BH}$	AV	CD

Now take the above 96 Central-Horizontal squares and apply the flip operation

0	\bigcirc	\bigcirc	0
O ↑	0	0	O ↑
ò	0	0	ò
0	0	0	\bigcirc

and you'll get 96 more Central-Horizontal squares. All squares so far found are perfect.

There are 112 more Central-Horizontal squares that are only 1/4 -or 1/8-perfect. They can be found by taking any of the seven squares:

6	10	5	9		14	2	13	1						10	t	14	5
13	12	3	2	$a_{\rm B}$	5	4	11	10					с,*	13	8	7	2
0	7	8	15	~~	8	15	0	7	0	13	2	15	9		15	0	11
11	1	14	4		3	9	6	12	11	8	7	4		3	6	9	12
$\downarrow d$ $\downarrow d$						14	3	12	1				↓ c				
12	5	10	3		13	4	11	2	5	6	9	10	×	12	11	4	3
11	9	6	4	a_1	10	8	7	5					cN	1	8	7	14
0	14	1	15	~~	1	15	0	14						2	5	10	13
7	2	13	8	.	6	3	12	9						15	6	9	_0_
$\frac{1}{4}$ -perfect										<u>1</u> 8	-perf	ect					

and applying any combination of the four operations:



Now take the 14 squares of Fig. 45 and apply any combination of complementation and the last three of our operations and you'll get a total of 224 squares, 56 of each of the types

Adjacent Central (AC), Broken Central (BC), Central Adjacent (CA), and Central Broken (CB).

6	9	4	11	14	1	12	3	2	13	0	15	10	5	8	7					
8	14	1	7 a ₈	0	6	9	15	b 12	6	9	$3 a_8$	4	14	1	11					
3	5	10	12	11	13	2	4	↔ 11	1	14	4	3	9	6	12	14	1	8	7	
13	2	15	0	5	10	7	8	5	10	7	8	13	2	15	0	2	12	3	13	
		d			,	ļ d				↓ d			ļ	l d		4	5 11	4	10	
12	3	8	7	13	2	9	6	4	11	0	15	5	10	1	14	9	6	15	0	
1	13	2	14 a_1	0	12	3	15	9	12	3	$6 a_1$	8	13	2	7			↓d		
6	10	5	9	7	11	4	8	7	2	13	8	6	3	12	9	13	2	1	14	
11	4	15	0	10	5	14	1	10	5	14	1	11	4	15	0	. 4	9	6	11	
	,	d				↓ d				↓ d				↓ d		10	7	8	5	
9	6	1	14	11	4	3	12	8	7	0	15	10	5	2	13	3	12	15	0	
2	11	4	13 a ₂	0	9	6	15	3	9	6	12 a_2	1	11	4	14	-				
12	5	10	3	14	7	8	1	14	4	11	1	. 12	6	9	3					
7	8	15	0	5	10	13	2	5	10	13	2	7	8	15	0					
							1/2-p	erfect									<u>1</u> -pe	rfect		

Figure 45. Adjacent and Broken Central and Central Adjacent and Broken Squares.

There remain only 16, rather irregular, squares to be found. You can get them by applying any combination of complementation and the last two of our operations to the two 1/2-perfect squares

1	14	9	6		2	13	3	12
10	3	4	13		5	6	8	11
7	8	15	0	$\stackrel{d}{\rightarrow}$	14	1	15	0
12	5	2	11		9	10	4	7

and they're 8 each of the types

Adjacent Knighted (AK), Broken Knighted (BK).

There are various permutations of the 16 numbers that occasionally lead from one magic square to another, namely

 a_n : nim-add n, for example $a_6 = (0 \ 6)(1 \ 7)(2 \ 4)(3 \ 5)(8 \ 14)(9 \ 15)(10 \ 12)(11 \ 13)$

b: the big swap $(0\ 12)(1\ 13)(14\ 2)(15\ 3)$

c: circle $(0\ 10\ 12)(1\ 11\ 13)(14\ 4\ 2)(15\ 5\ 3)$

d: double, mod 15 $(1 \ 2 \ 4 \ 8)(3 \ 6 \ 12 \ 9)(5 \ 10)(7 \ 14 \ 13 \ 11)$

and we've indicated some of these in the figures.

The Magic Tesseract



We'll leave it to you to rediscover the many remarkable relations between the 48 BD squares, sometimes called **pandiagonal** or **Nasik** squares, and our **Magic Tesseract** in which the vertices of every square add to 30. By projecting this along three different directions, you can find three magic cubes in which each face adds to 14. These are the duals of the three octahedral dice found by Andreas and Coxeter. Alternate vertices in the magic tesseract are the odious and evil numbers, and if you replace each odious number by its opposite (nim-sum with 15) you'll see how the tesseract was made.

Adams's Amazing Magic Hexagon

Starting from the pattern



can you reorder the numbers from 1 to 19, taking less than 47 years, so that all five rows in each of the three directions have the same sum?

Strip-Jack-Naked, or Beggar-My-Neighbour **1

Another problem that took almost 47 years to solve concerns this old children's game. Each of the two players starts with about half of the cards (held face-down), which they alternately turn over onto a face-upwards "stack" on the table, until one of them (who's now "the commander") first deals one of the "commanding cards" (Jack, Queen, King, or Ace).

After one of these has been dealt, the other player (now "the responder") turns over cards continuously until EITHER **2 a new commanding card appears (when the players change roles **3) or respectively 1, 2, 3, or 4 non-commanding cards have been turned over. In the latter case, the commander turns over the stack and ajoins it to the bottom of his hand. The responder then starts the formation of a new stack by turning over his next card, and play continues as before.

A player who acquires all the cards is the winner and in real games, it seems that someone always does win. The interesting mathematical question, posed by one of us many years ago, was "is it really true that the game always ends?" Marc Paulhus has recently found the answer to be "no!". About 1 in 150,000 games (played with the usual 52 cards) goes on forever.

We are fairly confident that no one person has played the game anything like that number of times, so the chance (with random shuffling) of experiencing a non-terminating game in a lifetime's play must be very small indeed.

Just as surely, however, the total number of times this game has been played by the World's **4 children must be significantly larger than 150,000, so many of them will have been theoretically non-terminating ones. We imagine, though, that in practice most of them actually did terminate because someone made a mistake.

The Great Tantalizer

This is a tantalizing puzzle which surfaces every now and then with a new alias. We've chosen one of the older names. An early American version was the Katzenjammer puzzle, but most recently it has emerged under yet another name, Instant Insanity. The manufacturers seem to be very good at selecting new names, but they never change the underlying puzzle.
The problem



Figure 46. Pieces for The Great Tantalizer.

is to assemble the four cubes of Fig. 46 (in which the outer letters refer to the hidden faces) into a vertical $1 \times 1 \times 4$ tower in which each wall displays all four "colors", A, B, C, D. If you don't go instantly insane on playing with the cubes, you'll probably be greatly tantalized by them.



Figure 47. The Tantalizer Solved?

T.H. O'Beirne seems to have been the first to publish a general way of solving such problems and we think his solution is still the best. Let's imagine the problem solved and concentrate on the North and South walls of the tower (Fig. 47). Then X, Y, Z, T will be A, B, C, D in some order, as will x, y, z, t. Write the four letters A, B, C, D on a piece of paper and join

X to x, Y to y, Z to z and T to t.

What you'll get will probably be a way of joining ABCD into a circuit, but it might perhaps be several circuits which together include each letter just once. For example if

are

ABCDDCAB

we get the single circuit



while if they were

you'd get two circuits of different lengths



There will be a similar circuit, or system of circuits, for the East-West walls. Each of the two systems will contain every vertex just once and have one edge for each cube.

It's now easy to solve the puzzle by drawing the following graph (Fig. 48). The vertices of the graph are the colors A, B, C, D and the *i*th cube yields three edges labelled *i* joining pairs of vertices corresponding to its pairs of opposite faces. All you have to do is to select from this graph the two separate systems of circuits which each use all four numbers and all four vertices just once.



Figure 48. Solving the Tantalizer.

What are the possibilities for such circuit systems in the example? By considering each possibility

for the circuit lengths, you'll rapidly conclude that both systems must consist of a single 4circuit which can only use the letters in the cyclic order ACBD. There is only one way of selecting two such systems without using any edge twice:



894

So the Great Tantalizer has a unique solution (up to reordering the cubes and rotating or inverting the whole tower). You can get it by pushing the cubes of Fig. 46 together left to right and tipping the result on end.

O'Beirne takes as his basic example a five cube puzzle of this type which dates from the first World War (Fig. 49) and uses the flags of the allies Belgium, France , Japan, Russia and the United Kingdom. You might like to check his assertion that this has just two essentially different solutions.



Figure 49. The "Flags of the Allies" Puzzle.

Polyominoes, Polyiamonds and Searching Policy

A domino is made of two squares stuck together, so S.W. Golomb has suggested the words tromino, tetromino, etc. for the figures that can be made by sticking 3, 4, or more equal squares together. He has registered the particular names pentomino (5 squares) and polyomino (n squares) as trade-marks. Unfortunately few of the puzzles that have been proposed have hidden secrets, so they yield to nothing better than trial and error (or systematic search). As Rouse Ball says about Tangrams in early editions of *Mathematical Recreations* and *Essays*, "the recreation is not mathematical and I reluctantly content myself with a bare mention of it".

Here is the type of puzzle that arises. Up to rotations and reflexions there are just 12 pentominoes, for which you'll find our naming system in Chapter 25, with a total area of 60 square units. Which of the candidate rectangles

 $3\times 20 \quad 4\times 15 \quad 5\times 12 \quad 6\times 10$

can be packed with them? Figure 50 shows a way of solving two of these problems at once, and also, if the pieces are regarded as made of five cubes each, of packing the $2 \times 5 \times 6$ box (they



Figure 50. Packing Pentominoes.

will also pack a $3 \times 4 \times 5$ box). Such problems are peculiarly susceptible to idle computers and the 6×10 pentomino rectangle was one of the first to be tackled in this way when C.B. Haselgrove found its 2339 solutions in 1960.

Noting that two equilateral triangles can form a diamond, T.H. O'Beirne has proposed the terms triamond, etc., for figures made from three of more. Counting reflexions as distinct this time we find there are 19 hexiamonds, named in Fig. 51, which will pack into the shape of Fig. 52 in many thousands of different ways. The packing shown in the figure is the most symmetric about the North-South line. (See page 920.) We'd like to see a similarly symmetric one for the East-West line.



Figure 51. The Nineteen Hexiamonds.

This prompts a few remarks about sensible search procedures when solving puzzles or finding strategies for games that may be too large for complete discussion. Even when you have a large computer it's wise to have some idea where to look. Symmetry is usually a valuable consideration. For instance the (nearly) left-right symmetric solutions of the hexiamond



Figure 52. The Most Symmetric Hexiamond Solution?

puzzle admittedly form only a small corner of the space to be searched, but this one is likely to be a profitable one because the constraints on opposite sides of the board are satisfied simultaneously. However, symmetry is not the only consideration. In analyzing a game it's wise to try to find out what the players are really fighting for (the game's hidden secrets). For example the French Military Hunt game on the Small Board is small enough that you can give an exhaustive analysis without needing to understand what's really going on. But when you've discovered that the players are really just fighting over the opposition you can extend the analysis to much larger boards for which a complete analysis would be prohibitive, even by computer.

Many of the analyses in Winning Ways were found in this way. Only when we realized that Dots and Boxes was really more concerned with parity than with box counting were we able to make any headway. And it's impossibly complicated to evaluate a reasonably sized position in Hackenbush Hotchpotch exactly, but we got a head start when we realized that often the atomic weight was the only thing that really mattered. In Peg Solitaire the hidden secret turned out to be the notion of balance represented by α and β in Chapter 23.

Even though polyomino type problems may have no hidden secrets, some people are much better at them than others because they subconsciously search in more likely places. Experienced polyominists don't undo their good work by repeatedly starting from scratch but keep most of the puzzle in place while fiddling with just a few pieces at any time. When they've found one solution, they can usually transform it into others by similar manipulation. For example, from Fig. 50 you can obtain another solution by repacking pentominoes R and S, and in Fig. 52 we can interchange the two (f, h) pairs or rotate the central (A,D,E, j, j')hexagon.

Exercise for Experts: For what values of n can you pack n^2 copies of hexiamond A into a replica of A on n times the scale?

Alan Schoen's Cyclotome



Figure 53. Dissections of 2n-gons into Rhombs, n = 5 and 6.

Alan Schoen is patenting the interesting sequence of puzzles he derived from the well-known dissections of 2*n*-gons into $\binom{n}{2}$ rhombs of angles $\pi k/n$, $1 \le k \le n-1$ (Fig. 53). He takes

one of each of the $\lfloor n/2 \rfloor$ shapes of rhomb and one of each of the shapes you can make by joining two rhombs in every possible way to form a hexagon. The hexagon must not contain a straight angle, since he observes that no packing of rhombs in the 2*n*-gon contains a pair of parallel edges, except those which form the rungs of the "ladders" which run between each pair of opposite sides in every packing. This non-convexity condition is similar to that imposed by Piet Hein in designing the Soma pieces, but here it arises naturally. Reflexions, are not counted as different. This set of rhombs and hexagons (cyclotominoes?) will pack into the original 2*n*-gon. In fact for

> n = 2 3 4 5 and 6 there are 1 1 3 14 and more than 150

essentially different packings. Schoen gave one of us a set of pieces for n = 8 and we were able to assemble them as in Fig. 54. We've numbered the pieces with the values of k, where $\pi k/n$ is the smaller angle of the rhomb. Where two shapes of piece are made from the same pair of rhombs, the one with the straighter reflex angle has its digits in natural order.

Solutions can be obtained from one another much as in O'Beirne's Hexiamond, or as on our Somap. In Fig. 54 the pieces 4 and 22 may be rotated or exchanged with 2 and 24, which in turn can be rotated or reflected. After this exchange, with 2 touching 11 and 34, we have a rotatable decagon, 1,2,34,11,32,3 & 4 of which the last four pieces form a rotatable octagon. As 3 & 4 are contiguous, they will exchange with 34, after which 4 & 1 and 2 & 3 are contiguous, and will swap with 14 and 23. After the original exchange, 2 may instead have two sides in common with 13 and these two will rotate, after which 21 and 12 may be interchanged if 1 & 2 are moved as well. Or again, 2 may touch 23 & 24, so that after the 34 exchange, 2 & 3 will swap with 32, and then 2 & 11 form a symmetric hexagon. And so on and on, yielding well over a hundred solutions.



Figure 54. Schoen's 16-piece 16-gon. A Century or So of Solutions.

MacMahon's Superdominoes

How many pieces are there in a set of cyclotominoes? According as n = 2m or 2m + 1, there are $m^2 - m$ or m^2 hexagons, and m rhombs in either case, so there are m^2 or $m^2 + m$ altogether. You can use sets for a variety of games and puzzles, ranging from Tangramlike pictures (Fig. 55) to quite sophisticated packing problems. It's early to say if these last contain any hidden secrets (though Alan Schoen has noted the one about parallel edges); there's perhaps a better chance since there is more structure in the shapes than there was in polyominoes and polyiamonds.

Many pleasing patterns can be produced: for example, take r^2 sets of pieces and pack them in nesting 2n-gons of side lengths $1, 2, \ldots, r$.

The exponential difficulty of this sequence of puzzles prompts us to add another remark about searching. A typical combinatorial puzzle or search of "size" n takes something like n! trials to complete, and this is much more like n^n than c^n , no matter how big you take c. On the other hand the number of solutions may only be c^n , and while this goes up fast, your chance of finding one of them is only $(c/n)^n$ and this gets very small very fast as soon as n is bigger than c.



Figure 55. Schoen-Shapes Made with a Sixteen Set: Rooster, Serpent and Gosling.

MacMahon's Superdominoes

In his *New Mathematical Pastimes*, MacMahon proposed a different kind of generalized domino, got by dividing a regular polygon into colored triangles. We'll discuss just two examples. If we use just four colors, there are exactly 24 ways of coloring a triangular superdomino, and the standard problem is to pack these into a regular hexagon with an all black perimeter and adjacent colors alike, as in Fig. 56.

In this case it's hard to keep the secret hidden for very long. There are barely enough black edges to go round, and once you've found a suitable arrangement for them the rest is fairly easy.

When we consider the 24 three-colored square superdominoes, with which the usual problem is to make a 4×6 rectangle under similar conditions, the black edge problem is much more subtle. It can be shown that every solution to this problem has a column of four squares in which every horizontal edge is black (the **ladder**). In Fig. 57(a) the ladder occupies the second column and in Fig. 57(b) it occupies the third. In the Extras you'll find every possible configuration for the black edges.



Figure 56. MacMahon's Four-Colored Triangular Superdominoes.



Figure 57. Three-Colored Square Superdomino Solutions Showing the Ladder.

MacMahon's superdomino problems can be made into jigsaw puzzles by using differently shaped edges in place of colors. Thus for the three colors in MacMahon's square problem one can use either the three edge shapes of Fig. 58(a) or those of Fig. 58(b) (which alter the matching condition).



Figure 58. Two Ways of Making a MacMahon Jigsaw Puzzle.



Figure 59. Conway's Christmas Card, 1968.



Figure 60. Paterson's Wrestling Match.

Quintominal Dodecahedra

903

Some years ago one of us sent out a Christmas card (Fig. 59) in the form of a jigsaw puzzle based on Fig. 58(b). The assembly in Fig. 59 is *not* a solution because it contains heads connected directly to hands and necks connected directly to arms. Can you turn it into an anatomically correct solution? Figure 60 is M.S. Paterson's modification of this idea, using another shape system. You must rearrange the pieces so that each wrestler has a properly connected body consisting of one head, one torso, one pair of shorts, two arms and two legs!

Quintominal Dodecahedra

The MacMahon superdominoes with five or more sides have not received much attention, but here's a nice little problem. There are 12 different **quintominoes** if we use five different colors once each and allow turning over. Can you fit them, colors matching, onto the 12 faces of a regular dodecahedron?

The Doomsday Rule

Here's an easy way to find the day of the week for an arbitrary date in an arbitrary year. The day of the week on which the last day of February falls in any given year will be called the **doomsday** for that year. For instance, in year 1000, doomsday (Feb. 29) was a Thursday (THOUSday). Then the following dates in *any* year are all doomsdays:

(the second alternative in leap years), otherwise for even months,

 $Apr \ 4 \quad Jun \ 6 \quad Aug \ 8 \quad Oct \ 10 \quad Dec \ 12 \\$

(the number of the month in the year), and for odd ones,

Mar 3 + 4 May 5 + 4 Jul 7 + 4 Sep 9-4 Nov 11-4

(add 4 for the 31-day, $\mathbf{long},$ months; subtract 4 for 30-day, $\mathbf{short},$ ones). Here's a summary with memos.

Jan	Feb	Mar	Apr	May	Jun	Jul	Aug	Sep	Oct	Nov	Dec
31/32	28/29	7	4	9	6	11	8	5	10	7	12
"last"	last	long 3	${\rm even}\ 4$	long 5	${\rm even}\ 6$	long 7	$\mathrm{even}\ 8$	short 9	$even \ 10$	short 11	even 12

You should get used to finding other doomsdays in each month by changing the given one by weeks or fortnights; for example, since

Jul 11 is a doomsday, so is Jul 4 (Independence Day),

and since

Dec 12 is a doomsday, so is Dec 26 (Boxing Day).

so these are all the same day of the week (Sunday in 1937, for example). On what day of the week was May-Day in the year 1000? May 9, and so May 2, were doomsdays (Thursdays in year 1000), so May 1 was a Wednesday. It's easy to go wrong when adding numbers to days, so we suggest you use our mnemonics

NUN-day ONE-day TWOS-day TREBLES-day FOURS-day FIVE-day SIXER-day SE'EN-day Sunday Monday Tuesday Wednesday Thursday Friday Saturday Sunday

Let's suppose we want Michaelmas Day (Sep 29) in the year 1000: we say Sep 5 (short 9) and so Sep 26 are doomsdays (Thursdays—FOURS-days) so Sep 29 is 3 on FOURS-day = SE'EN-day (Sunday).

To find doomsday for any year in a given century, you should add to the doomsday for the century year,

the number of *dozens* after that year, the *remainder* after this, and the number of *fours* in the remainder.

For example, for the year 1066 we say

$$\begin{array}{c} \text{THOUS} \\ \text{Thurs} \\ \text{FOURS} \end{array} \right\} \text{day, 5 dozen, 6 and 1,} \\ \begin{array}{c} 600 & (\text{remainder}) & (4' \text{s in 6}) \\ 1 & 1 & \pm 2, \mod 7, \end{array} \right.$$

doomsday in 1066 was a TWOS-day, and so the Battle of Hastings (Oct 14) was fought on a

4 on TWOS-day = SIXER-day (Saturday).

Let's do some years in our own century, given that 1900 = Wednesday = TREBLES-day.

Nov 11, 19—18 4 on TREBLES-day, 1 dozen, 6 and 1 = 15-day = ONE-day (Monday).

Of course, whole weeks can be cancelled, so the parentheses in

(4 on TREBLES) 1, (6 and 1)

can be forgotten, making the answer immediate.

In the Julian calendar (as instituted by Julius Caesar) each century was one day earlier than the last, and so

	0	100	200	300	400	500	600
	700	800	900	1000	1100	1200	1300
	1400	1500	1600	1700			
were							
	Sunday	Saturday	Friday	Thursday	Wednesday	Tuesday	Monday.

904

The Doomsday Rule

But in the modern, Gregorian, calendar (as reformed by Pope Gregory XIII)

$\begin{array}{c} 1600 \\ 2000 \end{array}$	$\begin{array}{c} 1700 \\ 2100 \end{array}$	1800 2200	1500 1900
Tuesday	Sunday	Friday	Wednesday

because each century year that is *not* a multiple of 400 drops its leap day, and so is *two* days earlier than the previous one. In practice, remember that 1900 was a Wednesday, and that each step *backwards* to 1800, 1700, 1600 *adds* two days.

Thus, since Jul 4 is a doomsday,

was

are

Jul 4, 17—76

exactly Sunday, \mathscr{G} dozen, 4 and \mathscr{Y} = Thursday.

Various countries adopted the Gregorian reform by omitting various days; for example,

in Italy, France and Spain, in Britain and the American colonies, elsewhere, various dates between 1583 (Poland) and 1923 (Greece).

You should also remember that the start of the year has not always been Jan 1. For some time before 1066 it was Christmas Day of the previous year, and for several centuries it was Mar 25 (so called Old Style dating, which was abolished in 1752). Such things are ignored in the Doomsday Rule, but, along with varying national conventions, must be accounted for in subtle examples:

Apr 23, 1616 (England) = 2 off Friday, 1 dozen, 4 and 1 = Tuesday (Shakespeare's deathday), Apr 23, 1616 (Spain) = 2 off Tuesday, 1 dozen, 4 and 1 = Saturday (Cervantes' deathday), Feb 29, 1603 (England) = exactly Friday, 0 dozen, 4 and 1 = Wednesday (Whitgift's deathday).

This "1603" must obviously be 1604 (New Style). Archbishop Whitgift was Queen Elizabeth's "worthy prelate" and first chairman of the commission which eventually produced the Authorized Version of the Bible.

The ambiguous days from Jan 1 through Mar 24 in years between about 1300 and 1752 were usually written in the "double dating" convention; e.g. Queen Elizabeth's deathday was Mar. 24, 1602/3 for which we find "3 on Fri + 3" = Thursday.

When calculating a B.C. date, it's best to add a big enough multiple of 28 (or 700) years to make it into an A.D. one, remembering that there was no year 0 (1 B.C. was immediately followed by 1 A.D.). Thus, in the Julian system we add 4200 to

Oct 23, 4004 B.C., getting Oct 23, 197 A.D. (not 196),

and giving

1 off SIXER-day, 8 dozen, 1 (and 0) = SE'EN-day = Sunday

for the day of Creation, according to Archbishop Ussher.

Problem: 1. On what weekday is the 13th of the month most likely to fall in the Gregorian calendar?

... and Easter Easily

A number of sources give more or less complicated rules for determining Easter. These usually apply only over limited ranges and are sometimes incorrect, even in reputable works, because they neglect the exceptions in the simple rule below.

Easter Day is defined to be the first Sunday strictly later than the **Paschal full moon**, which is a kind of arithmetical approximation to the astronomical one. The Paschal full moon is given by the formula

 $(Apr 19 = Mar 50) - (11G + C)_{mod 30}$

except that when the formula gives

you should take Apr 18

and when it gives

Apr 18 and $G \ge 12$, you should take Apr 17.

In the formula,

G(the Golden number)	 $Year_{mod 19} + 1$ (never forget	to add the 1!)
C(the Century term)	 + 3 for all Julian years	
	-4 for 15xx, 16xx)
	-5 for 17xx, 18xx	Gregorian
	-6 for $19xx$, $20xx$, $21xx$	J

The general formula for C in a Gregorian year Hxx is

Apr 19

$$-H + |H/4| + |8(H + 11)/25|$$
.

The next Sunday is then easily found by the Doomsday rule. Example

 $1945 \equiv 7$, mod 19 so G = 8 and we find for the Paschal full moon: Mar 50 - (88 - 6)_{mod 30} = Mar 50 - 22 = Mar 28.

Because this is a Doomsday, it's very easy to work out that it is

"exactly Wed (+3 + 9 + 2)".

Easter Day, 1945, was therefore Mar 32, April Fool's Day. For 1981 ($\equiv 5$, mod 19) the formula gives

Apr
$$19 - (66 - 6)_{\text{mod } 30} = \text{Apr } 19,$$

so the Paschal full moon is

Apr 18 = Doomsday, 1981 = Saturday,

so Easter Sunday, in 1981, was Apr 19.

Here is an example in the Julian system:

1573: P.F.M. = Mar $50 - (176 + 3)_{\text{mod } 30}$ = Mar 50 - 29 = Mar 21 = Saturday,

so Easter Day, 1573 was Mar 22. Since this date is still in the Old Style 1572, we can say that that year contained two Easters!

You should use the Julian system even today if you want to know when the Orthodox churches celebrate Easter. Example:

Julian P.F.M. $1984 = Apr 19 - (99 + 3)_{mod 30} = Apr 7.$

The next Doomsday is Apr 11, which is, still in the Julian system,

Tuesday, 7 dozen = Tuesday, (Julian 1900)

so that Orthodox Easter Day, 1984 is the Julian date, Apr 9. Since the Julian calendar is now 13 days out of date, this is Apr 22 in the Gregorian system.

Differences between Julian and Gregorian dates:

15xx,	16xx,	17xx,	18xx,	19xx,	20xx,	21xx,	
10 days,	10 days,	11 days,	12 days,	13 days,	13 days,	14 days,	

How Old is the Moon?

If you stand on the earth and watch the sun and moon going round you, you'll see that they take about $365\frac{1}{4}$ [$365 \cdot 242199$] and 30 [$29 \cdot 530588$ or $29\frac{5}{9}$] days to do so, on average [brackets like these contain better approximations to various numbers].

From these facts you can deduce that the number of days that have passed since the last new moon is approximately:

(day number) + (month number) + (year number) + (century number),

all reduced mod 30 $\left[29\frac{5}{9}\right]$.

The **day number** is the number of the day in the month.

```
The month number
```

for	Jan	Feb	Mar	Apr	May	Jun	Jul	Aug	Sep	Oct	Nov	Dec
					5							
	$[2\frac{2}{3}]$	4	$2\frac{1}{3}$	$3\frac{9}{8}$	$4\frac{4}{9}$	6	$6\frac{5}{9}$	8	$9\frac{5}{9}$	$10\frac{1}{9}$	$11\frac{5}{9}$	$11\frac{8}{9}$]

(or just remember that the rule is about $\frac{1}{2}$ a day late/early in the long/short odd months). The **year number** for a year whose last two digits are congruent, modulo 19,

$_{\rm to}$	0	± 1	± 2	± 3	± 4	± 5	± 6	± 7	± 8	± 9
			± 22							
	[0	$\pm 10\frac{8}{9}$	$\pm 21\frac{7}{9}$	$\pm 3\frac{1}{9}$	± 14	$\pm 24\frac{8}{9}$	$\pm 6\frac{2}{9}$	$\pm 17\frac{1}{9}$	± 28	$\pm 9\frac{1}{3}$]

[with an additional

in years 4n (after leap day) 4n+1 4n+2 4n+3 4n+4 (before leap day)].

The century number for the Gregorian centuries

and, for the Julian centuries

To remember these,

the day number is easy,

the month number also, except for Jan = 3, Feb = 4.

the year number's tens digit is its units digit reduced, modulo 3,

the centuries 14xx and 19xx are +1 and -4; and a short century (36524 days)

drops back $5\frac{1}{3}$ days, while a long century (36525 days) drops back $4\frac{1}{3}$ days (because 1273 lunations take $36529\frac{1}{3}[36529\cdot 337]$ days).

Thus (using only the rough numbers) on Christmas Day, 1984, the moon will be

 $25 + 12 + (+28) - 4 \pmod{30} = 1$ day old,

since $84 \equiv +8$, mod 19 and $8 \equiv 2$, mod 3. But, applying the formula to New Year's Day 1985 we find

$$1 + 3(!) + (+09) - 4 = 9$$
 days old

despite the interval of exactly 7 days. The true motion of the moon is very complicated, and such a simple rule can only hope to give answers to within a day or so. If you're watching the moon late at night, for instance, remember that 11:00 p.m. is nearer tomorrow than today because the rule is attuned to the start of the day.

Of course a moon's age of about

 $0, 7\frac{1}{2}, 15, 22\frac{1}{2},$

days corresponds to

New Moon, First Quarter, Full Moon, Last Quarter.

Those who like to keep mental track of the moon throughout a year should remember the total number for that year, e.g.

in 1998, day number + month number -1, in 1999, day number + month number + 10 etc.

2000,	2001,	2002,	2003,	2004,	2005,	2006,	2007,	2008,	2009,	2010
-8,	+3,	$\pm 15,$	-4,	+7,	-12,	-1,	+10,	-9,	+2,	+13

Jewish New Year (Rosh Hashana)

Here's how to calculate the date of the Jewish New Year (Rosh Hashana) for and year Y = 1900 + y in the range 1900 (inclusive) to 2100 (exclusive). You first compute $F = (12G)_{mod 19}$ and then use the formula (in which you usually ignore the fraction by which it exceeds an integer):

$${\rm Sep}\ 6 \,+\, \frac{3}{2}F \,+\, \frac{F+1}{18} \,+\, \frac{1}{4}y_{\rm mod4} \,-\, \frac{2y}{630}$$

However it must be postponed from any

			TUE	MON
SUN	WED	\mathbf{FRI}	if fraction $\geq .633$	if fraction $\geq .898$
			and if $F > 6$	and if $F > 11$
to the followi	ng			
MON	THU	SAT	THU (not WED)	TUE

For years outside this range, you should replace

Sep 6 by Sep 6 $\{[Y/100] - [Y/400] - 9\}$ in Gregorian years, and by Aug 24 in Julian.

$$\frac{-2y}{630} \text{ by } -\frac{2(Y-1900)-1}{630} - \frac{F}{760} = \frac{11(2100-Y)+7(F-1)}{3447360}$$

.633 by $\frac{1367}{2160} = \cdot 63287037..., \quad \cdot 898 \text{ by } \frac{23269}{25920} = \cdot 897723765...$

(the last three replacements seldom affect the answer).

Extras

Blocks-in-a-Box

The key to this puzzle is that every piece except the three $3 \times 1 \times 1$ rods occupies as many "black" cells as "white" in every layer. The rods must therefore be arranged so as to correct the color compositions in all fifteen layers simultaneously. It turns out that there is a unique arrangement which does this. Figure 61 also shows the only three dispositions for the $2 \times 2 \times 2$ cube and $2 \times 2 \times 1$ square. With these five pieces in place, the puzzle becomes easy.



Figure 61. Were You Able to Fit the Blocks-in-a-Box?

A much harder puzzle is to pack 41 $1 \times 2 \times 4$ planks (together with 15 $1 \times 1 \times 1$ holes) into a $7 \times 7 \times 7$ box (see reference to Foregger, and to Mather, who proves that 42 planks can't be packed.)

The Somap

The Soma pieces 1 = W, 2 = Y and 4 = O, while themselves symmetrical, may appear on the surface of the cube in either the *dexter* fashion





The diamond's gory secrets are seven seas away!

Figure 62. The ...



... Somap.

or the *sinister* one



so you can tell which of these pieces are dexter by giving the sum of their numbers, which we call the **dexterity** of the solution. The symbols

DC DC DC na nb nc

refer to different solutions having deficient piece D, central piece C and dexterity n, a single capital letter indicating that the same piece is both deficient and central. Thus

RL	RL	RL	RL
5a	5b	5c	5d

are four solutions in which Red is deficient, bLue is central and pieces 1 and 4 are dexter (1 + 4 = 5), while

 $\begin{array}{cccc} B & B & B \\ 6a & 6b & 6c & \dots \end{array}$

are solutions in which Black is deficient and central while 2 and 4 are dexter.

Along with the solutions in Fig. 62, there are their reflexions whose names are found by interchanging R and L and replacing n by

in the cases	3-n,	6-n,	7-n,
In the cases	O central,	W central,	otherwise.

When two solutions are related by changing just two pieces, P and Q, this is indicated by a solid line PQ. Some three-piece changes are indicated by dashed lines in a similar way. So all that's left for you to do is to find a suitable solution which you can locate on the Somap, and this will then lead you to all the others except R7d.

Solutions to the Arithmetico-Geometric Puzzle

Figure 63 shows how we indicate layers in this puzzle by using a or α , according to orientation, for an a-high block, etc. The 21 solutions to Hoffman's puzzle are exhibited in Table 1 in this notation. When, as usual, only the middle layer is shown, another layer is separated from it by a letter S, and the remaining one is the special layer of Fig. 63. The meanings of the other letters in Table 1 are:

- R: reflect the special layer across the dotted diagonal,
- S: swap the two non-special layers,
- S': swap two adjacent layers in a different direction,
- T: tamper with a $2 \times 2 \times 2$ corner, not involving the special layer,
- T': tamper with a $2 \times 2 \times 2$ corner, which does involve the special layer.



Figure 63. The Special Layer.

We'll leave it to you to work out why this gives just 21 solutions, and to verify that of these, exactly 17 have **duals**, obtained by replacing the dimensions a, b, c by c, b, a. Just one of the solutions (which?) is self-dual. This solution has the remarkable property that it can be repeatedly transformed (into rotations of itself!) by transporting either of two special faces to the opposite side.

Raphael Robinson and David Seal have found ways of combining solutions to the Arithmetico-Geometric puzzle in various dimensions to produce higher-dimensional ones. For example, if

$$a = a_1 + a_2 + a_3$$
 and $b = b_1 + b_2 + b_3$

we know how to pack 27

 $a_1 \times a_2 \times a_3$ or $b_1 \times b_2 \times b_3$

blocks into an

 $a \times a \times a$ or $b \times b \times b$

cube. The Cartesian product of these gives us a way of packing $27^2 = 729$

 $a_1 \times a_2 \times a_3 \times b_1 \times b_2 \times b_3$

6-dimensional hyperblocks into a single

$$a imes a imes a imes b imes b imes b$$

hyperblock. But now the Cartesian product of three copies of Fig. 7 gives us a way to pack $4^3 = 64$ of these

$$a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} b \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm} a \hspace{0.1 cm} imes \hspace{0.1 cm} a \hspace{0.1 cm}$$

hypercube.

In general the method combines m-dimensional and n-dimensional solutions to give an mn-dimensional one. We hope Omar will tell us how to deal with dimensions 5, 7, 11 and so on.



Table 1. The 21 Solutions to Hoffman's Puzzle.

and	One for	r "Th	ree'' T	00!						
	0	0	2		1	1	0	1	0	0
	0	0	2		1	2	1	1	2	2
	2	2	1		0	1	1	0	2	2
There's	only one c	ther s	olution	Hint	bhe ·	$x \pm u$	± ~			

There's only one other solution. Hint: add x + y + z.

Hares and Tortoises

Make the moves in this order (jumps are bold):

$H, \mathbf{T}, T, \mathbf{H}, \mathbf{H}, H, \mathbf{T}, \mathbf{T}, \mathbf{T}, H, \mathbf{H}, \mathbf{H}, T, \mathbf{T}, H.$

If you move only one kind of animal for as long as you can before moving the other kind, you'll soon see how to swap 57 Hares with 57 Tortoises.

Solutions to the other coin problems (heads are **bold**) are:

Start from 012 345 ;	move	01 to 67, 56 to 89 and 23 to 56;
	or	01 to 76, 23 to 98 and 56 to 65 .
Start from 01234567;	move	12 to 89, 45 to 12, 78 to 45 and 01 to 78;
	or	67 to 98, 01 to 76, 34 to 43 and 78 to 87.

M. Delannoy has shown that the first problem with n pairs of coins can always be solved in just n moves. However the second problem, due to Tait, requires n + 1 moves if n > 4. For some reason which we don't understand, we have always found these little problems confusing and can never remember their solutions!

The last little coin puzzle is one of the simplest examples we know of a psychological block. You notice that four coins are already in position (Fig. 64(a)), so you're reluctant either to move one of them (Fig. 64(b)) or to waste time by replacing it (Fig. 64(c)), but that's the only way to get to Fig. 64(d) in three moves. There's a four-move version in which you start with a triangle.



Figure 64. How to Infuriate Your Friends.

The Lucky Seven Puzzle

The Lucky Seven Puzzle has a solution in which just seven discs are slid down the bridge, alternately from the left and right sides:

Top Face Alterations for the Hungarian Cube

We give the shortest known sequences for all permutations (Table 2) and for all combinations of flips and twists (Table 3) in the top layer. The numbers are the numbers of moves, but not counting any final top turns (U^k) which can all be saved to the end. David Seal has proved that most of these are best possible.



Table 2. Top Layer Permutation Sequences. (the lower sequence of a pair refers to the left-right reflected picture)

	0	no moves required	c ● c ● ● 0	12	RU'L'UDB ² D'R ² U ² LU'RU' B'UF'U ² B ² DL ² U'D'FUB'U
••• • e • e ●	12	$F'U'F^2DRUR'D'U'F^2U^2FU'$	с ● с ● е ● е с	11	LR'U'R ² B'R'B ² U'B'U ² L' LU ² BUB ² RBR ² URL'
• e • • •	13	LF'UL'FB'UR'FU'RF'BU'	сес е ● ●● с	11	F'U ² R'U'R ² B'R'B ² U'B'F F'BUB ² RBR ² URU ² F
● e ● ● e ●	13	$L^2F^2L^2U^2R'LFL'RU^2L^2F^2L^2U$	с ● с € ● ● е с	14	$\frac{F^2R^2FU'F^2R'F'R^2U'R'U^2F^2R^2F^2}{F^2R^2F^2U^2RUR^2FRF^2UF'R^2F^2}$
•• a • • •• c	12	R'U'LU ² R'F ² RF ² U'RU ² L'U ² RU ² L'UB ² L'B ² LU ² R'ULU ²	сес ● е ●●¢	8	B'U'B ² L'B'L ² U'L'U ² RUR ² FRF ² UFU ²
●● a ● e ● e c	13	RU'LU ² R ² F'U'FUR ² U ² R'L'U BFU ² F ² U'L'ULF ² U ² B'UF'U'	сес • • • сс	12	$\begin{array}{l} BLU^2B'U'B^2L'B'L^2U'L^2B'\\ BL^2UL^2BLB^2UBU^2L'B' \end{array}$
● e a ● c ● € c	13	R'UL'U ² R ² BUB'U'R ² U ² RLU' F'B'U ² B ² ULU'L'B ² U ² FU'BU	c● c e e ●● c	12	R'F ² U'F ² R'F'R ² U'R'U ² FR R'F'U ² RUR ² FRF ² UF ² R
●● a e ● ● e c	14	F ¹ DF'UFD'FL'U'LUF ² U ² F'U LU ² L ² U'B'UBL'DL'U'LD'L ² U'	сес е е •ес	12	L'U'B'UBLBLUL'U'B'U ² FURU'R'F'R'F'U'FURU ²
● e a e ● ● ● C	14	B ² D'BU'B'DB'LUL'U'B ² U ² BU' L'U ² L ² UFU'F'LD'LUL'DL ² U	a ● c ● ● a ● c	14	$R^2B^2R^2U'RL'DL^2U^2LD'L^2U^2R'U$
● c a ● ● ● c c	14	$RUR^2F^2D'R^2BL'B'R^2DF'RF'U^2$ BL'BD'L ² FRF'L ² DB ² L ² U'L'U ²	a• c • e a e c	12	LUFU'F'L'R'U'F'UFR R'F'U'FURLFUF'U'L'
●● a e e ●● c	14	LF'D'L'BL'B ² U ² L'BDF'L ² F ² U ² B ² R ² BD'F'RU ² F ² RF'RDBR'U ²	a e c ● e a ● c	12	RBUB'U'R'L'B'U'BUL L'U'B'UBLRUBU'B'R'
● c a c = c ● c c	15	$\begin{array}{l} BU^2BR^2FD^2FLFL^2F^2D^2FR^2B^2U\\ R^2F^2LD^2L^2B^2L'B'L'D^2L'F^2R'U^2R'U'\end{array}$	acc ● ● aec	14	F'U'F ² UL'U'RLU'R'UF ² U'F'U ²
a ● ● ● ● c	11	$R'BD^2B'RU^2R'BD^2B'RU^2$	a ● c e e a ● c	14	R'U'F'L'R'U'F'UFRLUFR
a●● ● e ● e c	14	B'UFU'BU ² F ² L'U'LUF ² U ² F' LU'R'UL'U ² R ² BUB'U'R ² U ² R	aec e c aec	13	$\mathrm{LU}^{2}\mathrm{F}\mathrm{U}\mathrm{T}^{2}\mathrm{L}\mathrm{F}^{2}\mathrm{L}'\mathrm{F}^{2}\mathrm{U}\mathrm{F}\mathrm{U}^{2}\mathrm{L}'\mathrm{U}^{2}$
a●● ● ● c	13	$\begin{array}{l} L^{\prime}UBL^{\prime}D^{\prime}BD^{2}R^{2}D^{\prime}B^{2}L^{2}UF^{2}U^{2}\\ B^{2}U^{\prime}R^{2}F^{2}DL^{2}D^{2}F^{\prime}DRF^{\prime}U^{\prime}RU^{2} \end{array}$	c ● a ● ● a ● c	15	R'U ² RU ² R ² B'D'R',F.R ² F'DBU'R'U'
a é ● ● ● ● e c	13	$R'B^2F'L'DF'L^2FD'LFB^2RU^2$	с•а • с асс	15	R'F ² U'D'L'F'LFDULFL'FR R'F'LF'L'U'D'F'L'FLDUF ² R
a•• e e ••c	13	BR ² LFD'LF ² L'DF'L'R ² B'U ²	cea ● ● aec	14	$\begin{array}{l} R^{\prime} UB^{\prime}R^{2}U^{\prime}R^{U}R^{2}BU^{2}R^{2}B^{\prime}R^{\prime}BU^{2}\\ F^{\prime}LFL^{2}U^{2}F^{\prime}L^{2}U^{\prime}L^{\prime}UL^{2}FU^{\prime}LU^{2}\end{array}$
a e ● e c ● e c	15	$F^{\prime}L^{2}D^{\prime}B^{2}D^{\prime}R^{2}B^{\prime}R^{\prime}F^{\prime}R^{2}D^{\prime}F^{2}D^{\prime}L^{2}B^{\prime}U$	сеа е е аес	14	B'D²+U²RF'D²BLU²FU²F'LU

Table 3. Top Layer Flip and Twist Sequences. (the lower sequence of a pair refers to the picture with a and c swapped)

The Century Puzzle

The Century Puzzle is so called because it takes exactly 100 moves, and the Century-and-a-Half takes 151 according to the official rules, but since the first and last are only half-moves, we can obviously count it as 150 whole moves. You can see solutions to both puzzles in Fig. 65, and by turning the book upside-down you'll see the only other 100-move solution to the Century.



Figure 65. Solutions to the Century and Century-and-a-Half Puzzles.

Adams's Amazing Magic Hexagon



In Martin Gardner's Sixth Book of Mathematical Games you can read the remarkable history of Clifford W. Adams's discovery and of Charles W. Trigg's uniqueness proof. It's easy to see that a diameter d magic hexagon uses the numbers from 1 to $(3d^2 + 1)/4$, which add to

$$\frac{1}{2}\left(\frac{3d^2+1}{4}\right)\left(\frac{3d^2+5}{4}\right) = \frac{1}{32}\left(9d^4+18d^2+5\right),$$

so that each of its d columns must add to

$$\frac{1}{32}\left(9d^3 + 18d + \frac{5}{d}\right)$$

which can only be an integer if d divides 5. Frank Tapson has discovered that one William Radcliffe had already found the 'Adams Hexagon' and registered it at Stationers Hall, London, in 1896. We thank Victor Meally for this information.

Flags of the Allies Solution

If you use the O'Beirne method you will find the two pairs of 5-circuits



which lead to the solutions shown in Fig 49 and Fig. 66.



Figure 66. The Other Solution to the Flags of the Allies Problems.

All Hexiamond Solutions Found

In May 1996 Marc Paulhus wrote a program that used only a few days of computer time to find all solutions to O'Beirne's Hexiamond puzzle. Independently, in 1999, Donald Knuth also ran a program (in 5 hours and 21 minutes on a 450 MHz Pentium II) which found one more solution than Paulhus had found. Paulhus reran his program and this time really found them all. The initial run must have contained a machine error!

The numbers of solutions, classified according to how far out the Hexagon appears: A in a corner, B in a side, C, ..., F, G in the centre, are

A 75490 B 15717 C 6675 D 7549 E 11447 F 5727 G 1914 with a total of 124519 solutions, of which Figure 52 is indeed the most symmetrical.

The Three Quintominal Dodecahedra

The three quintominal dodecahedra should be recoverable from



Answer to Exercise for Experts



We have a rather complicated proof that n^2 copies of hexiamond A can be used to replicate A on a larger scale only if $n \equiv 0$ or $\pm 1 \mod 6$. Our proof establishes that these are the only values of n for which the relations (look at the foot of the previous page)

 $U^2D^2 = DUDU$ and $D^2U^2 = UDUD$

imply

$$\mathbf{U}^{2n}\mathbf{D}^{2n} = \mathbf{D}^n\mathbf{U}^n\mathbf{D}^n\mathbf{U}^n.$$

We've also shown that none of the usual kinds of coloring argument excludes other values of n.

Where Do the Black Edges of MacMahon Squares Go?

Round the outside, of course, but there are six more inside. These can be arranged in 20 different ways. In the first two the ladder is in the *third* column, otherwise it's in the second. The last row of Fig. 67 contains 6 + 6 + 2 arrangements: the dotted lines are alternative positions for the sixth black edge.



Figure 67. The Twenty Black Edge Arrangements for MacMahon Squares.

A Doomsday Answer

A tedious enumeration shows that in the 400 years of the Gregorian cycle Doomsday is

	Sun	Mon	Tue	Wed	Thu	Fri	Sat	
for	43	43	43	43	44	43	44	ordinary years
and	13	15	13	15	13	14	14	leap years

From this you can work out that the 13th day falls on

	Sun	Mon	Tue	Wed	Thu	Fri	Sat	
$_{ m in}$	687	685	685	687	684	688	684	months,

verifying B.H. Brown's assertion that the 13th of a month is just a little bit more likely to be a Friday than any other day of the week!

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-25-

What is Life?

Life's not always as simple as mathematics, Abraham! Mrs. Abraham Fraenkel.

Life's too important a matter to be taken seriously. Oscar Wilde.

... in real life mistakes are likely to be irrevocable. Computer simulation, however, makes it economically practical to make mistakes on purpose. If you are astute, therefore, you can learn much more than they cost. Furthermore, if you are at all discreet, no one but you need ever know you made a mistake.

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Most of this book has been about two-player games, and our last two chapters were about one player games. Now we're going to talk about a no-player game, the **Game of Life**! Our younger readers won't have learned much about Life, so we'd better tell you some of the facts.

Life is a "game" played on an infinite squared board. At any time some of the cells will be **live** and others **dead**. Which cells are live at time 0 is up to you! But then you've nothing else to do, because the state at any later time follows inexorably from the previous one by the rules of the game:

- **BIRTH.** A cell that's *dead* at time t becomes *live* at t + 1 only if *exactly three* of its eight neighbors were live at t.
- **DEATH by overcrowding**. A cell that's live at t and has four or more of its eight neighbors live at t will be dead by time t + 1.
- **DEATH by exposure**. A live cell that has only one live neighbor, or none at all, at time t, will also be dead at t + 1.

These are the only causes of death , so we can take a more positive viewpoint and describe instead the rule for

SURVIVAL. A cell that was live at time t will remain live at t + 1 if and only if it had just 2 or 3 live neighbors at time t.

Just 3 for BIRTH 2 or 3 for SURVIVAL

A fairly typical Life History is shown in Fig. 1. We chose a simple line of five live cells for our generation 0. In the figures a circle denotes a live cell .

Which of these will survive to generation 1? The two end cells have just one neighbor each and so will die of exposure. But the three inner ones have two living neighbors and so will survive. That's why we've filled in those circles.



Figure 1. A Line of Five Becomes Traffic Lights.

What about births at time 1? There are three cells on either side of the line that are dead at time 0, but have exactly three live neighbors, so will come to life at time 1. We've shown these prospective births by dots in the figure.

So at time 1 the configuration will be a solid 3 x 3 square. Let's briefly follow its later progress.

- Time 1-2: The corners will survive, having 3 neighbors each, but everything else will die of overcrowding. There will be 4 births, one by the middle of each side.
 - 2-3: We see a ring in which each live cell has 2 neighbors so everything survives; there are 4 births inside.
 - 3-4: Massive overcrowding kills off all except the 4 outer cells, but neighbors of these are born to form:
 - 4-5: another survival ring with 8 happy events about to take place.
 - 5-6: More overcrowding again leaves just 4 survivors. This time the neighboring births form:
 - 6-7: four separated lines of 3, called **Blinkers**, which will never interact again.
 - 7-8-9-10-... At each generation the tips of the Blinkers die of exposure but the births on each side reform the line in a perpendicular direction.
Still Life

The configuration will therefore oscillate with period two forever. The final pair of configurations is sufficiently common to deserve a name. We call them **Traffic Lights**.



Figure 2. If Three Survive, They'll Make a Blinker or a Block.

The Blinker is also quite common on its own (Fig. 2a). Most other starting configurations of three live cells will blank out completely in two moves (Fig. 2(b)). But if you start with three of the four cells of a 2×2 block, the fourth cell will be born and then the **Block** will be stable (Fig. 2(c)) because each cell is neighbored by the three others.

Still Life

It's easy to find other stable configurations. The commonest such **Still Life** can be seen in Fig. 3 along with their traditional names. The simple cases are usually loops in which each live cell has two or three neighbors according to local curvature, but the exact shape of the loop is important for effective birth control.



Figure 3. Some of the Commoner Forms of Still Life.

Life Cycles

The blinker is the simplest example of a configuration whose life history repeats itself with period > 1. Lifenthusiasts (a word due to Robert T. Wainwright) have found many other such configurations, a number of which are shown in Figs. 4 to 8.



Figure 4. Three Life Cycles with Period Two.



Figure 5. Two Life Cycles with Period Three, (a) Two Eaters Gnash at Each Other, (b) The Cambridge Pulsar CP 48-56-72.





Figure 6. A Flip-Flop by the Gosper Group.

The Glider and Other Space Ships

When we first tracked the r-pentomino (you'll hear about that soon) some guy suddenly said, "Come over here, there's a piece that's walking!" We came over and found Fig. 9.

You'll see that generation 4 is just like generation 0 but moved one diagonal place, so that the configuration will steadily move across the plane. Because the arrangements at times 2, 6, 10, ... are related to those at times 0, 4, 8, 12, ... by the symmetry that geometers call a glide reflexion, we call this creature the glider. But when you see Life played at the right

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Figure 7. (a) Catherine Wheel, (b) Hertz Oscillator. Still Life Induction Coils Keep Field Stable.



Figure 9. The Glider Moves One Square Diagonally Each Four Generations.

speed by a computer on a visual display, you'll see that the glider walks quite seductively, wagging its tail behind it. We'll see quite a lot of the glider in this chapter.



Figure 10. (a) Lightweight, (b) Middleweight, (c) Heavyweight Spaceships.

It was at just such a visual computer display that one of us first noticed the **spaceship** of Fig. 10(a) (and was very lucky to be able to stop the machine just before it would have crashed into another configuration). This **lightweight spaceship** immediately generalizes to the **middleweight** and **heavyweight** ones (Figs. 10(b) and (c)) but longer versions turn out to be unstable. It was later discovered, however, that arbitrarily long spaceships can still travel provided they are suitably escorted by small ones (Fig. 11). All the spaceships, as drawn, move Eastwards.



Figure 11. An Overweight Spaceship Escorted by Two Heavyweight Ones.

The Unpredictability of Life

Is there some way to foretell the destiny of a Life pattern? Will it eventually fade away completely? Or become static? Oscillate? Travel across the plane, or maybe expand indefinitely? Let's look at what should be a very simple starting configuration—a straight line of n live cells

n = 1 or 2	fades immediately,			
n = 3	is the Blinker;			
n = 4	becomes a Beehive at time 2,			
n = 5	gave Traffic Lights (Fig. 1) at time 6,			
n = 6	fades at $t = 12$,			
n = 7	makes a beautifully symmetric display before terminating in Honey Farm (Fig. 3) at $t = 14$;			
n = 8	gives 4 blocks and 4 beehives,			
n = 9	makes two sets of Traffic Lights,			
n = 10	10 turns into the pentadecathlon, with a life cycle of 15,			
n = 11	becomes two blinkers,			
n = 12	makes two beehives,			
n = 13	turns into two blinkers,			
n = 14				
$\left.\begin{array}{c}n=14\\\\\text{and}\\n=15\end{array}\right\}$	vanish completely,			
n = 15				
n = 16	makes a big set of Traffic Lights with 8 blinkers,			
n = 17	a = 17 becomes 4 blocks,			
n = 18				
$\left.\begin{array}{c}n=18\\\text{and}\\n=19\end{array}\right\}$	fade away entirely,			
n = 19)				
n = 20	makes just 2 blocks,			
and so on.				

What's the general pattern? Even when we follow the configurations which start with a very small number of cells, it's not easy to see what goes on. There are 12 edge-connected regions with 5 cells (S.W. Golomb calls them pentominoes). Here are their histories:

.



Once again, it doesn't seem easy to detect any general rule.

Here, in Figs. 12 and 13 are some other configurations with specially interesting Life Histories, for you to try your skill with.

Can the population of a Life configuration grow without limit? Yes! The \$50.00 prize that one of us offered for settling this question was won in November 1970 by a group at M.I.T. headed by R.W. Gosper. Gosper's ingenious **glider gun** (Fig. 14) emits a new glider every 30 generations. Fortunately it was just what we wanted to complete our proof that

> Life is really unpredictable!



Figure 12. Exercises for the Reader.



Figure 13. Mainly for Computer Buffs.



Figure 14. Gosper's Glider Gun.

Gardens of Eden

There are Life configurations that can only arise as the initial state, because they have no ancestors!

We'll prove that if n is sufficiently large, there is some configuration within a 5n - 2 by 5n - 2 square that has no parent. It will suffice to examine that part of a prospective parent that lies in the surrounding $5n \times 5n$ square (Fig. 15). If any one of the component 5×5 squares is empty, it can be replaced as in Fig. 15(b) without affecting subsequent generations. So we need consider only

$$(2^{25}-1)^{n^2} = 2^{24 \cdot 999999957004337...n^2}$$
 of the 2^{25n^2}

configurations in the 5n by 5n square. However, there are exactly

$$5^{(5n-2)^2} = 2^{25n^2-20n+4}$$

possible configurations in the 5n-2 by 5n-2 square, so that if

$$24 \cdot 999999957004337 \dots n^2 < 25n^2 - 20n + 4,$$

one of these will have no parent! We calculate that this happens for n = 465163200, so that there is a Garden of Eden configuration that will fit comfortably inside a

2325816000 by 2325816000 square!



Figure 15. Location of the Garden of Eden.

This type of argument was first used by E.F. Moore in a more general context. More careful counting in the Life case has brought the size down to 1400 by 1400. However, using completely different ideas and many hours of computer time the M.I.T. group managed to produce an explicit example (Fig. 16).



Figure 16. An Orphan Found by Roger Banks, Mike Beeler, Rich Schroeppel, Steve Ward, et al.

Life's Problems are Hard!

The questions we posed about the ultimate destiny of Life configurations may not seem very mathematical. After all, Life's but a game! Surely there aren't any difficult mathematical problems there?

Chapter 25. What is Life?

Well, yes there are! Indeed we can prove the astonishing fact that *every* sufficiently wellstated mathematical problem can be reduced to a question about Life! Those apparently trivial problems about Life histories can be arbitrarily difficult!

Here, for instance, is a tricky little problem that kept mathematicians busy from the time Pierre de Fermat proposed it over 300 years ago until Andrew Wiles solved it in the 1990s. Is it possible for a perfect nth power to be the sum of two smaller ones for any n larger than 2? Despite many learned investigations by many learned mathematicians we still don't know! But if you had an infallible way to foretell the destiny of a given Life configuration, you'd be able to answer this question!

The reason is that we can design for you a finite starting pattern P_0 which will fade away completely if and only if there is a way of breaking an *n*th power into two smaller ones. If you had a mechanical method which would accept as input an arbitrary finite Life pattern P, and is guaranteed to respond with

FADE, if the rules of Life will eventually cause P to disappear completely, and STAY, if not,

then you could apply it to P_0 and settle Fermat's question.

Even better, we could design a pattern P_1 which will tell you what those perfect powers are. If

$$a^n + b^n = c^r$$

> is the first solution of Fermat's problem in a certain dictionary order, then eventually P_1 will lead to a configuration in which there are

a gliders, travelling North-West, b gliders, travelling North-East, c gliders, travelling South-West, n gliders, travelling South-East,

and nothing else at all! We can do the same sort of thing for other mathematical problems.

Making a Life Computer

Many computers have been programmed to play the game of Life. We shall now return the compliment by showing how to define Life patterns that can imitate computers. Many remarkable consequences will follow from this idea.



Figure 17. Gliding Pulses.

Good old fashioned computers are made from pieces of wire along which pulses of electricity go. Our basic idea is to mimic these by certain lines in the plane along which gliders travel (Fig. 17). (Because gliders travel diagonally, from now on we'll turn the plane through 45°, so they move across, or up and down, the page.) Somewhere in the machine there is a part called the **clock** which generates pulses at regular intervals and most of the working parts of the machine are made up of logical **gates**, like those drawn in Fig. 18. Obviously we can use Glider Guns as pulse generators. What should we do about the logical gates? Let's study the possible interactions of two gliders which crash at right angles.



Figure 18. The Three Logical Gates.

When Glider Meets Glider

There are lots of different ways in which two gliders can meet, because there are lots of different possibilities for their exact arrangement and timing. Figure 19 shows them crashing (a) to form a blinker, (b) to form a block, (c) to form a pond, or (d) in one of several ways in which they can annihilate themselves completely. This last may seem rather unconstructive, but these **vanishing reactions** turn out to be surprisingly useful!



Figure 19. Gliders Crashin' in Diverse Fashion.

How to Make a NOT Gate

We can use a vanishing reaction, together with a Glider Gun, to create a NOT gate (Fig. 20). The input stream enters at the left of the figure, and the Glider Gun is positioned and timed so that every space in the input stream allows just one glider to escape from the gun, while a glider in the stream necessarily crashes with one from the gun in a vanishing reaction (indicated by *). Figure 20 shows the periodic stream

 $1 \ 1 \ 0 \ 1 \ 1 \ 0 \ 1 \ 1 \ 0 \ \dots$

being complemented to

 $0 \ 0 \ 1$ 0 0 1 0 0 1 110110110... being complemented to 001001001.... Å. Glider * .4 Gun s, ۸. \$ ٨ 1 0 11 ** * 0 1 1 0 ţ ÷ ŝ 꾰 Input <u>8-1</u> . 0 . 0 ×. 1 . 0 . 0 x. 1



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Output-

· 0 . 0 ÷

The Eater

Fortunately there are several vanishing reactions with different positions and timings in which the decay is so fast that later gliders from the same gun stream will not be affected (Fig. 21). This means that we can reposition a glider stream arbitrarily by turning it through sufficiently many corners (Fig. 22).



Figure 21. A Variety of Vanishing Reactions Between Crashing Gliders.

The Eater

What else can happen when glider meets glider ? Lots of things! One of them is to make an eater (Fig. 23) and an **eater** can eat lots of things without suffering any indisposition. The eater, which was discovered by Gosper, will be very useful to us; in Fig. 24 you can see it enjoying a varied diet of (a) a blinker, (b) a pre-beehive, (c) a lightweight spaceship, (d) a middleweight spaceship, and (e) a glider. If it attempts a heavyweight spaceship it gets indigestion and leaves a loaf behind; if it tries a blinker in the wrong orientation it leaves a baker's shop!

Sometimes glider streams are embarrassing to have around, so it's especially useful then it just sits there and eats up the whole stream!



Figure 22. Repositioning and Delaying a Glider Stream.



Figure 23. Two Gliders Crash to Form an Eater.



Figure 24. The Voracious Eater Devours a Varied Meal.



Figure 25. (a) Blockbusting Glider, (b) Glider Dives into Pond and Comes Up With Ship. (c) Glider Crashes into Ship and Makes Part of Glider Gun.



Figure 26. Thirteen Gliders Build Their Own Gun.

Gliders Can Build Their Own Guns!

What happens when a glider meets other things? We have seen it get eaten by an eater. It can also annihilate a block (and itself! Fig. 25(a)). But more constructively it can turn a pond into a ship (Fig. 25(b)) and a ship into a part of the glider gun (Fig. 25(c)). And since gliders can crash to make blocks (Fig. 19(b)) and ponds (Fig. 19(c)), they can make a whole glider gun! The 13 gliders in Fig. 26(a) do this in 67 generations. Figures 26(b,c,d) show the positions after 8, 16 and 44 generations. The extra glider then slips in to deal with an incipient beehive, and by 67 generations (Fig. 26(e)) the gun is in full working order and launches its first glider 25 generations later.

The Kickback Reaction

Yet another very useful reaction between gliders is the **kickback** (Fig. 27(a)) in which the decay product is a glider travelling along a line closely parallel to one of the original ones, but in the opposite direction. We think of this glider as having been *kicked back* by the other one. Figure 27(b) shows our notation for the kickback.



Figure 27. The Kickback.

All the working parts of our computer will be moving glider streams, meeting in vanishing and kickback reactions. The only static parts will be glider guns and eaters (indicated by G and E in the figures).

Thinning a Glider Stream

The glider streams that emerge from normal guns are so dense that they cannot interpenetrate without interfering. If we try to build a computer using streams of this density we couldn't allow any two wires of this kind to cross each other, so we'd better find some way to reduce the pulse rate.

In Fig. 28 the guns G_1 and G_2 produce normal glider streams in parallel but opposite directions. But there is a glider g which will travel West until at A it is kicked East by a glider from the G_1 stream. The timing and phasing are such that at B it will be kicked back towards A again, so that it repeatedly "loops the loop", removing one glider from each of the two streams per cycle. After this every Nth glider is missing from each of these streams. We don't want the G_1 stream, so we feed it into an eater, but we feed the G_2 stream into a vanishing reaction with a stream from a third gun G_3 . Every glider from G_2 now dies, but every Nth one from G_3 escapes through a hole in the G_2 stream! So the whole pattern acts as a **thin**



Figure 28. Thinning a Glider Stream.

gun, producing just one Nth as many gliders as the normal gun. To get the phasing right, N must be divisible by 4, but it can be arbitrarily large and so we can make an arbitrarily thin stream. Now two such streams can cross without interacting as in the right hand part of Fig. 28, provided things are properly timed. So from now on we can use the word **gun** to mean an arbitrarily thin gun. Perhaps a thinning factor of 1000 will make all our constructions work.

Building Blocks for Our Computer

In Fig. 29 we see how to build logical gates using only vanishing reactions (we've already seen the NOT gate in more detail in Fig. 20). But there's a problem! The output streams from the AND and OR gates are *parallel* to the input, but the output stream from the NOT gate is at *right angles* to the input. We need a way to turn streams round corners without complementing them, or of complementing them without turning them round corners. Fortunately the solution to our *next* problem automatically solves *this* one.

The new problem is to provide several copies of the information from a given glider stream, and we found it a hard problem to solve. To get some clues, let's see what happens when we use one glider to kick back a glider from a gun stream.



Figure 29. (a) An AND Gate. (b) An OR Gate. (c) A NOT Gate.

We suppose that the gun stream, the **full stream**, produces a glider every 120 generations (a quarter of the original gun density; N = 4 in the previous section). Then it turns out that when we kick back the first glider, the effect is to remove just three gliders from the stream! This happens as follows:

- (i) The first glider is kicked back (Fig. 27) along the full stream.
- (ii) The second glider crashes into the first, forming a block (somewhat as in Fig. 19(b)).
- (iii) The third glider annihilates the block (Fig. 25(a)).
- (iv) All subsequent gliders from the full stream escape unharmed.

We can use this curious behavior as follows. Suppose that our information-carrying stream operates at one tenth, say, of the density of the full stream, so that the last 9 of every 10 places on it will be empty, while the first place might or might not be full. If we use 0 for a hole and block the places in tens, our stream looks like

```
\dots 00000000D \ 00000000C \ 00000000B \ 00000000A \rightarrow
```

We first feed it into an OR gate with a stream of type

 $\dots 0000000g0 \ 0000000g0 \ 0000000g0 \ 0000000g0 \to$

the g's denoting gliders that are definitely present. The result is a stream

 $\dots 0000000gD \ 0000000gC \ 0000000gB \ 0000000gA \rightarrow$

in which every information-carrying place is definitely followed by a glider g. This stream is used to kick back a full stream whose gliders are numbered:



Figure 30. Copying a Glider Stream.

Auxiliary Storage

If glider A is *present*, it will obliterate gliders 1, 2 and 3 of the full stream and the following glider g can escape in the confusion. But if A is *absent*, then full stream glider 1 escapes and gliders 2, 3, 4 are removed instead by the following glider g. So the stream which emerges is definitely empty except for the second of every ten places and these places carry a copy of the input stream. The original full stream now manages to carry the information *twice*, in the first and fourth digits of each block, the first digit carrying the *complemented* version (which has not been turned through a right angle). By feeding this stream into vanishing reactions with suitably thin streams we can recover the original stream either complemented or not, and freed from undesirable accompanying gliders! Figure 30 shows these techniques in action.

From here on it's just an engineering problem to construct an arbitrarily large finite (and very slow!) computer. Our engineer has been given the tools—let him finish the job! We know that such computers can be programmed to do many things. The most important ones that we will want it to do involve emitting sequences of gliders at precisely controlled positions and times.

Auxiliary Storage

Of course the engineer will probably have designed an internal memory for our computer using circulating delay lines of glider streams. Unfortunately this won't be enough for the kind of problem we have in mind, and we'll have to find some way of adjoining an *external* memory, capable of holding arbitrarily large numbers. To build this memory, we'll need an additional static piece (the block).

Had Fermat's problem been still unsolved, we might ask the computer to compute

$$a^n + b^n$$
 and c^n

for all quadruples (a, b, c, n) in turn and stop when it finds a quadruple for which

$$a^n + b^n = c^n.$$

We don't know how big a, b, c and n might get, and they'll almost certainly get too large even to be written in the internal memory.

So we're going to adjoin some auxiliary storage registers, each of which will store an arbitrarily large number. Figure 31 shows the general plan. Each register contains a block, whose distance from the computer (on a certain scale) indicates the number it contains. In the figure, register A contains 3, B contains 7, C contains 0 and D contains 2. When the contents of a register is 0, the block is just inside the computer. All we have to do is to provide a way for the computer to

> *increase* the contents of a register by 1, *decrease* the contents of a register by 1, and *test* whether the contents are 0.

Fortunately each of these can be accomplished by a suitable fleet of gliders. One such fleet is off to increase register B by one! And another glider is about to discover that register C contains 0.



Figure 31. Auxiliary Storage.

How We Move Blocks

To find these fleets we studied the six possible glider-block crashes. One of them does indeed bring the block in a bit, but unfortunately by a knight's move. However the block can be brought back onto the proper diagonal by repeating the process with a reflected glider on a parallel course. The combined effect of this pair of gliders is to pull the block back three diagonal places (Fig. 32).

Unfortunately there is no single glider-block crash which moves the block further away, but there is a crash which produces the arrangement of 4 behives we call a honey farm , and two of these four are slightly further away, and so we can send in second, third and fourth gliders to annihilate three of the beehives, and then a fifth glider which converts the remaining beehive back into a block. The total effect again pushes the block off the proper diagonal, but a second team of five gliders will restore this, resulting in a block just one diagonal place further out! (Fig. 33).



Figure 32. Two Gliders Pull a Block Back Three Diagonal Places.



Figure 33. Ten Gliders Move a Block Just One Diagonal Place.

We therefore choose a diagonal distance of 3 to represent a change of 1 in a register and can decrease the contents of a register using a pair of gliders, or increase it using 3 flotillas of 10.

Apart from the difficulty discussed in the next section we have now finished the work, for Minsky has shown that a finite computer, equipped with memory registers like the ones in Fig. 31, can be programmed to attack arbitrarily complicated mathematical problems.

A Little Difficulty

But now comes the problem. Every glider in our finite computer has at some time been produced by a glider gun, so how could we arrange to send those gliders along closely parallel, but distinct paths? Surely one gun would have to fire right through another (Fig. 34)? Our technique of **side-tracking** uses three computer controlled guns G_1 , G_2 , G_3 as in Fig. 35. These are programmed to emit gliders exactly when we want them to.



Figure 34. How Can One Gun Fire Through Another?



Figure 35. Side-tracking.

Firstly, G_1 emits a glider g travelling upwards,

Secondly, G_2 emits a glider at just the right time to kick g back downwards,

Thirdly, G_3 kicks g back up again,

and so on, alternately, until at a suitable time G_2 fails to fire and g is released. By controlling the number of times G_2 and G_3 fire, the same guns can be used to send a succession of gliders along distinct parallel paths.

Mission Completed--Will Self-Destruct

Side-tracking can be used for a much more spectacular juggling act! We can actually program our computer to throw a glider into the air and bring it back down again. In Fig. 36, G_1 , G_2 , G_3 behave as before and can be programmed to arrange that a glider g ends up travelling Eastwards arbitrarily far above the ground. But G_4 has been arranged to emit a glider which will be kicked back down by g. We could even arrange to kick it back up again, then down again, then up again,... as suggested by the dotted lines in Fig. 36.



Figure 36. Double Side-tracking.

Using such techniques we can design a program for our computer which will send large numbers of gliders far out into space and then turn them round so that they head back towards the computer along precisely defined tracks (Fig. 37).

Now comes the clever part. Figures 38(a), 38(b) and 25(a) show that the eaters, the guns' moving parts, and blocks, can all be destroyed by aiming suitably positioned gliders from behind their backs. If the computer is cleverly designed we can even destroy it completely by an appropriate configuration of gliders!

Here's the idea. We design the computer so that every glider emitted by a gun or circulating in a loop would, if not deflected by meeting other gliders, be eventually consumed by an appropriately placed eater. Then we design our attacking force of gliders to shoot the computer down, guns first. After each gun is destroyed we wait until any gliders it has already emitted have percolated through the system and either been destroyed by other gliders, or swallowed by eaters, before attacking the next gun. When all the guns are destroyed we shoot down the eaters and blocks.



Figure 37. Self-attack from All Directions!



Figure 38. (a) The Eater Eaten! (b) The Gun Gunned Down!

The whole process requires some care. Each gun G_i must have a matching eater E_i , and G_i and E_i lie in a strip of the plane which contains no other static parts of the computer (Fig. 39). The gliders g_1, g_2, g_3, \ldots with which we shoot down a given gun can be arbitrarily widely spaced in time provided they come in along the right tracks. Moreover we can arrange to shoot down the successive guns, eaters and blocks after increasingly long intervals of time.

 \diamond



Figure 39. Arranging Destroyable Guns.

However, it *can* be done! We intend to use it like this. Program the computer to look for a solution of an arbitrarily hard problem, such as Fermat's. If it never finds a solution it will just go on forever. However, if it *does* find a solution we instruct it to throw into the air a precisely arranged army of gliders, then reduce all its storage numbers to zero (this brings all the blocks inside the computer), switch off, and await its fate. The attacking glider army, of course, is exactly what's needed to obliterate the computer, leaving no trace. It's important to realize that a *fixed* computer can be programmed to produce many different patterns of gliders and in particular the one required to kill itself. The information about this glider pattern can be held by the numbers in the memory of the computer and not in the computer's design.

Since mathematical logicians have proved that there's no technique which guarantees to tell when arbitrary arithmetical problems have solutions, there's no technique that's guaranteed to tell even when a Life configuration will fade away completely. The kind of computer we have simulated is technically known as a *universal machine* because it can be programmed to perform any desired calculation. We can summarize our result in this answer to our chapter heading:

LIFE IS UNIVERSAL!

Extras

Life is Still Being Lived!

The game of life is very much alive on the Internet, where a search will soon reveal many web pages devoted to it, which have found many new things. We shall only mention Dean Hickerson's page: http://www.math.ucdavis.edu/dean/life.html and Mark Niemiec's extended counts of still-lifes (i.e., stable objects) and flip-flops (i.e., period 2 oscillators):

size	still-lifes	pseudo still-lifes	flip-flops	pseudo flip-flops
3	0	0	1	0
4	2	0	0	0
5	1	0	0	0
6	5	0	3	0
7	4	0	0	0
8	9	1	1	0
9	10	1	1	0
10	25	7	1	1
11	46	16	1	2
12	121	55	6	14
13	240	110	3	17
14	619	279	20	46
15	1353	620	29	78
16	3286	1645	98	225
17	7773	4067	199	484
18	19044	10843		
19	45759	27250		
20	112243	70637		
21	273188	179011		
22	672172	462086		
23	1646147	1184882		

Life Computers Can Reproduce!

Eaters and guns can be made by crashing suitable fleets of gliders, so it's possible to build a computer simply by crashing some enormously large initial pattern of gliders. Moreover, we can design a computer whose sole aim in Life is to throw just such a pattern of gliders into the air. In this way one computer can give birth to another, which can, if we like, be an exact copy of the first. Alternatively, we could arrange that the first computer eliminates itself after giving birth; then we would regard the second as a reincarnation of the first. There are Life patterns which behave like self-replicating animals.

There are Life patterns which move steadily in any desired rational direction, recovering their initial form exactly after some fixed number of generations.

Genetic Engineering

We've now shown that among finite Life patterns there is a very small proportion behaving like self-replicating animals. Moreover, it is presumably possible to design such patterns which will survive inside the typical Life environment (a sort of primordial broth made of blocks, blinkers, gliders, ...). It might for instance do this by shooting out masses of gliders to detect nearby objects and then take appropriate action to eliminate them. So one of these "animals" could be more or less adjusted to its environment than another. If both were self-replicating and shared a common territory, presumably more copies of the better adapted one would survive and replicate.

Whither Life?

From here on is a familiar story. Inside any sufficiently large random broth, we expect, *just by chance*, that there will be some of these self-replicating creatures! Any particularly well adapted ones will gradually come to populate their territory. Sometimes one of the creatures will be accidentally modified by some unusual object which it was not programmed to avoid. Most of these modifications, or **mutations**, are likely to be harmful and will adversely affect the animal's chances of survival, but very occasionally, there will be some *beneficial* mutations. In these cases the modified animals will slowly come to predominate in their territory, and so on. There seems to be no limit to this process of evolution.

It's probable, given a large enough Life space, initially in a random state, that after a long time, intelligent self-reproducing animals will emerge and populate some parts of the space.

This is more than mere speculation, since the earlier parts are based on precisely proved theorems. Of course, "sufficiently large" means very large indeed, and we can't prove that "living" animals of any kind are likely to emerge in any Life space we can construct in practice. It's remarkable how such a simple system of genetic rules can lead to such far-reaching results. It may be argued that the small configurations so far looked at correspond roughly to the molecular level in the real world. If a two-state cellular automaton can produce such varied and esoteric phenomena from these simple laws, how much more so in our own universe?

Analogies with real life processes are impossible to resist. If a primordial broth of aminoacids is large enough, and there is sufficient time, self-replicating moving automata may result from transition rules built into the structure of matter and the laws of nature. There is even the very remote possibility that space-time itself is granular, composed of discrete units, and that the universe, as Edward Fredkin of M.I.T. and others have suggested, is a cellular automaton run by an enormous computer. If so, what we call motion may be only simulated motion. A moving particle in the ultimate microlevel may be essentially the same as one of our gliders, appearing to move on the macrolevel, whereas actually there is only an alteration of states of basic space-time cells in obedience to transition rules that have yet to be discovered.

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Glossary

 $A = ace = \{0 | tiny\}, 357$ $\bar{\mathbf{A}} = -\mathbf{ace} = \{\min \mathbf{y} \mid 0\}, 359$ $A = \{on | A | | 0\}, 359$ $A + = \{0 || \bar{A} | off \}, 359$ ℵ₀, aleph-zero, 329 Δ , also, slow join, 286 \wedge , and, join, 278], "ceiling," least integer not less than, 485 $(x)_n, \{x \mid -y\}_n$, Childish Hackenbush values, 238a < b >, class a and variety b, 363 $\bar{1} \clubsuit = \{ \clubsuit \mid 0 \}, \text{ clubs}, 359$ $\mathbf{A} = 0 \mathbf{A} = \{1 \mathbf{A} \mid 0\}, \text{ clubs}, 359$ $1 \clubsuit = \{ deuce | 0 \}, clubs, 359 \}$ • • • () (in the content of the cont γ° , degree of loopiness, 361 $2 = \{0 | ace\} = ace + ace = deuce, 357$ $\bar{1}\diamondsuit = \{\bar{\mathbf{J}}|0\}, \text{ diamonds}, 359$ $\diamond = 0 \diamond = \{ \mathbf{ace} \mid \bar{1} \diamond \}, \text{ diamonds}, 359 \}$ $1\diamondsuit = \{0|\diamondsuit\}, \text{ diamonds}, 359$ $\downarrow = \{\downarrow * | 0\} = \downarrow + \downarrow$, double-down, 68, 69, 71 $\uparrow = \{0 \mid \uparrow *\} = \uparrow + \uparrow$, double-up, 68, 69, 71 $\uparrow * = \{0 \mid \uparrow\} = \uparrow + \uparrow + *, \text{ double-up star, 71}$ $\downarrow = \{*|0\}, \text{ down, } 64$ $\downarrow_2 = \{\uparrow * | 0\},$ down-second, 235, 236 $\downarrow_3 = \{\uparrow + \uparrow^2 + * | 0\}, \text{ down-third, } 235, 236$ $\downarrow_{abc} \ldots, \downarrow_{abc} \ldots, 261$ V, downsum, 336, 357

 $dud = \{ dud | dud \}, deathless universal draw,$ 337 e = 2.7182818284..., base of natural logarithms, 610 ϵ , epsilon, small positive number, 328 $\varepsilon =$ **over**, 674, 675, 688, 689, 691–702 \doteq , equally uppity, 242, 247, 249 | |, "floor," greatest integer not greater than, 51, 311, 485 ||0, fuzzy, 29, 32, 37 G, general game, 28–31 $G \parallel 0, G$ fuzzy 2nd player wins, 29 G < 0, G negative, R wins, 29 G > 0, G positive, L wins, 29 G = 0, G zero, 1st player wins, 29 G + H, sum of games, 31 G^L , (set of) L option(s), 31 G^R , (set of) R option(s), 31 $\mathcal{G}(n)$ nim-value, 82 $G. \uparrow = \{G^{\mathrm{L}}, \uparrow + \uparrow \ast \mid G^{\mathrm{R}}, \uparrow + \Downarrow \ast\}, 247,$ 255, 258 >, greater then, 32 \geq greater than or equal, 32 \triangleright , greater than or incomparable, 32, 35 \geq , at least as uppity, 242, 246 $\frac{1}{2} = \{0|1\}, \text{ half, } 7, 20$ $\overline{1}\heartsuit = {\heartsuit | 0}, \text{ hearts, } 359$ $\heartsuit = 0 \heartsuit = \{1 \heartsuit | \mathbf{A}\}, \text{ hearts, } 359$ $1\heartsuit = \{0 | \mathbf{joker}\}, \text{ hearts}, 359$

Glossary

 $hi = \{on || 0 | off\}, 355$ $hot = \{on | off\}, 355, 710$ ||, incomparable, 35 $\infty = \mathbb{Z} || \mathbb{Z} | \mathbb{Z}, 329, 334, 391$ $\pm \infty = \infty | -\infty = \mathbb{Z} | \mathbb{Z} = \int^{\mathbb{Z}} *, 334$ $\infty \pm \infty = \infty | 0 = \mathbb{Z} | 0, 334$ $\infty + \infty = 2 \cdot \infty = \mathbb{Z} \mid \mid \mathbb{Z} \mid 0$, double infinity, 334 $\infty_0 \ \bar{\infty}_2$, unrestricted tallies, 314 $\infty_{abc...}$, 387–395 $\infty_{\beta\gamma\delta...}$, 333 ∫, integral, 167–180 $\mathbf{J} = \{0|\bar{\mathbf{A}}+\} = \mathbf{ace} \wedge (-\mathbf{ace}) = \mathbf{joker}, 358$ $\bar{\mathbf{J}} = \{\bar{\mathbf{A}} - |0\} = \mathbf{ace} \, \forall (-\mathbf{ace}) = -\mathbf{joker}, 359$ L, Left, 2 LnL, LnR, RnR, positions in Seating games,44, 45, 132, 133, 260 <, less than, 32 < less than or equal, 32 \triangleleft , less than or incomparable, 32, 35 $lo = {on|0|| off}, 355$ D, loony, 397, 407 $\gamma, \gamma^*, \gamma^+, \gamma^-$, loopy games, 335 s&t loopy games, 336 $-1 = \{ |0\}, \text{ minus one, } 19$ $-_{on} = \{ on | 0 || 0 \}, miny, 353$ $-_{\frac{1}{4}} = \{\frac{1}{4} | 0 | | 0\},$ miny-a-quarter, 134 $-_{x} = \{x \mid 0 \mid \mid 0\}, \text{ miny-x}, 126$ * nim-product, 475–477 + nim-sum, 58 off = { |off|, 336–340, 357, 674, 710 $\omega = \{0, 1, 2, \ldots \mid \}, 329 - 333$ $\omega + 1 = \{\omega \mid \}$, omega plus one, 329–333 $\omega \times 2 = \{\omega, \omega + 1, \omega + 2, \dots \mid \} = \omega + \omega, 329 - 333$ $\omega^2 = \{\omega, \omega \times 2, \omega \times 3, \dots \mid \} = \omega \times \omega, 329 - 333$ $on = \{on | \}, 336 - 341$ $1 = \{0|\}, \text{ one, } 7, 19$ $1/\mathbf{on} = \mathbf{over}, 341, 673, 674, 686$ 1 over, 674, 677, 681, 686 **ono** = {**on**|0}, 355–356 $oof = \{0 | off \}, 355-357$ **over** = $\{0|$ **over** $\} = \frac{1}{\text{on}}$, 341, 674, 688, 696–695 $\pi = 3.141592653..., pi, 328-329, 610$ $\pm 1 = \{1 | -1\},$ plus-or-minus-one, 120–122 (± 1) , $\uparrow = \{\uparrow \uparrow | \Downarrow \}, 247$ $\Psi = \{ \psi * | 0 \} = 4$, \downarrow , quadruple-down, 347 $1 = 4. \uparrow$, quadruple-up, 71, 347, 693, 706

 $\frac{1}{4} = \{0|\frac{1}{2}\},$ quarter, 6, 20 $\frac{\hat{1}}{4} = \frac{1}{4} \cdot \uparrow = \{ \uparrow * | 1 \frac{1}{2} \cdot \downarrow + * \},$ quarter-up, 236 $\frac{\hat{1}}{4} * = \frac{1}{4} \cdot \uparrow + * = \{ \uparrow | 1\frac{1}{2} \cdot \downarrow \}, \text{ quarter-up-star}, \}$ 236R, Right, 2 $\frac{*}{2} = \{*, \uparrow \mid \downarrow *, 0\}, \text{ semi-star}, 370$ $\frac{1}{2} = \frac{1}{2} \cdot \uparrow = \{ \Uparrow * | \downarrow * \}, \text{ semi-up, } 236, 247$ $\frac{1}{2}*=\frac{1}{2}$, $\uparrow+*=\{\uparrow\downarrow\downarrow\}$, semi-up-star, 236 $\frac{3}{2} = 1\frac{1}{2}$, $\uparrow = \{* | *\}$, sesqui-up, 236 sign(), 348-350 |, ||, |||, ... slash, slashes, ... separate L and R options, 6-7, 128-129, 366 $2 = \{ \bar{\mathbf{A}} | 0 \}$, spades, 359 $\bar{1} \spadesuit = \{0 | \bar{2} \spadesuit\}, \text{ spades}, 359$ $\mathbf{A} = 0\mathbf{A} = \{0|\bar{1}\mathbf{A}\}, \text{ spades}, 359$ $1 \spadesuit = 0 | \spadesuit$, spades, 359 🕅 , far star, remote star, 230–232, 244–251 $* = \{0|0\}, \text{ star}, 38$ $*2 = \{0, *|0, *\},$ star-two, 41 $n = \{0, *, \dots, *(n-1) | 0, *, \dots, *(n-1)\},$ starn, 41 $*\alpha$, star-alpha, 333 $\hat{*} = *. \uparrow = \{ \uparrow * | \downarrow * \}, \text{ starfold-up}, 236, 248$ $*\bar{n}$, all nimbers except *n, 397 $*n \rightarrow .$, all nimbers from *n onwards, 397 $\odot = 0 * \rightarrow$, sunny, 397–401, 404 $\uparrow_{abc...} = -\downarrow^{abc...}$, superstars, 261 2130, tallies, 300-326 $\frac{3}{4} = \{\frac{1}{2}|1\}, \text{ three-quarters, } 17$ $+_{on} = \{0|oof\} = \{0||0|off\} = tiny, 353, 357$ $+\frac{1}{4} = \{0 \mid |0| - \frac{1}{4}\}, \text{ tiny-a-quarter, } 126$ $+_{2} = \{0 \mid | \mid 0 | -2\}, \text{ tiny-two, } 126$ $+_{x} = \{0 \mid | \mid 0 \mid -x\}, \text{ tiny-}x, 126$ $tis = {tisn | } = 1\&0, 342, 374$ $tisn = \{ |tis\} = 0\& -1, 342, 374 \}$ $(l, r)(l, r)_c$, Toads-and-Frogs positions, 127, 136, 368. 375. 376 $\Psi = \{ \downarrow * | 0 \} = 3. \downarrow, \text{ treble-down, } 71, 347$ $\uparrow \uparrow = \{0 | \uparrow *\} = 3. \uparrow, \text{ treble-up}, 71, 347, 693$ $3 \triangleq \{0 | \mathbf{deuce}\} = \mathbf{trey} = \mathbf{ace} + \mathbf{deuce}, 357$ \triangle , triangular number, 253 U, ugglies, ugly product, 483–487 $2 = \{1|\}, \text{two}, 7, 19$ under = {under |0} = -over, 341 ∨,**union**, or, 300

964
$\uparrow = \{0|*\}, \text{ up, } 64, 71, 261$ $\uparrow^{\alpha}, \text{ up-alpha, } 341$ $\uparrow^{2} = \{0|\downarrow *\}, \text{ up-second, } 235, 236$ $\uparrow^{*} = \{0, *|0\}, \text{ up-star, } 65, 229$ $\uparrow^{3} = \{0|\downarrow + \downarrow_{2} + *\}, \text{ up-third, } 235, 236$ $\uparrow_{abc...}, 261$ **upon** = {**upon**|*}, 341, 375 upon* = {0, upon * |0}, 341
∧, upsum, 336, 357
∨, ur, urgent number, 312
[a|b|c|...]_k, Welter function, 506–514
ℤ = {..., -2, -1, 0, 1, 2, ...}, the set of integers, 334
0 = { | } = 0* = 0. ↑, zero, 7, 41

965

 \diamond

Index to Volumes 1–4

abacus positions, 512–515 Abacus strategy, 513, 514 abnormal move, 325 absorbancy, downsum, 360 Abt, Clark C., 959 accounts-payable, 127, 826, 827, 838 ace, 357 acrostic games, 482, 485, 487 Mock Turtle Fives,, 487 product, 482, 483, 486 Turnips, 483 Twins, 487, 473, 474 action, in hottest game, 167 active postion, 149, 150 acute triangles, 254 Adams, Clifford W., 892, 919 Adams, E. W., 117 Adam's Magic Hexagon, 892, 919 Adders = \cdot **73**, 429, 444 Adders-and-Ladders, 386 addition misère, 419 of Checkers, Chess, Domineering, Go, 757 of games, 30-33 of loony games, 399 of loopy games, 390 of switches, 120

nim-, 58. 59, 73, 74, 90, 109, 110, 116, 199, 246, 390, 418, 473, 475, 476, 497, 498, 504, 507, 508, 510, 556, 558, 605, 606, 752, 860, 886 ordinal, 219, 220 to switches, 122 two kinds, 191 Additional Subtraction Games, 395 additives of no atomic weight, 225 adjacency matrix, 222 age moon's, 907 Air on a \mathcal{G} -string, 96 ajar, 408-410 Akin, E., 766 Albert, M., 690, 757, 763 Algorithm Secondoff, 535, 537 Zeckendorf, 535 Alice in Wonderland, 1–3 56, 57, 229, 461–463– 473all small games, 229, 230, 757 All Square, 405 All the King's Horses, 277–283, 286–289, 292– 293Allemang, Dean Thomas, 454 Allen, J. D., 764 Allis, L. V., 740, 764, 766, 767 almonds, 379, 393 also-ran, 286

alternating group, 867 alternating moves, 46, 47 Althöfer, I., 86, 117 altitude, 31, 693, 687-689, 692, 694, 702, 703, 705altitude decrement, 700, 701, 703 exceptional, 704 amazing jungle, 208, 221 Amazons, 15, 16, 756, 757, 761, 762 ambient temperature, 164, 168, 188 ambitious distraction, 483 ambivalent Nim-heaps, 426 Amendment, tactical, 697 American colonies, 905 anatomy of Toads-and-Frogs, 63 AND gate, 941, 948 Anderson, William N., 767 Andersson, Göran, 119 Andreas, J.M., 891 Andrews, W.S., 923 angel, 643, 665 anger, fit of, 317 Animal Farm, 146, 812 animals dead, 135 Grundy's wild 431 tame, 425-430 tracking, 202-210, 221-222 annihilation games, 224 Anshelvich, Vadim V., 744, 767 Anthony, Piers, 607 Antipathetic Nim. 493–496 499 Antonim, 493–496, 499 Applegate, David, 602, 607 Arbib, Michael A., 959 Archangel, 9, 30, 1916 Argument, Poker-Nim, 55, 56 Argument, Snakes-in-the-Grass, 191 Argument, Tweedledum and Tweedledee, 2, 3, 35, 72, 349-350 Arithmetico-Geometric Puzzle, 847, 848, 913-915arithmetic periodicity, 99, 113-117, 144, 187, 752army, 95, 821-823, 835 arravs Nimstring, 581

arrows, 551, 741, 835, 849, 854 Ars Amatoria, 736, 769 Artful Arrow, 850, 854 Arthur, 541, 542, 545, 551, 554, 569, 574, 575 assets, 824, 825 asymmetrical heating, 173 atom. 229 atomic weight = uppitiness, 200-202, 204, 206, 208-209, 211, 220-222, 225, 230-232, 234, 236-240, 242-244, 246, 248, 251, 253, 256, 258, 259, 261 calculus, 231, 232, 234, 237, 239, 242, 247, 249, 256, 259 of lollipops, 238 of nimbers and of up, 232 atomic weight, 690, 757 eccentric, 231, 232, 237, 251 fractional, 234, 236 high, 706 hot, 706 properties of, 236 rules, 242, atoms Lucasta, 588-594 superheavy, 333 Austin, A.K., 923 Austin, Richard Bruce, xiv, xviii, 52, 96, 109, 117, 538average versus value, 10 averages, playing the, 167, 173 Ayo, 761 Babbage, Charles, 767 baby-sitting, 413 Bach, Clive 369 back-handed compliment, 405 backbone, 700, 702, 704 sequence, 700, 701, 704 values, 700, 701 backfield, 681, 695, 700, 710 Backgammon, 14 Backsliding Toads-and-Frogs, 367, 375–376 backwards playing, 817 thinking, 384 bad child, 430 bad move, 16, 547, 818 Baduk (see Go)

\diamond

Baiocchi, Claudio, 533 Baked Alaska, 321, 323 baker, bakery, 936 balance, 381, 824, 896 sheet, 823 ball, 96, 752 Ball, W.W. Rouse, 117, 538, 895, 923 Ball and Chain, 850 Balsdon, J.P.V.D., 767 Banks, E. Roger, 939, 959 bargain, 159 barge, 929 Barmy Braid, 853 Baseball, 15 Basketball, 15 Battle, xvii(2)battle, 3 hot, 145, 299-307, 310, 322, 326 Nim-like, 230 Battleships, 15 beacon, 930 bead = bivalent node, 562Beal, D. F., 764, 765 Beasley, John D., 18, 640, 805–807, 837–840 Beasley's Exit Theorems, 829, 830, 837, 838 bed, redwood, 211, 213-217, 222 bee, queen, 937 beetle, 600 behavior for Princes, 537 beehive, 929, 934, 935, 943, 953 Beeler, Michael, 939 Belgium, 894 Bell, A. G., 767 Bell, Robert Charles, 666, 668, 710, 737, 767 Belladonna, 528 Benson, David B., 764 Benson, David J., 871, 875, 876 Benson, E. C., 117 Bentz, Hans J., 766 Berge, Claude, 78 Berghold, Ernest, 805, 806, 837, 839, 840 Berlekamp, Elwyn, 16, 18, 77, 78, 123, 144, 187, 188, 224, 225, 318, 538, 578, 584, 607, 693, 710, 756, 757, 762-764 Berlekamp's Rule, 77, 78 Bertha, 541–545, 547, 551, 554, 555, 569, 574 Bessy, Frenicle de, 886

Bicknell-Johnson, Marjorie, 79 Biclock, 937 big firms, 423 big game, 65, 75, 202, 206, 222, 417 Biloaf, 937 binary (base 2), 73, 99, 481,940 tree. 23 bipartite graphs, 217, 224, 583 birth, 927 control. 929 birthday, 318, 417, 771, 903, 962 bit = binary digit, 99, 940Bjørne, Anders, 766 Bjornsson, Y., 764 Black, Farmer, 145, 146 Black, Larry, 746 Black Path Game, 746, 747 black heap, 532 black stones = blocking stones, 642blanker, 929 Blass, Andreas, 643, 748 Blass, U., 78 blatantly loopy, 391 blatantly winning ways, 410 Blinker, 928, 929, 935, 937, 943 block, 941, 947, 948, 951, 952, 955-957 puzzles, 843-848, 864, 877-885, 892-898, 910-915, 919-922 Blockbusting, 187 blocking stones, 642-664 Blocks-in-a-Box, 844, 910 blossom, 198, 199, 240 bLue edge, 2, 198, 230, 237 Blue Flower Ploy, 199, 201, 240, 242 Blue Jungle Ploy, 201, 210 bLue tinted nodes, 48, 204 Blue-Red Hackenbush, 1-6, 15-17, 20, 22, 23, 27-33, 35-38, 43, 77, 78, 179, 200, 211-217, 224 Blue-Red-Green Hackenbush, see Hackenbush Hotchpotch Board Continental, 817, 835 English, 803 board games, 641-842 sizes, 642, 649, 654, 657, 663, 664

boards, F&G, 686 Boardman, Mike, 818 boardwidth large, 710 boat, 929, 935, 936 Bodlaender, H. L., 117, 225 body, 502, 503, 609, 741, 807, 900 bogus Nim-heap, 57-58, bomb, time, 937 Bond, James = $\cdot 007$, 94 Bono, Edward de, 384, 408, 412 bonus move, 405 booby prize, 533 bottle, 218, 225 boundary left, 154-156, 158, 164, 169 right, 154-156, 164, 169, 170 booby prize, 533 Bounded Nim, 517, 518 Bouton, Charles L., 42, 52, 78, 452 Bouwkamp, C.J., 923 box, 488, 501, 541-545, 547-551, 554-556, 558, 561, 564, 566, 569, 571, 572, 575, 577-579, 592, 844, 845, 896 Boxing, 488 Boxing Day, 903, 905 Boy Leaves Girl, 856 boys by billions, 547 braiding paper, 853 branch, 22, 34, 191, 555, 564, 588-590, 592, 594, 606 Brandt, J., 766 brazils, 393 breakthrough, 188 Breuker, D. M., 144 Bridge, 17 bridge, 189, 194-196, 685-687, 693-695, 866, 916 Bridgit, 744-746 Britain, 905 Broline, D., 766 Brousseau, Bro. Alfred, 923 Brouwer, Andreas E., 740, 768 Brown, B.H., 923 Brown, T.A., 923 Browne, Cameron, 744, 767 Browne, Patrick, xvi

Brualdi, Richard A., 767 Bruckner, Gottfried, 767 Bruijn N.G. de, 492, 840 Brussels Sprouts, 603, 607 Buckingham, David, 936 Bugs, 597 Bulgarian Solitaire, 761 bulls 145, 146 bullseve, 681-684 Bumble-bee Problems, 827 Bundy, A., 765 Buro, Michael, 762, 765 Bumby, Margaret, 875 Bushenhack, 604, 606, 607 Busschop, Chevalier Paul, 840 Bűvös Kocka, 868–876, 917–918 Bynum's Game = Eatcake, 136, 234, 289Bynum, James, 136, 234, 289 bypassing reversible options, 60, 62–64, 66, 70, 71.75.77 Byrom, John, 53 Byron, Henry James, 81 CABS, 388-393 cabs, 364-366 Cabbages, 597 Caesar, Julius, 585 cage, 709 Caines, Ian 117 cake, 26, 51, 196, 220, 221, 284-288, 289-292, 294-297, 299-303, 308-310, 321-323 Calculus, Atomic Weight, 231, 232, 234, 237, 239, 242, 247, 249, 256, 259 calendar Gregorian, 905–909 Julian, 904-909 Calistrate, Dan, 188 candelabra, 936 Cannings, C., 766 canonical form for numbers, 22 Canterbury puzzle, 767 capturable coin, 555, 556, 558 capture, custodian, 666 captures, 665, 710 Cardano, Girolamo, 923 card sharp, 16 cards, shuffling, 14 Cards, House of, 357-361

970

 \diamond

carousal, 369, 380, 382 Carpenter, 461, 462 Carpets, 478-480 Fitted, 479 carpet, Greatly-Valued, 480 Carroll, Lewis, 229 Carteblanche, Filet de, 525, 538, 923 cash flow, 126 cashews, 393 Cashing Cheques, 122–124, 145, 158, 159, 240 Catherine Wheel, 932 Cat, Cheshire, 936 catalyst, 807-809, 812-813 caterpillar, 502, 503, 597 ceiling, 51, 485 Celoni, James R. 226 central formation, 695 central Solitaire, 804-807, 809 central Soma piece, 847 centralizing switches, 123 centred King, 657, 658 Century Puzzle, 877-878, 881-885, 919 Century-and-a-Half, 877-878, 881, 885, 919 Cervantes's deathday, 906 cgsuite, 687, 690, 702, 703, 705 chain , ball and, 851 green, 40 long, 543, 545–547, 549, 550, 552–556, 561, 563, 564, 569, 574, 575, 577-579, 829 Lucasta, 588-590, 594 short, 547, 574 snapping, 564,565, 567 Snort, 147, 149-151, 153, 156-158, 161, 167, 168, 177, 180-183 chain of boxes, 542–547, 549–550, 552–555, 561 chair redwood, 211 swivel, 349 chalk-and-blackboard game, 1 chance moves, 14 Chandra, Ashok, K., 224, 227 Chang, Alice, 765 change of heart, xvii(2)change, phase, 167, 168 charge, electric, 253, 254 Charming Antipodean Beauty Spot, 391

Charming Charles, 525, 526 Charosh, M., 840 Chas., 641-643, 663, 664, 767, 923, 924 Chas. Scribner's Sons, 539 Checkers = Draughts, 18, 224, 226, 227, 757, 763Hawaiian = Konane, 690, 759, 760, 765cheque-market exchange, 158, 159 Cherries, 757, 763 Cheshire Cat, 936 Chess, 14, 18, 224, 226, 667, 755, 757, 763 complete analysis, 755 Dawson's 89–92, 101, 109 Japanese (See Shogi) problems, 658 Chessgo, 641, 643 chesspersons, 88, 641-643, 667, 710 child bad 430 good 427 Childish Hackenbush, 43, 52, 157 Childish Hackenbush Hotchpotch, 236–238 childish lollipops, 237-240 childish picture, 43 children's party, 132, 170 China 16 Chinese Nim = Wythoff's Game, 427 Chinese Rings, 856-861 chilling, 758, 764 chocolate bar, 632, 635 Chomp, 632, 635 chopped, 675 chopping, 564, 604 Christmas, 902, 908 Chvátal, Vasek, 737 Chvtie, M. P., 227 circle, 682-684, 691 double, 691 higher, 691 isolated, 684 lower, 691 single, 691 circled F&G position, 705, 681-685, 687, 694, 696 circled FOXTAC values, F&G, 691 circumscribed games, 470 class and variety, 362-363

972

class, outcome 28, 84 classes, Riess's Solitaire, 816 Claus = Lucas, 862clean and dirty fuses, 936 Clef, Double Treble, 848, 850 climbing bars, 211 clique, 613-615 635 Clique Technique, 613, 615 Clobber, 690, 757, 758, 761, 763 clock, 669 closed, 408-410 cloud, 31, 36, 121, 149, 150 clubs, 359, xxi(3) coalition, 15, 534 cockroach, 600 Cocoons, 597 Codd, E.F., 959 code digits, 92, 93, 101, 103-105, 107, 108, 113, 117code genetic, 605, 606 Gray, 858-861 of behavior, 537 coin sequence game, 635 coin-sliding, 539 coin-sliding games, 539, 764 Coinage Sylver, 15, 539, 609-631, 635-640 coins, 123-124, 464, 466, 467, 469, 470, 472, 473, 476-478, 481-483, 486-488, 491-494, 496, 499, 506-508, 516, 529, 550-552, 554-556, 609, 610, 617, 711, 863, 916, 923 colored, 532 Col, 37-39, 47-51, 67, 68, 75, 145, 224, 502 cold game, 299-300, 307, 316, 317 cold position, 301, 304-306, 322 cold war, 299-302, 304, 306, 307 cold work, 316 Coldcakes, 300 colon notation, 243, 244 Colon Principle 191, 193, 194, 220, 355 coloring, 38, 145, 147, 187 Commandment, Lukewarmth, 307, 310 Markworthy, 317 common coset, 110 common values, 110

comparing games, 35, 36, 122, 348 compendium, 109, 387 complementary position, 816 complementing effect, 405 complete, exptime, 763 complete graphs, 583 complete, NP, 224, 227 complete, PSPACE, 224, 227, 760 completing a box = complementing more, 541 complete in exponential time, 224 complete in Pspace, 224 complete information, 14 complicated value, 707, 708 complimenting move, 379, 405-407, 541, 552 component, 20, 22, 31, 32, 35, 278, 281, 286, 288-289, 299-302, 307, 312, 325-326, 396-398, 415cold, 299, 300 hot, 299, 300 loopy, 410 tepid, 327 compound conjunctive, 278, 284, 286 continued conjunctive, 286 disjunctive, 289 impartial, 396 selective, 299, 300 severed selective, 312 shortened selective, 312 subselective, 396, 533 compound game, 31, 299 compound thermograph, 164 computable function, 630 computing power, 163 confused, 31, 68-71, confusion interval, 121, 149–151, 158, 163 congruence modulo 16, 507 conjecture, 112, 537, 620 conjunctive compound, 278, 286 Connect-Four. 764 Connell, Ian G., 78 Continental board, 817, 835 continued conjunctive compound, 286 Contours, 587, 588 contract, 126 control, 544-546, 574, 575 convention, normal play, 14

Conway, Elena, 876 Conway, J.H, 18, 22, 52, 78, 117, 123, 144, 188, 225, 262, 298, 377, 446, 454, 489, 539, 598, 668, 752, 847 Cook, Stephen A., 224, 225 Coolcakes, 308, 309, 321 cooling, 151, 152, 154, 167, 179, 231, 690 formula, 151 Copper, Mark, 607 coprime, 404 Corderman, Charles L., 937 Corinthians I, 13.12, 69 cork-screw, left-handed, 379 corner, 437, 654-663, 822-823, 829-830 defence. 654-657 tactics, 654, 661 cornered King, 658–661 Corners, Turning, 473-475, 478 cosets, common, 110 costs, 161 counters, heaps of, 41 Couples, Seating, 44, 45, 132, 133 coupons stack of, 693 cousin, 101, 103-105, 107, 109, 114, 116 coverlet, 223 Cowley, Abraham, 119 cows, 145-147 Coxeter, Harold Scott Macmillan, 78, 117, 538, 539, 923 Cram = Impartial Domineering, 141, 143, 298, 502 - 506cricket, 15 criminal, minimal, 194, 214 critical position, 687, 690 critical rank, 710 critical temperature, 167, 168, 171 Cross, Donald C., 840 cross, 582, 585, 603, 607 Crosscram = Domineering, 119, 298crosses, tendrilled, 582 Csirmaz, László, 767 cube Hungarian, 768-776, 817-818 magic, 768-776, 817-818 Rubik, 768–776, 817–818 Culberson, J., 766

cul-de-sac, 832 Curtis, Robert Turner, 468, 469 custodian capture, 666 Cutcake, 25, 26, 31, 32 Hickerson's, 51 Cutcakes, 284-285, 292-294, 300 cutting, 284, 292 cycles, 192-194, 213-214, 316, 551, 930-932 Cyclotome, Alan Schoen's 898-899 D.A.R., 361 Dad's Puzzler, 877-879, 885 D'Alarcao, Hugo, 607 D.U.D.E.N.E.Y, 521, 523 Damß, J. E., 225 daggered position, 677, 710 danger, 670 dark positions, 677 darkening, 684, 694, 686 darkening move, 677, 678, 685 date, 395, 410 Davies, D.W., 767 Davis, Harry O., 839, 840 Davis, Morton, 766 Dawson's Chess = $\cdot 137$, 89–92, 101, 109 Dawson's Kayles = $\cdot 07$, 15, 90, 93, 95, 101, 109, 261, 438, 444 Dawson's vine, 566, 569, 576 Dawson, Thomas Rayner, 89, 117, 923 De Parville, 924 dead animals, 135 dead cell, 927 Deader Dodo Problems, 827 deadly dodge, 548 death, 927 deathday, 905 Death Leap principle, 127–130, 135 deceptive defence, 548 decomposing, 568, 569 decrement altitude, 700, 701, 703 decremented index, 705 deficient Soma piece, 847 deficit accounting, 827 Rule, 826, 831 Definite F&G boards, 674, 685–686

degree cooling by one, 179 of loopiness, 361 of **upon**, 375 Delannov, Henri-Auguste, 916, 923 deleting dominated opinions, 62, 63, 75, 77 delphinium 47, 199, 242, 695 lush, 675, 677, 681, 684 trimmed, 686 Welton's, 695 Delta, 697 Region, 697 Demaine, Erik D., 755, 763, 767 Demaine, Martin L., 755, 763, 767 Denim, 534 deriders of zero, 483 Descartes, Blanche, 96, 97, 117 desert, 821, 822 deuce, 357 devastating U-turns, 754 devil square-eating, 643 devil's label, 195 dexterity, 910, 913 diamonds, 359 dice octahedral, 891 paradoxical, 886 dictionary Col, 48-51 Cram, 502-506 Domineering, 120, 138-142 Nimstring, 559, 560, 565, 581–583 Snort, 147, 180-183 Difference Rule, 74, 404 digits binary, 99, 940 code, 92, 93, 98-106, 113, 117 Dim, 98, 442 with Tails, 404 disarray, 95 discount, 161 disentailing, 401 disguise, 95, 462, 586, 733 disincentive, 148 disjunctive compound, 278, 289 dissection, 128, 129, 134

disassociation, thermal, 168 distributive law, 475, 483 Dividing Rulers, 436-437, 469 do or die donation, 548 dodecahedra, quintominal, 923 Dodgem, 749-751 Dodgerydoo, 751 Dodie Parr = odd parity, 545-550, 558, 569, 570, 572, 575, 578-580 Dodo Problems, 827 dodos, 820 dog with leftward leanings, 4 dogmatism, 377 dogs, 711 Dollar Game, Silver, 491, 492, 535 dominated option, 62, 63, 75, 77, 126, 149 Domineering = Crosscram, 119-122, 125, 138-142, 153, 177, 178, 298, 366, 690, 757 Domineering, Impartial = Cram, 141, 143, 144, 298, 502-506 Dominoes, 119 Domoryad, A.P., 923 Donkey Puzzle, 877, 878, 880, 881, 885 Don't-Break-It-Up Theorem, 213, 214, 216 Doomsday Rule, 903-906, 923 Doors, 478, 481, 482 Dots-and-Boxes, 15, 95, 225, 539, 897 dots+doublecrosses=turns, 546, 571 Double Circle, 691 Double Duplicate Nim, 114 Double Hackenbush, 343 double infinity, 334 Double Kayles, 99 Double Treble Clef, 848, 850 double-cross, 541, 543, 544, 571 double-dealing, 542–544 double-down, \Downarrow , 68, 69, 71 double-six, 405 double-up, ↑, 68, 69, 71, 242 doubling, 620 of nim-values, 93, 98 down ↓, 68, 151 downon*, 708 down-second, 236 downsum, 336, 355, 357, 360 Draughts=Checkers, 18, 224, 226, 227, 757, 763

draw \neq tie, 14 drawn, 14, 335 dud = deathless universal draw, 337–338, 353Dress, Andreas, 117 Dudeney, Henry Ernest, 82, 117, 521, 539, 767, 805, 807, 825, 840, 923 Duffus, Dwight, 312 Duffy, Adam, 758 duke, 641, 642, 644-646, 757 Dukego, 644–646 Duncan, Anne, 732 Duplicate Kayles, 99, 444 Duplicate Nim, 114, 116 duplication of nim-values, 94, 98, 99, 114, 116, 444 - 445Durer, Albecht, 888 Duvdevani, N., 79 Dyen, Louis, 713 Eagle, Edwin, 768 early F&G values, 689 early sequence, 688, 689 early values, 700 early-stage, 689 earwig, 600 Easter, 905 Eatcake = Bynum's Game, 233, 234, 235, 292Eatcakes, 286, 289-291, 285-297, 308 Eater, 930, 943, 944 eating, 289-292, 643, 930, 943, 944 eccentric cases of atomic weights, 231, 232, 237, 239, 249, 251 economy, underlying, 151 Eddins, Susan, 607 Eden, Garden of, 938 edge attack, 646, 650 corner attack, 650–652 defence, 647-649 edges, 40, 43, 135, 345 bLue and Red, 2-6, 29, 30, 77, 198, 201, 230, 237, 329, 343 grEen, 29, 30, 33, 40, 41, 190-196, 198-202, 204-206, 210, 211, 218, 220, 221, 225, 237, 251, 330 pink and pale, 343–344 Edmonds, Jack, 767 effective computability, 640

Eggleton, Roger Benjamin, 522 Eight-in-a-row, 740 El Acertijo, 756 electric charge, 251, 253, 254 Eliot, Thomas Stearns, 461, 534 Elizabeth, Queen, 905 Elkies, Norm, xiii, xvii, 757, 763 Emperor Nu, 609, 610 empty numbers, 471 empty set, 82, 399 encirclement, games of, 641-729 end, 562 quiet, 609, 617-620, 629 ender, 329, 330, 353, 610 end-position, 616, 617 Endgame, 416 endgames Go, 187 loony, 577, 579 ending condition, 14, 35, 46, 47, 329 England, 905 English Solitaire board, 710, 803 enlarged flow, 204, 311 Enough Rope Principle, 16, 547, 736 entailing, 379, 396-405 Epp, Robert J., 539 Eppstein, David, 755, 767 Epstein, Richard A., 518, 668 Epstein's Game, 518–520, 535, 536 equally favorable, 35 equally uppity, 242, 245 equitable, 157-161, 169-172 equivalences Nimstring, 565 Twopins, 503, 567, 568 Erdős, Paul, 767 Erickson, Jeff, 690, 766 Ernst, M. D., 690, 765 escapade, 677, 679-681, 684, 703, 704, 709 eternal games, 46, 379 Étienne, D., 766 etiquette, 631 Evans, Ronald J., 767, 768 even, 279, 281, 282, 287 evicts!, 306 timer, 303, 305, 307, 314, 317 Even Alteration Theorem, 511

 \diamond

Even, Shimon, 224, 225, 744 Evie Parr = even parity, 545-550, 558, 569, 570, 572, 575, 578-580 evil = even, 287, 307evil numbers, 110, 111, 463, 464, 891 exactly periodic, 86 exceptional altitude decrements, 704 exceptional values, 90-92, 101, 108 excitable, 157-161, 169-172 excluded tolls, 306 excluded values, 111 exemptions, tax, 151 exit move, 829-830, 838 Exit Theorems, Beasley's, 829, 830, 837, 838 Ex-Officers Game = $\cdot 06$, 101–103, 445, 504 exp time, 763 explosive nodes, 49, 50 exponential-time algorithms, 224 exposure, death by, 927 extended thermograph, 161–162 extended thermography, 759 Extras, 14, 46, 73, 101, 133, 180, 219, 255, 292, 324, 369, 408, 442, 488, 535, 571, 607, 632, 665, 690, 725, 756, 834, 910, 958 F&G = Fox-and-GeeseMuseum, 705, 706-708 position, circled, 682, 683 Values, initial, 697, 699–701 fair board, 646, 664 fair position, 642, 645, 646 Fair Shares and Unequal Partners, 380, 394 Fair Shares and Varied Pairs, 379, 410 fairy chess, 117, 641 fairy tale, 312 Fajtlowicz, S., 632 Falada, 312-314, 319-321, 325-326 Falkener, Edward, 767 Fano's Fancy Antonim Finder, 495 far star = remote star, 230-232, 236-240, 243-251, 258, 259 farm, 146, 488, 812, 934, 952 favorite, 262, 278, 282, 284, 289 Felton, G. E., 768 fence, 488 Fencing, 488

Ferguson's Pairing Property, 86

Ferguson, Thomas S., xiii, xvii, 86, 117, 430, 454, 539, 766 Fermat powers of two, 476 Fermat, Pierre de, 476, 940 Fermat problem, first solution of, see margin of page 940 ferz = fers, 641, 642Fibonacci Nim, 517 Fibonacci numbers, 517, 518, 520, 535 Fibulations, 520, 535, 537 fickle, 413, 423-425, 429, 432, 434, 514, 533 field, 145, 312-314, 325-326, 483 Fifteen Puzzle, Sam Loyd's 864, 867 fifth column, 411 figure eight, 932 finalist, 811, 812, 815, 834 fine print, 126 finicky figures, 323 finished product, 824 finishing line, 407 firm, 413, 423-425, 429, 432, 514 first bite, 323 cousin, 101, 103-105, 107, 109, 114, 116 eaten strip, 295 home, 278, 283, 320, 322 horse stuck, 281 off, 312 one-by-one cake, 284 player wins, 28–30 strip, 295 fit 22, 251 Fitted Carpets, 479 Five-in-a-row, 378, 740 Fives Acrostic Mock Turtle, 487 Ruler, 470 Staircase, 499 Triplet, 470 fixed 335 Flags of the Allies Puzzle, 895, 920 Flanigan, James Alan, 377, 445, 539 Flanigan's Game = .34, 504 flare path setter, 936 flat, 430 Fleischer, R., 763 flip-flop, 931

floor, 51, 75, 311, 485 flow, 202, 204-206, 208, 209, 222 flow, cash, 126 Flow Rule, 201, 202, 204, 210 flower, 29, 30, 33, 35, 36, 47, 66, 67, 190, 195, 199-201, 210, 230, 240, 242, 245, 246, 364 flower garden, 190, 199, 229, 230, 240 flowerbed, 242, 244, 246 flowerstalk = stem, 36flowers, Hackenbush, 675, 695 Fool's Solitaire, 826, 835 foot, 211, 212 Football, Philosopher's, 752–755 Fondanaiche, Philippe, 533 Ford, Lester R., 205, 212, 225 Foregger, T.H., 910, 923 forging, 389 fork, 332 threat, 740 form canonical, 22 simplest, 22 standard, 101, 103-105, 107-109, 114 formation, 677, 686, 702, 705 central, 695 circled, 687 migrating geese, 703, 709 Formula, Cooling, 151 foundations for thermographs, 155 Four-in-a-Row, 737, 738 Fox Game = Hala-Tafl, 666Fox's Safe Dancing Haven, 695 Fox-and-Geese, 15, 666, 669-709 value of, 673 Fox-Flocks-Fox, 673, 674, 697, 709 FOXSTRAT, 684-687, 692, 693, 697 FOXTAC, 681-684, 687, 692 fractional atomic weights, 234–236 multiples, 255-257 Fraenkel, Abraham, 927 Fraenkel, Aviezri S., xvii, 18, 78, 224, 226, 412, 489, 517, 539, 763 France, 736, 894, 905 free, 335, 715 free your fetters!, 565 freezing point, 154, 168–172

Fremlin, David, 881 French Military Hunt, 15, 711-713, 897 Fried, Kati, 875 frieze patterns, 509-512, 535 Frevstafl, 666 Frogs, see Toads FTOZOM, 598 Fulkerson, Delbert Ray, 205, 212, 225 Full Moon, Paschal, 907 function computable, 630 pagoda, 818-823, 831 remoteness, 279-286 ruler, 98, 436, 437, 470 score \neq pagoda, 837–838 Steinhaus, 279 suspense, 286-289 Welter, 193, 506-514 fundamental insects, 598, 600 Fundamental Theorem of Zeroth Order Morbundity, 598 Funkenbusch, William, 768 furniture, redwood, 211, 212, 213-217, 222 fuse, 192, 936 fusion, 192–196 Fusion Principle, 192–195 fuzzy flowers, 29, 30 fuzzy games, 28–33, 35, 36, 39, 42, 239 fuzzy positions, 28, 32, 33 G-ness, 422 G-raph, 205–207 \mathcal{G} -sequence = nim-sequence \mathcal{G} -string, Air on a, 97 G-string, 524 \mathcal{G} -value = nim-value Gale, David, 117, 632, 640, 744, 768 gallimaufry, 68, 75, 757 Galvin, Fred, 253 galvinized games, 253, 254 game acrostic, 482-487 additional subtraction, 395 annihilation, 224 big, 65, 75, 202, 206, 222 birthdays, 417 cheap, 387 coin sequence, 635

 \Diamond

cold, 145, 299, 300, 307, 316, 317 comparisons of, 25 compendium, 109, 387 compound, 31, 299, 312 coolest, 173 eating, 136, 234, 286, 289-292, 295-297 entailed, 379, 396-405 equitable, 157-161, 169, 170, 172 eternal, 46, 379 excitable, 157–161, 169–172 fairly hot, 307 finite, 46, 115 fuzzy, 28-33, 35, 36, 39, 42, 239 galvinized, 253, 254 half-tame, 423, 435-437, 444, 445 hard, 211, 217, 223 hexadecimal, 116, 117 hot, 125, 133, 145-174, 176-185, 187, 225, 300-308, 316, 326 identification, 65 impartial, 14, 40, 56, 82, 84, 196, 220, 281, 283, 284, 289, 291-294, 296, 297, 330, 379, 396impartial loopy, 275, 379-412 impartial misère, 275, 413, 446 in the jungle, 204, 208, 210 locator, 101, 103, 464 loopy, 15, 275, 327, 334-377, 396, 758, 764 many-dimensional, 220, 488, 742 map-coloring, 38, 145 misère 281, 286-288, 290, 291-294, 312, 413-451misère Grundy's, 416-420 misère Kayles, 446–451 misère octal, 413–451 misère Welter's, 514 negative of, 33-35 no-player, 927 NP-hard, 224, 225 octal, 101-116 of encirclement, 665 of Life, 927-960 of pursuit, 15, 669-729 one-horse, 278 ordinal sum, G: H, 220partizan, 15, 65, 187, 275, 292, 295, 312, 376, 379

reduced, 446 reserves, 414 restive, 425-426, 432-438 restless, 432-435, 443 short hot, 225 simplifying, 60, 62, 63, 75, 77 subtraction, 84, 86, 87, 98, 395, 430, 442, 529 - 532switch, 121-125 take-and-break, 81, 82, 84, 86, 87, 89–96, 98, 99, 101, 103-105, 107-117 take-away, 82, 87, 100, 319 tame, 417, 422-438, 443-446, 514 tameable, 425, 446 tartan, 476-482 tepid, 306, 308, 316, 327 tiniest, 125, 126 tracking, 65, 222 trees, 40 two-dimensional, 332, 333, 473-487 wild, 430 with cycles, 376 zero, 4, 9, 33, 41 Game (see also game) Black Path, 746, 747 Bynum's = Eatcake, 136, 289 Epstein's, 518–520, 535, 536 Ex-Officer's Game = .06, 101–103, 445, 504 Falada, 320-321 Flanigan's Game = .34, 504 Fox = Hala-Tafl, 666G4G4G4G4, 757 Grundy's, 15, 96, 112, 310, 434, 439-440, 444, 690Kenyon's = .3f, 116, 117 L-, 384, 408, 412 Lewthwaite's, 747 Northcott's, 55 One-Star = 4.07, 102, 103, 604 Ovid's, 736 Put-or-Take-a-Square, 518–520 Ruler, 469, 470, 478, 483, 521 Sato's Maya = Welter's, 427Shannon Switching, 744–746 Silver Dollar Game, 491, 492, 535 The 37 Puzzle, 521

Welter's, 427, 506-515 Wythoff's, 15, 60, 74, 427 gaming tables, 492, 525, 526 Gao Xin-Bo, 765 garden 33, 199, 229, 230, 240, 241, 525 of Eden, 705, 938 Gardner, Martin, v, xvi, 18, 52, 143, 298, 502, 607, 640, 668, 713, 729, 736, 767, 768, 839, 840, 885, 919, 923, 959 Garev, Michael R., 217, 224, 226 Gasser, Ralph, 737, 764, 768 gates, logical, 941, 948 gathering, 354 Gauss, Karl Friedrich, 254 gee-up, 247 Geese, see Fox-and-Geese Geese's Landing, 697 Geese's Strategic Landing Plan, 696 GEESESTRAT, 686, 693, 697 GEESETAC, 686 Generalized Geography, 224, 760 Generalized Hex, 224 Generalized Kayles, 224 generalized thermography, 759, 764 genetic codes for Nim, 605, 606 genetic engineering, 958 genus, 422-446, 501, 504, 587, 588, 590, 597, 604, 607 Geo., 641, 642-644, 646-648, 650, 652, 654, 657, 658, 661, 663, 664 Geography, 517, 539 geraniums, 47, 199, 242, 675 Gerritse, Richard, 621, 624 Gibberd, R.W., 924 gift horse, 72, 77 Gift Horse Principle, 72, 77 Gijlswijk, V. W., 384, 412 Ginny, 387-391 giraffe, 205-206 girl meets boy, 854 glass of wine, 379-380 glass, magnifying, 151 glider, 931, 933, 935, 940–943, 944–947 gun, 935, 937, 942, 944 stream, 947-949 Go complete analysis, 755

My First, 731 stones, 642, 644-646, 648, 650-652, 654-658, 661 Go-Bang, 740 Go-Moku, 14, 740, 743 Goats, see Sheep, 863 Göbel, Frits, 118, 539, 729, 769 Godd, 195, 226 Godd's label, 195 Golav, M.J.E., 959 Golav code, 467-469 Gold Moidores, 464, 465 Goldbach, Christian, 401 Goldbach position, 382 Goldbach's Nim, 401 Golden number, 906 golden number (ratio), 75, 115, 535, 822-823 Golden Pagoda, 819 Goller, Nicholas E., 386, 408 Golomb, Solomon, xvii, 502, 642, 644, 741, 768, 895, 923, 924, 934 GONC = More) Games of No Chance, 18, 78, 118, 143, 187, 188, 225, 227, 489, 539, 540, 584, 640, 668, 710, 762-769 good = odd, 287, 307good child, 427 good move, 16, 22, 196, 397 Good, Irving John, 226 Goodell, John D., 764 GOOSESTRAT, 677, 678, 687, 695 GOOSETAC, 675, 676, 695 Goose Girl, 312 Gordon, Pritchett, 607 gosling, 899 Gosper, R. William, 931, 935, 938 Goto, David, 875 Gozreh, F., 840 grafting plumtrees, 354 Grantham, Stephen Brian, 273 graph, 145-147, 205-207, 551-553, 556, 558-562bipartite, 222, 583 complete 583 spanning tree of, 217, 222 Nimstring, 552-569, 581-583 non-planar, 551 graphic picture of farm life, 145, 146

979

grass, 40, 42, 199, 240, 242 Gray code, 858-861 Great Hall, 372 Great Tantalizer, The 892 greatly-valued carpet, 480 Great-Aunt Maude, 387 Greaves, John, 922 Grecos, A.P., 924 Greek gift, 72, 77, 547 Greek letters, loopy, 335 Green, Trevor, 533 green chain, 40 edges, 29, 30, 33, 41, 190-196, 198-202, 204-206, 210, 211, 218, 220, 221, 235, 237, 251, 330 girl, 192, 193 Hackenbush, 39-42, 190-196, 225 jungle, 198, 199, 201, 220, 221 snake, 40 tinted nodes, 204 tracks = paths, 202-210trees, 191-193 greenwood trees, 34 Gregorian calendar, 905 grey heap, 532 grin, 936 Gros, L., 924 Gross, Oliver, 744 Grossman, J. P., 755, 757, 763, 768 ground (=earth), 193, 550, 551 grounding cycles, 193, 196 group alternating, 867 symmetric, 867 grown-up picture, 43 Grundy scale 87, 90, 91, 94, 96, 101, 398-399, 406 Skavles, 91 Grundy, Mrs., 310 Grundy, Patrick Michael, 42, 56, 79, 117, 220, 221, 333, 417, 444, 454, 924 Grundy's Game, 15, 96, 112, 434, 439-440, 444, 690 misère, 416-420 wild animals, 431 Grunt, 472, 473, 481

Guibas, Leo J., 885, 924 Guiles = $\cdot 15$, 94, 101, 103, 436, 444 gun glider, 935, 937, 942, 944 thin, 947 Guy, Michael John Thirian, 111, 464, 520, 535, 621, 827, 845 Guy, Peter Richard Thirian, 841 Guy, Richard Kenneth, 18, 89, 98, 109, 117, 118, 123, 144, 386, 489, 539, 621, 640, 768, 951, Haber, Heinz, 840 hack, 675, 684, 693, 695 Hackenbush Blue-Red, 1-6, 17, 20, 23, 27, 28, 77, 78, 197.211-217 Childish, 43, 52, 157, 237 Double, 343 Green, 33, 34, 189–196 Hotchpotch, 29, 33, 36, 37, 47, 66-68, 197-210, 218-222, 229, 230, 237, 238, 242-246, 251,897 infinite, 327-332, 344, 345 is hard!, 211, 217 loopy, 343-345 number system, 78 picture, 1-5, 189-195 string, 22, 23, 77, 78, 194, 195, 327-331 von Neumann, 606 Hackenbush companion, 693 Hackenbush, flowers, 675, 695 hacking toll, 686 Haigh, J., 766 Hajtman, Béla, 924 Hala-Tafl, 666 Hales, Alfred W., xvii, 740–744, 748, 768 half-move, 4, 7, 9, 19, 20 half-off, 356 half-on, 356 half-perfect square, 886, 890 half-tame, 414, 423, 435, 444 halving, 620 of nim-values, 195 Hamlet's Memorable Problem, 826 Hammer and Sickle, 835, 836 handouts, 547 Handsome Hans, 525

 \diamond

Hanner, Olof, 188 Hanoï, Tower of, 861 Harary, Frank, 748 Harborth, H., 768 hard game, 211, 217, 222, 223 hard problems, 223-225 hard redwood bed, 217, 222 hard-hearted, 547 hard-headed, 179 hardness, 211, 217, 223, 225 Hare and Hounds, 711–729 Hares and Tortoises, 862, 915 harmless mutation, 562, 564 Harrocks, 517 harvester, 936 Harvey, Sir Paul, 924 Harry Kearev, 318 Haselgrove, C.B., 896, 924 Haselgrove, Jenifer, 924 Hashimoto, T., 762 havoc 807, 937 Hawaii, 16 Hawaiian checkers, 759 head, 807 animals' 205 girl's 193, 196 losing your, 222 severed, 205, 220-222 shrunken, 220-221 heaps, see also Nim-heaps heap, see also Nim-heaps black, 532 colored, 532 grey, 532 quiddity, 534 white, 532 heart, change of, xvii(2) hearts, 359 heat, 125, 132, 145, 299 heat, latent, 307 heating, 167, 173, 690 height, 675, 686, 687, 698, 703 Hein, Piet, 744, 843, 898 Hensgens, P., 762 Hentzel, Irvin Roy, 840 hereditarily tame, 425 Hermary, 840

Hertz oscillator, 932 heuristic discussion, 158, 159 Hex, 224, 226, 743, 744, 767, 768 hexadecimal games, 116, 117 hexagon, 744, 854, 855, 896-899, 919 Hexiamond, O'Beirne's, 896-898 Hexomino, 748 Hi, 609, 610, 617, 631 Hi-Q, 804 hi, 355, 356 Hickerson, Dean, 51, 521 Hickory, Dickory, Dock, 521 hidden secrets, 842–844, 846, 895, 897, 899 hierarchy, 95 high atomic weights, 706 high central region, 697, 698 high region, 688, 692,693, 700, 701 high scrimmage region, 697 higher ground, 209 highway, 364, 391 hilarity, 380 Hilbert Nim, 333 Hillman, A. P., 79 Hirshberg, Susan, 765 Hnefatafl, Saxon, 666 Hockey, 15 Hogan, M. S., 517, 539 Hoey, Dan, 112 Hoffman, Dean, 845, 913 Hoffman, Professor = Lewis Angelo, 768Hoffman's Puzzle, 845-846, 913-915 Hofstadter, Douglas R., 924 Hoggatt, Verner E., 79 Holladay, John C., 118, 584 Hollosi, A., 762 hollyhocks, 36 home, 277, 280, 281, 286, 289, 320, 322, 386-387 Honest Joe, 158, 159 Honsberger, Ross, 840 Hopscotch, 737 Horrocks, D. G., 539 horse, 28, 277–284, 286, 288, 289, 292, 293, 312, 320, 325, 406, 407 also-ran, 286 favorite, 278, 282, 284, 289 gift, 72, 77

outsider, 286, 289 racing, 288 remote, 278 slow, 286, 289 working out, 28 Horsefly, 395, 411 Hoshi, Y., 766 hot, 300, 301, 322, 326, 355-356, 710 hot, 125, 133, 145, 149, 151, 171, 173, 225, 300, 304, 307, 316 battle, 145, 299-308, 310, 322, 327 coating, 322 component, 299–300 game, 125, 133, 145, 171, 173, 225, 300, 307, 316, 317, 326 position, 149, 304 work, 316 Hot Atomic Weight, 706 Hotcakes, 299-304 Hotchpotch, Hackenbush, 38, 47, 66-68, 198-202, 204-206, 208-211, 225, 230, 238, 251, 896 Hotstrat, 188 Hound-Dog Position, 724, 729 Hounds, see Hare, 711–729 house and garden, 33 House of Cards, 357-361 Howells, D. F., 539 Huddleston, Scott, 640 Hudson, Paul D.C., 534, 539 Hungarian Cube, 868-876, 917-918 Hunter, J.A.H. "Fun with Figures", 924 Huntjens, M. P. H., 740, 767 Hutchings, Robert L., 615–617, 631, 827 Hyperspace Tit Tac Toe, 768 Icelandic sagas, 666 idempotent, 693 Igusa, K., 761, 766 Iida, Hiroyaki, 762, 765, 766 Iida, Takahito 766 illegal, 320, 322, 404 Imai, Hiroshi, 766 imminent jump, 9-11 impartial, 15, 40, 56, 82, 84, 196, 220, 291, 330, 379, 396 Cutcakes, 284, 292, 294

Domineering = Crosscram, 142, 144, 298

982

horse-moving, 283, 289 infinite tolls, 314, 315, 317, 319, 320, 326 loopy games, 359-363, 366-375 remoteness, 296 incentive, 147, 148, 256-259 incomparable, 35 indefinite F&G boards, 674, 686 Independence Day = Doomsday, 903induction, 115, 234 index decremented, 705 inequalities for stoppers, 351 Inequality Rule, 348–349 infinite delphinium, 47 ender, 329 frieze pattern, 512 geranium, 47 Hackenbush, 327, 332, 344 Nim, 330 nim-values, 636, 640 ordinal numbers, 329 remoteness, 381 repetition, 14, 384 Smith theory, 333 tolls, 315-320, 325-326 infinitesimal 36, 169, 170, 171, 173, 229, 539, 690 infinitesimal numbers, (See **over**) infinitesimals, in F&G, 690 loopy upon^{*}, 708 loopy positive, 693 negative, 708 pure, 706 infinitesimally close, 151, 152, 154 infinitesimally shifted, 177–179, 198 infinitesimally small, 158 infinity, 329, 334, 391 Inglis, Nick, 51 initial F&G Positions values of, 688 initial F&G Values, 698–701 initial values, 698, 699 ink, waste of, 382 input, 942, 944 Instant Insanity, 892 integral, 167–174, 176–179, 334, 366–367

Intermediate Value Theorem, 426, 438 interval, confusion, 121, 149-151, 158, 163 intriguing women, 524 inverting Welter's function, 510–512 invoices and cheques, 126 inward move, 677 irrational, 610 irregular values, 90-92, 101, 108, 187 ish = Infinitesimally-SHifted, 177-179, 198 Isidor, Bishop of Seville, 768 isolated circle, 684 isomorphism, 469, 473, 491, 529, 586, 604, 632, 732, 733 Isotopy Extension Principle, 853 Italy, 905 Itoh, H., 766 Iwata, S., 769 Iwata, Shigeki, 765 Jacobson, Guy, 602, 607 Jaffrey, A., 539 Jam, 732 Japan, 894 Japanese chess, (See Shogi) Jelly Beans = $\cdot 52$, 428, 444 Jam, 732 Jenkyns, Thomas Arnold, 539 Jewett, Robert I., 740-744, 748, 768 Jewish New Year, 909 jig-saw puzzles, 895, 900–903 Ji-Hong, Shen, 765 Jimmy, 387–391, 413 Jin-Bai, Kim, 764 Jocasta, 597, 607 Johnson, David S., 217, 224, 226, join, 277-298, 300 rapid, 286, 289 slightly slower, 283 slow, 286-291, 300, 302 joints, 562, 564, 566 joker, 358 jump, 8–12, 63–65, 75, 76, 127–131, 133-137, 803-812, 817, 820, 825-826, 828-832, 835, 837-838 jumpee, 11 jumper, 9 jungle warfare tactics, 210

jungle clearing, 222 green, 198, 201, 220, 221 parted, 201, 202, 209 sliding, 198 smart game in, 209 tracking, 204-208, 221, 222 unparted, 210 k-number, 471 Kajihara, Y., 762 Kano, M., 118 Kao Kuo-Yuen, 188 Karp, Richard M., 217, 224, 226 Kasai, Takumi, 765 Katzenjammer Puzzle, 892 Kayles = $\cdot 77$, 15, 81, 82, 88–92, 95, 98, 99, 109-112, 224, 417, 424, 425, 431, 444-452, 500, 504, 566-569 Kayles Dawson's = $\cdot 07$, 15, 89, 92, 95, 101, 260, 438, 439, 444, 500, 504, 566, 569 Double, 99 Misère, 446-452 Quadruple, 99 Triplicate Dawson's, 260, 271 Kayles-vine, 567, 568, 579, 690 Keller, Michael, 756 Kenyon's Game = $\cdot 3f$, 116–117 Kenyon, John Charles, 109, 115, 118 kickback reaction, 947 Kierulf, Anders, 764 killing mutation, 562 Kim, Yonghoan, 764 Kindervater, G. A. P., 384, 412 king centred, 657, 658 cornered, 658, 660, 661 edge-charging, 650, 652 sidelined, 661, 663 rook versus king, 667 King's Horses All the, 277-283, 286-289, 292, 293 Some of the = Falada, 312 King, Kimberley, 222 Kinggo, 646-664 Kite Strategy, 243, 244, 246 Klarner, David A., 539

984

knight, 278, 395, 411, 641 Knight, White, 56-59, 278 Knights of the Square Table, 756 Knuth, Donald Ervin, xvii, 18, 51, 473, 489 Kolpin, Garrett, 758 Konakis, 666 Konane, 690, 759, 760, 765 Körner, Thomas W., 665 Ko (see also loopiness), 377, 758 Kotzig's Nim, 515, 517 Kotzig, Anton, 516, 539 Kraitchik, Maurice, 710, 840, 924 Kriegspiel, 15, 667 Kuperberg, Greg, xvii Kutz, Martin, xvii, 748 L-game, 384, 408, 412 L-package, 809, 832 L-purge, 809, 837 Lacrosse, 15 ladder, 386-388, 504, 897, 899, 921 Lake, Robert, 763 Lam, T. K., 607 landing fox, 693-695 geese, 697 Landman, H. A., 764 L'Ane Rouge, 875 large boardwidth, 710 Largest Nim, 532 Lasker, Edward, 99, 118, 768 Lasker's Nim, 99, 113, 114 last cut, 294 last home, 286, 289 last horse, 286, 289 last move, 171 last mouthful, 289 last player losing, xvii(2), 413–452 last player winning, xvii(2), 2, 8, 12, 14 latent heat, 132, 307 latent loopiness, 375, 391 latent phase change, 168 lateral thinking, 384, 412 Latin squares, 497 Lee, Chester C., 738, 960 Left, 2 boundary, 154-156, 164, 165, 169, 170 excitable, 158-159

remoteness, 279, 281-283 stop, 149, 150, 152, 161 tally, 301, 302, 304, 310, 314, 327 Lefty, 25, 51, 234, 284, 289, 300, 308 leg, 211-213, 215, 222 Lehman, Alfred, 768 Lehman's switching game, 767 Leibniz, Gottfied WIlhelm von, 818, 840 Lemma Norton's, 220, 243, 244, 354 Snort, 180 Lemon Drops = $\cdot 56$, 428, 444 Lenstra, Hendrik Willem, 461, 489 Les Pendus, 736, 737 Lester, W.E., 923 Let them eat cake!, 289 Levine, J., 765 Levy, D. N. L., 764 Levy, Silvio, xvii Lewis, Anglo, 768 Lewthwaite, G.W., 747 Lewthwaite's Game, 747 Li Shuo-Yen, Robert, 343, 373, 377, 539 lice, infestation with, 600 Lichtenstein, David, 224, 226, 763 Life, 15, 53, 927-960 computer, 940, 958 configuration, 935, 940, 957 cycle, 930-932 environment, 958 history, 928, 935–937, 940 is universal!, 957 pattern, 934, 940, 958 sole aim in, 958 space, 959 still, 929, 932 Life's but a game, 939 Life's Problems are Hard, 939 Life's unpredictability, 934 Lifeline, 960 Lifenthusiasts, 930 light positions, 677 lightening, 684, 686, 694 lightening move, 677, 678, 685 lightning bolts, 50 limbs, stretching, 565

line of scrimmage, 687 real number, 23 wigglv, 560, 565, 582 Linnaeus, Carolus, 666, 668 Litchfield, Kav P., 924 little safe, 621 live cell, 927-929, 934 live spots, 598, 601, 602 Lo. 609, 610, 617, 631 lo. 355-356 Lo-Shu, 886 loaf, 929, 935, 943 Loeb, D. E., 766 logical gates, 941, 942, 948, 949 lollipops, 237, 238, 240 long, 546 barge, boat, 928 chain, 543, 545-547, 549, 550, 552-556, 561, 563, 564, 569, 574, 575, 577-579, 829 cycle = loop, 553-555, 565path, 562, 563 period, 108, 109 ship, 929 Long Chain Rule, 546, 549, 550 loony, 322, 397-407, 558, 561-564, 576, 577, 579, 580 loony endgames, 577 Loony Loop, 849, 852, 861 loop, 929 looping the, 947 long, 554 short, 554 loopiness (see also kos, superkos), 361–363 blatant, 391 degree of, 360 latent, 375, 391 patent, 376, 391 Loops-and-Branches = .73, 586loopy, 758 component, 390, 410 game, 327, 334–377, 396, 758, 764 Hackenbush, 343–345 infinitesimal (See over, upon^{*}), 708 number (See over), 341, etc. option, 410 positive infinitesimal, 693

position, 389, 408, 758 value, 387, 389-391, 758 lose control, 574 lose slowly, 278 lose your shackles!, 565 losing, last player, xvii(2), 413-452 Lost World, 413-414 louse, 600 Lovász, L., 766 low region, 688, 698 Lower region, 697 Loyd, Sam, 82, 118, 864 Lucas, Edouard, 588, 607, 713, 729, 768, 840, 862, 924 Lucasta, 585, 588-597 Lucky Seven Puzzle, 864-866, 876, 916 lucky star, 246, 249 Ludo, 14, 15 ludus terni lapilli, 736 Lukewarmth Commandmant, 307, 310 lush delphinium, 675, 677, 681, 684 Lustenberger, Carlyle, 768 m-plicate, 98 MacMahon, Major Percy A., 924 MacMAhon jig-saws, 900-902 squares, 899-902, 921 superdominoes, 899-902 triangles, 899 Macmillan, R. H., 768 Madachy, Joseph S., 924 Magic cube, 866-874, 918-919 Fifteen, 733 hexagon, 890, 920 Mirror, 749, 752–754, 756–758 Movie, 854–855 square, 733, 886-891 tesseract, 889 Maharajah and Sepoys, 710 Mahomet, 854 making tracks, 202, 204-209, 221, 222 management of cash flow, 126 Mancala, 761, 766 many-way Maundy Cake, 220, 221 map, 37, 145, 414, 879-884, 910-913 Mark, 4.24: Matthew 13.12, 405

Markeert, D., 225 markup, 161 Markworthy Commandment, 317 Martin, Greg, xvii, 644, 645, 668 Mason, R. E. A., 764 mast, 152, 154, 163 mast value = mean value, 165, 166mate, 506, 507 Mateescu, A., 764 Mathematical Go, 759 mathematician, 936 Mather, MIchael, 910, 923 Mathews, Charles, 764 Mathieu, Emil, 469 Mating Method, 506, 507 mattress, 215 Mauhin, Patrick, 27 Mauldon, James G. 79 Maundy Cake, 26, 27, 51, 196, 220, 221 max, 288 maxim, 278, 279, 287 maximal flow, 202, 204-206, 209 May-Day, 903 Maya Game, Sato's = Welter's Game, 427 Mayberry, John P., 533, 539 McCurdy, S., 763 Meally, Victor, xvii(3), 506, 920 mean value, 149, 151-154, 165-168, 172, 174, 178Meander, 747 Meier, Kersten, 924 Melencolia I., 888 men = Phutball players, 752Merrilees, 737 Method Magic Mirror, 849–854 Magic Movie, 854, 855 Mating, 506, 507 Methuselah, 937 mex = Minimum-EXcludant, 56, 82, 83, 418, 556, 558 migrating geese formations, 703, 709 Mill, 737 Miller, J.C.P., 924 Milnor, John, 188 Milton, John, 19, 145 minimal criminal, 194, 214

minimal spanning tree, 216, 217, 224 Minsky, Marvin L., 954 miny, 126-129 minv, 353, 356, 359, 360, 362 Miracle Octad Generator, 468, 469 Mirror, Magic, 849, 852-854, 856-858 misère birthdays, 417 Contours, 587, 588 Cram, 504–506 Cutcakes, 284, 292, 294 Eatcakes, 289, 291 Grundy's Game, 416–420 Kayles, 431-432 Loops-and-Branches, 586 Lucasta, 590-597 Mex Rule, 418 Nim, 413 octal games, 443–445 play, 15, 86, 278, 280-285 287-294, 312, 413-452, 610, 636 remoteness, 283 Rims, 586 Sprouts, 602 Stars-and-Stripes, 604 suspense, 289 theory, 415-419 Twopins, 501 unions, 312 Welter's Game, 514, 515 Wyt Queens, 427 mistake, inevitable, 382 mixed, 335 Möbius, August Ferdinand, 467 Möbius transformation, 467, 867 Mock Turtle, 461–465 Fives, 470 Theorem, 464–466 Mock Turtles, 463–465 Modular Nim, 539 Moebius, 464-467 Nineteens, 470 Moews, David, 52, 188, 224, 539, 764 MOG = Miracle Octad Generator, 468, 469 Mogul, 464-469 Moidores, 464, 465 Mollison, Denis, 598

 \diamond

986

 \diamond

money, 171 moneybag, 492 Monopoly, 15 moon, 379, 903-906 age of, 905 new, 906 paschal full, 906 Moore and More, 534 Moore, Eliakim H., 427, 539 Moore, Edward F., 939, 960 Moore, Thomas E., 607 Moore's Nim_k , 427, 533 Morelles, 737 Morgenstern, Oskar, 539 moribundity, 598 equation, 601 Morra, Three-Finger, 15 Morris, Lockwood, 225 Morton, Davis, 640 mosaic, 936 Moser, Leo, 667, 732, 742, 743, 768 motley, 469, 477, 483 Mott-Smith, Geoffrey, 502, 768 mountain, 7 purple, 197, 198 move 14, 40 abnormal, 325 alternating, 46, 47 bad, 16, 547, 818 bonus, 405 chance, 14 complimenting, 379, 405–407, 541, 552 consecutive, 405 darkening, 677, 678, 685 entailing, 379, 396-405 equitable and excitable, 161 exit, 829-830, 838 five-eights of a, 20 futile, 654 good, 16, 196, 397 half, 4, 19, 20 horse, 406 hotter, 173 illegal, 320-322, 404, 636 inward, 677

legal, 404 lightening, 677, 678, 685 loony, 322, 398, 400, 411, 561-564, 576-580 non-entailing, 397-398, 400 non-suicidal, 322, 562 normal, 325 outward, 677 overriding, 312, 314, 317, 319, 320, 326 pass, 281, 283, 284, 286, 289, 292, 293, 294, 338, 352, 355plausible, 654 predeciding, 312, 320, 321 quarter, 6, 20 repainting, 343 reversible, 55, 56, 60, 62-64, 66, 70, 71, 75, 77, 126, 212, 213, 415 reversible misère, 415 reverting, 425 strategic, 652 stupid, 636 suiciding, 312, 320 sunny, 397-401, 404, 407, 411 tactical, 652 temperature-selected, 132 three-quarter, 17 trailing, 402 worthwhile, 213-216 move set, 515-517 Movie, Magic, 853-854 Mr. Cutt and Mr. Shortt, 745, 746 Mrs. Grundy, 310 Mühle, 737 Müller, Martin, 18, 188, 762, 764 multiples of up, 71, 242, 247, 256, 258 multiples, fractional and non-integral, 236 multiplying pegs, 812–816 multum in parvo, 937 Munro, Ian, 226 Murray, H.J.R., 668, 737, 768 Muscovites, 666 museums, 705 F&G, 705-708 musical series, 96, 117 mutation, 562, 959 harmless, 562 killing, 562 Myhill, John, 960

n-dimensional k-in-a-row, 742 N-positions, 83, 279, 381-384, 408, 410, 439, 447-451, 466, 516, 517, 519, 520, 535, 537, 572, 573, 589, 592, 594, 613, 615-618, 620, 627, 631, 632, 635, 636, 638-640, 751 n-theorem, 615 Nagai, Avuma, 766 Nakamura, Teigo, 764 Nash, John F., 744 Nasik squares, 891 negative, 28 charge, 254, currency, 610 infinitesimals, 708 numbers, 19, 147, 148, 330 of a game, 33-35 positions, 28, 29, 68 negs = negpegs, 831neighbor, 601 Neumann, John von, 539, 606, 960 von Neumann Hackenbush, 606 Nevman, Abraham, 117 Nim 15, 40-42, 53-59, 82-84, 113, 114, 173, 240, 387, 454, 491-493, 508, 533, 534, 586, 604, 605 Antipathetic, 493 Bounded, 517, 518 Chinese = Wythoff's Game, 427Double Duplicate, 114 Duplicate, 114, 116 Entailing, 400 Fibonacci, 517 genetic codes for, 605, 606 Goldbach's, 401 Hilbert, 333 in hot games, 173 infinite, 330–333, 636, 640 Kotzig's, 515–517 Largest, 532 Lasker's, 99, 113, 114 misère, 413, 416, 418, 520 modular, 539 Moore's, 427, 533 Poker, 55 Similar Move, 496–498 Smallest, 532 Sympathetic, 494–496

Triplicate, 114, 116 two-dimensional, 332-333 Welter's, 427 Nim Addition Rule, 59, 73, 199, 390, 418, 556 nim-addition, 58, 59, 73, 74, 89, 90, 109, 110, 115, 116, 191, 196, 199, 246, 390, 418, 446, 463, 464, 469, 471, 473, 475, 476, 497, 506-508, 510, 556, 558, 605, 606, 752 Nim-heaps, 41, 42, 55-59, 82-84, 389, 393, 413, 418-424, 426, 430, 434, 436, 442, 462, 469, 491, 504, 514, 605 ambivalent, 426 bogus, 56, 57 nim-like game, 751 nim-multiplication, 475-478 Nim-position, 41, 240, 387-389, 425, 426, 446, 491, 507, 605 nim-product, 475-478, 481 nim-sequence, 82-87, 89, 93, 94, 98-102, 104-108, 116, 529, 530, 587 Nimstring, 552-556, 558-562, 564-567, 569 problem, 552 nimstring arrays, 581 nim-sum, 58, 59, 73, 74, 82, 89, 90, 91, 110, 112, 115, 462-464, 471, 473, 475, 481, 498, 507, 586, 605, 606, 751 nim-values, 82-94, 96-102, 104-117, 191-196, 230, 387-408, 410, 411, 425, 426, 430, 434, 436, 442, 444, 446-450, 462-466, 469-488, 498, 506, 507, 514, 521, 529-533, 537, 557-560, 567-569, 579-583, 586-590, 597, 604-606, 636, 640, 751, 752 addition, 390 doubling, 94 duplication, 94, 98, 99, 114, 116, 444-445 halving, 195 periodic, 84, 86, 91, 92, 94, 98, 99, 101, 103-105, 107-110, 112-117 reflected, 109 relevant, 396 replication, 98 nimbers, 40-42, 56-58, 65, 74, 110, 199, 200, 231, 258, 262, 397-401, 405, 406, 418, 556,558, 581, 690 adding, 42, 58 infinite, 330 Nim_k, Moore's, 427, 533

 \diamond

Nine-in-a-row, 740, 741 Nine Men's Morris, 736, 737, 768 No Highway, 364-367 Noah's Ark theorem, 423, 432–435, 443 no-player game, 927 node-disjoint cycles, 578 nodes Col, 47–51 explosive, 49–50 game positions, 42, 341-343, 354-358, 369-370 Hackenbush, 191–194, 201, 202, 204, 222 Nimstring, 554, 560–562, 564, 576, 578, 581 Snort, 146, 180, 181, 183 tinted, 204, 206, 208, 210 untinted, 204, 206, 209, 210 non-abacus positions, 514 non-arithmetic periodicity, 115 non-number, 147, 148, 159, 160 normal move, 325 normal play, 12, 14, 278, 279, 281, 282, 284, 286, 287, 291, 293, 300, 312 normal Soma pieces, 847 Northcott's Game, 55 Norton, Simon, 145, 168, 188, 220, 243, 244, 247, 256, 354, 496, 520, 535, 667 Noshita, Kohei, 766 NOST, 756 NOT gate, 941, 942, 948-949 Noughts-and-Crosses = Tic-Tac-Toe, 14, 731-736 novice, 382 Nowakowski, Richard Joseph, xvii(1), xiv(2), xviii(3), 516, 517, 539, 540, 584, 640, 668, 755, 757, 762-769 NP-complete, 217, 224-227 NP-hard, 224, 225, 577, 755 Nu, 609 nude, 675 Number Avoidance Theorem, 147–149, 183 number system Hackenbush, 78 tree and line, 24 numbers, 22, 119, 300, 314 canonical form, 22 empty, 471 evil, 110, 463, 464, 485, 891

Fibonacci, 517, 520, 535, 537 infinite ordinal, 329 infinitesimal (See over) k-, 471 loopy (See over) negative, 8, 19, 708 odious, 110, 111, 463, 464, 471, 472, 481, 891 overheated, 176 simplest, 19, 21, 22, 305, 307, 314 Surreal, 18 suspense, 286-289, 292, 294, 315, 317 thermographic thicket of, 176 triangular, 254 whole, 19 nutcrackers, impossible, 379 O-positions, 382, 384, 390–391, 408, 410 O'Beirne, Thomas H., 79, 118, 452, 515, 539, 729, 768, 769, 893, 924 obtuse triangle, 254 octad, 467-469 octal games, 101, 103-105, 107-116, 267, 443-445octal notation, 101, 464, 465, 586, 686 odd, 279, 281, 282, 287, 303, 305, 307, 314, 317 odd admits!, 306 odious numbers, 110, 111, 463, 464, 471, 472, 481, 485, 891 Odlyzko, Andrew M., 883, 924 off, 336-340, 353, 355-358, 674, 704, 710 Off-Wyth-Its-Tail, 402 Officers = $\cdot 6$, 95, 439–440, 444 offside, 336-337, 340, 344-345, 348, 354-355, 369 - 370Ohara, E., 766 O'Hara Frank, 876 Omar, 42, 72, 109, 137, 187, 262, 517, 518, 529, 533, 537, 664, 667, 688, 702, 839, 914 on, 336-338, 353, 673, 686 On-the-Rails, 406–407 ONAG = On Numbers and Games, 18, 22, 48, 52, 78, 117, 138, 143, 188, 225, 233, 235, 262, 278, 298, 377, 417, 431, 452, 475, 478, 489, 506, 510, 539 oNe = Weak or Strong place, 713One-for-you, Two-for-me, ..., 319 one-horse game, 278

One-Star Game = 4.07, 102, 103, 604 One-step, Two-step, 529 One-upmanship Rule, 242 ono, 355-357 onside, 336-337, 340, 344-347, 354-355, 369-370 oof. 355-358 open, 382, 384, 390-391, 408, 410 opposition, 714-717 option, 14, 154, 155 best, 303 dominated, 62, 63, 75, 77, 149 Left, 31, 32 loony, 558 loopy, 408 non-loopy, 389 questionable, 387 reversible, 60, 62-64, 70, 71, 75, 77, 387 Right, 31, 32 suicidal, 317, 558 worthwhile, 213-215 optional extras, 84 OR gate, 941, 948-949 ordinal numbers, 329 ordinal sum, 220 Orman, Hilarie, xvii, 749, 768 Othello, 760, 765 outcome, 28, 35 classes, 28, 84 of sum, 31, 32 output, 942, 944 outsider, 286, 289 outward move, 677 over 341, 353, 673, 674, 686, 693 overcrowding, 927 overheating, 174, 176, 187, 366, 690 overriding, 312, 314, 319, 326 Ovid, 736, 769 Ozanam, 924 Ozery, M., 79 P-positions, 83, 113, 279, 381–385, 408, 410, 411, 416, 417, 419-424, 426, 431, 447-451, 462, 464-469, 493-499, 508, 516-522, 526-529, 535, 537, 548, 572, 573, 587, 589-592, 594, 597, 602, 611, 613-616, 620-

639, 751, 752 p-theorem, 615 packages, 807-812, 815, 834, 836, 837 pagoda functions, 818-823, 831 Pairing Property, Ferguson's, 86, 442, 530, 531 pairs, restive and tame, 425 pale twig, 344 panacea, panache, 809 pandiagonal squares, 891 Păun, Gh., 764 Pandora, 346, 548 paradox, 534, 539 paradoxical dice, 886 pennies, 885 Parikh, 764 parity, 73, 191, 194, 526, 545, 698, 704, 897 Parity Principle, 191, 194 Parker, Richard, 535 parody, 286 Parotty girls, 545 parted jungle, 201, 202, 209 particles, 168-170 partizan, 15, 65, 187, 292, 295, 312, 376 party trick, 856 Parville De $\lambda = Lucas$?, 862, 924 Paschal Full Moon, 906 pass, 283, 284-286, 288, 289, 292-294 Patashnik, Oren, 769 patently cold and hot, 307 patently loopy, 376, 391 Paterson, Michael Stewart, 598, 898, 902, 903 path, 202, 562, 563, 746, 747 Path of Righteousness, 548 paths = tracks, 193, 202, 204, 222Patience = Solitaire, 15, 804Paul, Jerome L., 769 Paul Wolfgang J., 226 Paulhus, Marc, xvii paw mark = block, 936 pawns, 89 pearls, 522, 523 Peek, 224 Peg Solitaire, 15, 801–840, 897 Pegg, Ed, xvii Pegity, Pegotty, 740 pegs, 803-840, 859, 860 pencil-and-paper game, 1, 604 Penfield, Wilder, 414

 \diamond

990

 \Diamond

Penrose, Roger, 876 pentadecathlon, 934, 937 pentominoes, 749, 895-896, 934 perfect square, 886 periodicity, 89-91, 96, 99, 100-102, 104-109, 111-115, 117, 236, 292, 310, 311 arithmetic, 99, 113–117, 144, 187 Dawson's Chess, 90-91 Domineering, 144, 187 Eatcakes, 289, 291 exact. 86 Guiles, 94 Kayles, 91, 92 octal games, 101, 103-105, 107-109, 113-115subtraction games, 84, 86 ultimate, 99, 101, 112, 289, 522, 529, 752 petal, 33, 47, 66, 67, 199, 201, 240, 245, 344 Petrie, Douglas G., 937 pharisees, 601, 602 phase change, 167, 168 phase change, latent, 168 Philosopher's Football = Phutball, 224, 752-755Philpott, Wade E., 839 Phutball, 224, 752–755 picture, 1, 2, 190, 192, 467, 468 of farm life, 145, 146 piebald node 47-48 piebald spot, 145-147 pink twig, 344 place(Zero, oNe, Two), 712–719, 728, 729 placing plumtrees, 354 Plambeck, Thane E., 454 play, 334 misère, 15, 278, 280-295, 312, 413-452, 590-597, 610, 636 normal, 12, 14, 278, 279, 281, 282, 284, 286, 287, 291, 293, 300, 312 player first, 28-30 second, 28-30 symmetrical, 547 playing backwards, 817

playing the averages, 167, 173 Pless, Vera, 489 Plov Blue Flower, 199, 201, 240, 242 Blue Jungle, 201, 210 Red Jungle, 210 Plugg = Cram, 141, 143, 298, 502-506plumtrees, 352, 354, 357 poisoning, 754 Poker, 15 Poker-Nim, 53-55, 491 pole, barber's 936 Policy, Searching, 894-895, 898 Policy, Temperature, 124, 125, 131, 132 Pollak, Henry Oliver, 740 Polyiamonds, 895 polyominoes, 120, 139–142, 748, 749, 895 pond, 929 Pond, I. C., 539 Poole, D. J., 517, 539 PORN, 630, 631 position, 14 abacus, 513-515 active, 149, 150 ajar, 408-410 circled, 681, 684, 685, 696 closed, 408-410 cold, 301, 304–306, 322 complementary, 816 criticial, 687 daggered, 677 dark, 677 Domineering, 120, 121, 139, 142, 144, 153, 177 exceptional, 717 F&G. 682-683 fair, 642 fickle, 413, 424, 425, 429, 432, 434, 514 firm, 413, 423, 425, 429, 432, 435, 514 fuzzy, 28, 32, 33 Goldbach, 382 hot, 149, 304 Hound-Dog, 724, 729 light, 677 loony, 322, 397, 407, 557, 558 loopy, 389, 408

N-, 83, 279, 381-384, 408, 410, 439, 447-451, 466, 516, 517, 519, 520, 535, 537, 572, 573, 589, 592, 594, 613, 615-618, 620, 627, 631, 632, 635, 636, 638-640, 751 negative, 28-30, 68 non-abacus, 514 O-, 382, 384, 390-391, 408, 410 \mathcal{P} -, 83, 113, 279, 381–385, 408, 410, 411, 416, 417, 419-424, 426, 431, 447-451, 462, 464-469, 493-499, 508, 516-522, 526-529, 535, 537, 548, 572, 573, 587, 589-592, 594, 597, 602, 611, 613-616, 620-639, 751, 752 positive, 28, 68 Scare'm Hare'm, 719-721 stable, 705 starting, 14 sunny, 397-398, 401-404 tepid, 306 terminal 2, 7, 29, 41 unstable, 705 positive charge, 263 house, 38 house and garden, 33 positions, 28, 68 posy, 29, 30, 33 POT(S), POTSHOT, 738-740 poultry, 755 predecider, 312, 320, 321 pretending, 421, 439, 445 Price, B.D., 924 Prim, 98, 404, 442 Prince Charles, 525 Princes' Code of Behavior, 537, 538 Princess and the Roses, 524–528, 537, 538 Principle Bogus Nim-heap, 56 Colon, 191, 193, 194, 220, 355 Complimenting Move, 405 Death Leap, 127-130, 135 Enough Rope, 16, 547, 736 Fusion, 192-195 Gift Horse, 72, 77 Isotopy Extension, 855 Parity, 191, 194 Star-Shifting, 250, 258

Translation, 149, 153, 158 Uppitiness Exchange, 246 Pritchett, Gordon, 607 problem Amazon, 756, 761 Bumble-Bee, 827 Clobber, 758, 761 Deader Dodo, 827 Dots-and-Boxes, 570, 578-580 escapade, 681 Fox-Flocks-Fox hard, 223, 224, 827, 831, 948, 957 Hamlet's Memorable, 826 Konane, 759, 760, 762 Nimstring, 552 reversal, 809-810, 820, 826-827, 831-832, 839 Three B'ars, 827 unsolved F&G, 710 Prodigal Son, 825, 826 product acrostic, 482-487 finished, 824 Gross National, 826 nim-, 475-478, 481, 488 raw, 824 ugly, 483 productive, 824 professional boxer, 548 Professor Hoffman (Angelo Lewis), 768 profit, 161, 492, 843 profit-consciousness, 171 program cycle, 316 projective, 430 proof, 115, 147, 148, 165, 166, 183, 188, 201, 212, 213, 216, 217, 220, 221, 224, 236, 240, 248-250, 256-259, 340, 348-350, 370-373, 408, 435, 442-443 Proviso, Endgame, 416–417 pruned, 675 pruning plumtreees, 354 pseudocorner, 664 PSPACE, 224, 227, 760, 765, 766 -complete, 224, 226 -hard, 224 puffer train. 937 pulsar CP, 930, 936

pulses, 940, 941 purchasing contract, 127 pure infinitesimals, 706 purges, 807-809, 828, 834, 837 purple mountain, 198, 199 pursuit, 15, 298 Put-or-Take-a-Square = Epstein's Game, 518-520, 536 putative nim-value, 471, 472 Puzzle Century, 877, 878, 881-885, 919 Century-and-a-Half, 878, 885, 919 Donkey, 877, 878, 881, 885 Fifteen, 864, 867 Flags of the Allies, 895, 900-902 Hoffman's, 847, 848 jigsaw, 895, 900-903 Lucky Seven, 864-867 sliding block, 877-885, 919 Solitaire-like, 862 The Thirty-Seven, 521 Tricky Six, 867, 868, 876 wire and string, 849-861 Wir(e)s(tr)ing, 851 Quadraphage, 143, 298, 642, 643, 664, 668 Quadruple Kayles, 99 quality of quaternity, 528 Quam, 534 quantity beats quality!, 201 Quaquaversal Quadrimagifier, 888 quarter-infinite board, 284, 292, 293, 312, 642, 644, 654, 667, 751 quarter-move, 6, 91, 20 quarter-perfect square, 889-890 queen bee, 937 Queen Elizabeth, 905 quiddity heap, 534 quiet end, 609, 617 position, 617, 618, 620, theorem, 618-620, 624, 629 quietly excludes, 617 quietus, 619 quintessential quinticity, 528 quintominal dodecahedra, 921 quintominoes, 900 quotation marks = eccentric cases, 231, 232, 251 - 253

 \Diamond

rabbit, 711 Rabin, Michael O., 640 Radcliffe, William, 920 rademacher, rado, radon, 371 Rademacher, Hans, 571, 584 Rails, 585, 586 randomness reigns, 528 range, 497, 498 trifurcated, 705 rank, critical, 710 rapid join, 286, 289 rare values, 110-112 raw product, 824 reaction kickback, 947, 948, 953, 955 vanishing, 941-943, 947, 948, 951 Read, Ronald C., 924 reader. assiduous, 42 gentle, 706, 812 more mathematical, 362 persevering, 90 skeptical, 130 real number line, 24 rectangles, 25-27, 141, 234-236, 632, 664, 895 Red edges, 2, 198, 230, 237 Red tinted nodes, 48, 204 Red twig, 213 Red-Blue Hackenbush, see Blue-Red Hackenbush Red Jungle Ploy, 210 reduced game, 446 redwood bed, 211, 213, 217, 222 furniture, 211, 217, 222 tree. 214–217 twig, 213, 216 Reeve, John E., 924 References, 18, 52, 78, 117, 143, 188, 225, 262, 298, 377, 412, 452, 489, 538, 584, 607, 640, 668, 710, 729, 762, 840, 923, 960 reflexion of Nim-values, 109 register storage, 951 region, 692 delta, 697 high, 688, 692, 693, 700, 701 high central, 697, 698

region, low, 698 lower, 697 transition, 702 transitional, 701 trifurcated, 701, 702 Welton's, 688, 693, 695, 697, 699–702, 704, 705Reid, Michael, xvii Reisch, Stefan, 226 Reiss, M., 814-815, 818, 840 remote, 243 horse, 278, 280, 281-283 remote star, 230-232, 236, 237, 239, 240, 243-252, 258, 259 Remote Star Test, 246-249 remoteness, 173, 279-286, 289, 292, 293, 295-297, 315, 317, 381-386, 519, 520, 535-537, 725-728,740 even, 279, 281, 282, 381 horse's, 281-284, 292 infinite, 381-382 Left, 279, 281-283 misère, 280–285, 292, 293, 295 normal, 279 odd, 279, 281, 282, 381 Right, 279, 281–283 rules, 279, 281 repainting moves, 343 replication of nim-values, 98, 529 reproduction of computers, 958 resetting the thermostat, 185 reources, available, 824, 826 restive, 414, 425-426, 432-438, 597 restless, 414, 432-435, 443 Restricted Translation Rule, 261 reversal problem, 809, 810, 817, 818, 820, 826, 827, 831, 832, 839 Reverse Hex, 767 Reverse Othello, 765 Reversi, 765, 760 reversible moves, 55, 56, 60, 62-64, 66, 70, 71, 75, 77, 126, 212, 415 reverting moves, 425 rhombs, 897, 898 Richardson, J., 765

Right, 2 boundary, 154-156, 164, 169, 170 excitable, 158, 159 remoteness, 279, 281-283 slant, 158, 159 stop, 149, 150, 152, 163 tally, 301, 302, 304, 310, 314, 317 Rims, 585, 586 ring, 854, 855, 858, 860, 928 rings and strings, 850, 851 Chinese i =Scandinavian ?, 856–861 Rip Van Winkle's Game = Kayles, 82 ripening plums, 354 Rita, 25, 51, 234, 284, 289, 300, 308 Robertson, Edward, 226 Robinson, Raphael M., 914, 924 Robson, J. M., 763, 764 Rodgers, Tom, 607, 765 Rognlie, R., 762 Rolling Stones, 346 Romantica, 525, 528 Roode, Thea van, xvii, 78, 519 rook, 667 rooster, 899 rooted trees, 604 rose-garden, 525, 537, 538 Rosh Hashana, 909 Rosser, Barkley, 924 Roth, T., 925 round the world, 381 roundabout, 327, 369, 380 row-rank, 710 Roy, Constant, 713 Rozenberg, G., 764 Rubik's Cube, 868–876, 917, 918 Rubik, Ernő, 868 Ruchonet, 841 Ruderman, Harry D., 769 Rugs, 478 Rule Atomic Weight, 242 Berlekamp's, 77, 78 C.A.B.S., 388-393 Deficit, 826, 831 Difference, 74, 404 Doomsday, 903-906

Downsum Absorbancy, 361 Flow, 201, 202, 204, 210 Go-dialects, 16 Inequality, 348, 349 Long Chain, 546, 549, 550 loonv addition. 399 Mex, 56, 418, 556, 558 Misère Mex, 418 Misère Nim, 418 Misère Play, 15, 279, 281, 413 misère remoteness, 279, 281 Nim-Addition, 59, 73, 390, 418, 556 Normal Play, 12, 14, 278, 279 One-upmanship, 242 remoteness, 279, 281 Restricted Translation, 261 Simplicity, 22, 24, 25, 27, 39, 45, 48, 304, 305, 307, 314, 317 Smith's, 388-389, 392-393 suspense, 287, 303 Tally, 304, 305, 315, 316, 326 Two-ahead, 199, 200, 242, 246, 249 With, 256, 257 Without, 256, 257 Wythoff's Difference, 74, 404 Rule of Three, 811, 812, 816, 834 Rule of Two, 811 Ruler Eights, 470 Fifteens, 470 Fives, 470 Fours, 470 function, 98, 436, 437, 470 Game, 469, 470, 478, 483 Sevens, 470 Sixes, 470 ruler function, 98 rules 14 Li's loopy Hackenbush, 344 Russ, I., 767 Russia, 894 Sabidussi, G., 539 Sackson, Sidney, 729, 769 Safe Dancing Haven, Fox's, 695 Sakuta, Makoto, 766 Salomaa, A., 764 saltus, 99, 114, 117, 144, 186, 752

Sarsfield, Richard, 79 Sasaki, Nobusuke, 762, 766 Saskatchewan landscape, 156 Sato's Mava Game = Welter's Game, 427Saxon Hnefatafl, 666 scale, Grundy, 87, 89, 91, 94, 96, 100, 398, 399, 406Scare'm Hare'm Position, 719-721 Schaeffer, Jonathan, 18, 757, 763, 764 Schaefer, Thomas J., 109, 118, 224, 226, 227 Scarne, John, 769 Schaer, Jonathan, 43 Scheinerman, E. R., 539 Schocken, Wolfgang Alexander, 925 Schoen, Alan, 897-898, 923 Schroeppel, Rich, 939 Schuh, Prof. Frederick, 118, 539, 640, 710, 713, 725, 729, 769, 732, 769 Schuhstrings, 523, 524 Schwenk, Allen J., 539 Scissors-Paper-Stone, 15 score function, 837, 838 Scorer, Richard S., 923 scorpion, posing as insect, 600 Scott, Elizabeth Anne, 869, 871, 875 Scott, Katherine, xvii, 578, 584 scout, 821-823, 835 Scrabble, 760, 765 scrap-heap, 391-392 scrimmage formation, 687, 697 line of, 687 position, 694, 700, 701 region, high 697, 700, 701 sequence, 687-689, 692, 693, 697, 700, 705 values, 688, 697, 700 Seal, David J., 635, 871, 875, 914 seasoned campaigner, 164 Seating Boys and Girls, 132, 133, 179, 261, 310, 366Seating Couples, 44, 45, 95, 132, 133, 261 Seating Families, 95, 261 second cousin, 101, 102, 104–106, 114 second player wins, 28–30 Secondoff Algorithm, 535, 537

 \diamond

996

secrets gory, 911 hidden, 844-846, 848, 894, 896, 897, 899 Seemann, Markus, 768 Seki, Hirohisa, 766 Select Boys and Girls, 310–311 selective compound, 299, 300, 312, 396 selective compound, shortened = severed, 312 Selfridge, John Lewis, 522, 740, 767, 768 sente, 161, 187, 188 Seo, Masahiro, 766 separation, 703 Sepoys, see Maharajah, 710 sequence backbone, 700, 701, 704 early, 688, 689 scrimmage, 687-689, 692, 693, 697, 700, 705 serpent, 899 \mathbf{set} empty, 82, 399 move, 516 subtraction, 84, 86, 442 variation, 220, 221 Seven-in-a-Row, 740 Seven-up, 218, 225 severed head, 205, 220, 221, 312 severed selective compound, 312 sexopposite 132 significance of, 528 shackles, 565 Shader, Leslie E., 925 Shakespeare, William, 905 Shaki, Ahiezer, S., 377 Shannon, Claude Elwood, 740, 763, 767, 960 shatter, 380 She-Loves-Me, She-Loves-Me-Not = $\cdot 05$, 67, 101-103, 113, 116, 529, 607 She-Loves-Me-Constantly, e.g., • 51, 101, 103, 518Sheep and Goats, 863 see Wolves-and-Sheep, 665, 666 Shephard, Geoffrey C., 509, 539 shifting by stars, 250, 258 infinitesimally, 177, 198 ships, 928, 931, 933

Shogi, 760, 761, 765, 766 short chain, 547, 554, 562 hot games, 225 loops, 554 paths, 562 positions, 694 shortened selective = severed, 312 shortlist, 303, 304, 317 short-sighted view, 159, 160 Shortt, Mr., 745, 746 SHOT(S), 738, 739 Shurman, Jerry, 640 Sibert, William L., 446, 451, 452 Sicherman, George, xvii, 640 sickle and sickle, 836 side, 336-337, 340-341, 344-348, 350, 354-356, 369sidelined king, 661, 663 sidling, 338-342, 346, 350, 365, 371-377 Sidling Theorem, 340, 371–373 Siegel's cqsuite, 687, 690, 692, 697, 702, 703, 705, 766 Siegel, Aaron, xvii, 687, 690, 702, 703, 705, 766 sign, 348-351, 373 Silber, Robert, 79 Silver Dollar Game, 491, 492, 535 Silverman, David L., 668 Simões-Pereira, J.M.S., 607 Simonim = Similar MOve NIM, 496-498Simonson, S., 539 simplest form, 19, 22, 71, 370 simplest number, 19, 21, 22, 305, 307, 314 Simplicity Rule, 22, 24, 25, 27, 39, 45, 48, 304, 305, 307, 314, 316-317 simplifying games, 60, 62, 63, 75, 77 singleton, 413 Single Circle, 691 Singmaster, David, 871, 875, 925 sinister, 913 Sipser, Michael, 224, 226 Sisyphus, 346–348 Six-in-a-Row, 740 Six Men's Morris, 737 Ski-Jumps, 7, 9–11, 14, 15, 19, 27, 40 skittles, 81 slant, right = correct, 163

 \Diamond

slash, |, 6, 163, 366 slashes, ||, 130, 163, 366 Sleator, Daniel, 602, 607 sliding block puzzles, 877-885, 919 jungles, 199, 220, 221 slipper, 9 Slither, 767 Sloane, N. J. A., xvii(3), 489, 767 slow horses, 286, 288, 289 slow join, 286-291, 300, 302 slower join, 283 small, 36, 38, 229–262, Smallest Nim, 532 Smith Theory, 333, 379-396 Smith's Rule, 388-389, 392-393 Smith, Alvy Ray, 960 Smith, Arthur, 765 Smith, Cedric Austen Bardell, 88, 96–98, 109, 117, 118, 278, 279, 298, 312, 388, 389, 391-393, 412, 417, 452, 525, 528, 533, 923 Smith, Martin C., 765 Smith, Sally, xiii snakes, 40-42, 240, 929 Snakes-and-Ladders, 14, 386–388 Snaky, 748, 768 Snatzke, Raymond George, 756, 763 Snort, 47, 145-147, 149, 151, 153, 154, 156, 157, 161, 167-169, 177, 178, 180-183, 224, 502 dictionary, 147, 180-183 lemmas, 180 Solitaire, 15, 503-840 army, 821-822, 835 board, English, 710, 730 central, 805-809, 811, 837-839 Fool's, 826, 835 -like puzzles, 863 solutions to problems, 709, 761, 762 Soma, 843-846, 897, 910-913 Somap, 897, 910–913 sophistication levels, 569 Soulé, Stephen, 765 sound bound for a hound, 723 Sowing Games, 761, 766 spaceships, 931, 933, 936 spades, xix(1), 359-363 Spain, 905

span, 710 span-length, 194, 195 spanning tree of graph, 217, 224 spar, 488 spare move, 838 Sparring, 488 sparse space, 110–113 species, 435 Spight, William, xvii, 764, 765 Spinner, The, 833 spinster, 506, 507 splitting the atom, 234–236 spoiler, 161 spokes, 549 spot, 145, 146, 391, 585 Spots and Sprouts, 95 Sprague-Grundy Theory, 56, 117, 118, 220, 333, 556-558, 569, 636 Sprague, Roland Percival, 41, 56, 79, 220, 333, 506, 539, 540 Sprouts, 598-602 Brussels, 603, 604, 607 Squandering Squares, 254 square, 731-765 half-perfect, 886, 890 magic, 731, 886-891 Nasik, 891 pandiagonal, 891 perfect, 518, 886 quarter-perfect, 886, 889, 890 Strong, 713, 717, 718 Weak, 713, 715, 717-721 square-eater, 643 Squares, MacMahon, 899–902, 921–922 Squares Off, 319 Stability Condition, 363 stable positions, 705 stack of coupons, 693 stage, 392-395 Stage, early, 689 Staircase Fives, 499 stalemate = tie, 14 stalk, 29, 30 stalk = stem, 67Stalking = $\cdot 31$, 429, 444 standard form, 101, 103-105, 107-109, 114

star, 34, 39, 47, 48, 65, 68, 69, 120, 121, 125, 155, 156, 190, 229, 261 far, 230-232, 236, 137, 239, 243-251, 258, 259lucky, 246, 249 remote, 230-232, 236, 237, 239, 240, 243-251, 258, 259 thermorgraph of, 156 Star-Incentive theorem, 259 Star-Shifting Principle, 250, 258 Stars-and-Stripes, 603, 604 starting position, 14 startling value, 39 Stead, W., 925 Steingrímsson, E., 117 Steinhaus function = remoteness, 279 Steinhaus, Hugo, 279, 298, 393 stem, 36, 47, 67, 199, 200, 210, 230, 240, 242 step, 529 Stewart, Bonnie M., 841 still life, 929 Stiller, Lewis, 763 Stockmeyer, Larry, 223, 226, 227 stolid survivor, 829 Stone, John, 112 stones black = blocking, 642, 644-646, 654-657Go, 642, 646 lifting, 346 non-static, 661 resetting, 346 rolling, 346 static, 661, 663 strategic, 643, 650, 652, 654, 657, 658, 661 tactical, 643, 652 unlimited supply, 650, 664 useful, 663 wandering = white, 642, 644-646, 648, 650,654-658, 661, 663 well-placeed, 656 stop, Left and Right, 165, 166, 697 stopper, 337, 340-341, 351-352, 354, 356-357, 362, 369-370 stopping position, 149–153, 156–166 value, 149, 153, 173 Storer, James, 224

Strategic Landing Plan, 694 Geese's, 696 strategic stones, 643, 650, 652, 654, 657, 658, 661 strategy 173, 314 Abacus, 513 copying, 642 Goller's 408 Hare's, 718 kite, 243, 244, 246 misère Lucasta, 593 stealing, 406, 556, 616, 744 survival, 371–373 Swivel Chair, 349 symmetry, 3, 194, 547, 550 Thermostatic, 158, 163–167, 183–185, 188, 225Tweedledum and Tweedledee, 2, 3, 35, 349, 350winning, 28, 31, 32, 46, 314 streak, 483 Streaking, 483 stream full, 948 glider, 946 thin, 947 thinning, 946 string, 741-750, 849-858 air on a G-, 97 q-, 524 Hackenbush, 22, 24, 77, 78, 194, 195, 329 of pearls, 522, 523 Strings-and-Coins, 550-555 Strip and Streak, 483, 487 stripping, 708, 814 strong squares, 713 structure of periods, 109 subperiods of nim-values, 109 subselective compounds, 396, 533 subtraction game, 84, 86, 87, 98, 395, 430, 442, 529 - 532subtraction set, 84, 86, 442 succour the sucker, 826, 836, 837 suicider, 312, 317, 320, 558 sums eternal, 46 galvinized, 253

of games, 31, 120,121, 161, 163, 165, 166, 220, 277 of nimbers, 42 NP-hard, 225 ordinal, 220 Sun Xin-Yu, 640 sunny positions, 397-401, 404 superdominoes, MacMahon's, 899 superheavy atoms, 333 superkos (See also loopiness), 758 superstars, 261 surprise exam, 534 Surreal Numbers, 18 survival, 348, 371-373, 413, 927 survivor, stolid, 829 suspense, 173, 277, 287, 288 numbers, 286-289, 292, 294, 315, 317 rules, 287, 303 swanpan, 513 Swedes, 666 Swedish King, 666 Sweets and Nuts, 393–394, 410 Swirling Tartans, 476, 477 switch engine, 937 switches, 121-125, 157, 706 Switching Game, Shannon, 744-746, 767, 768 Swivel Chair Strategy, 349 Sylver Coinage, 15, 539, 609–631, 635–640 Sylvester, James Joseph, 640 Sym, 473, 479 symmetric group, 867 symmetrical player, 547 symmetry, 51 symmetry rule's OK, 548 symmetry strategy, 3, 194, 547, 550 Sympathetic Nim, 494 Sympler, 473, 479 Synonim, 494-496 T-move, 165, 184 tabella, 736 Tableux de Young, 766 Tablut, 666 tackle, Phutball, 754 Tactical Amendment, 697

move, 643 stone, 643, 652 tactically worthless, 652 tactics, corner, 654 tail, 807 tails, 396, 402, 404 Tait, Hilary, 310 Tait, Peter Guthrie, 916 Take-A-Square, 430, 442 take-and-break games, 81–117 take-away games, 82, 84, 86, 87, 98, 101 Taking Squares, 429 Takizawa, Takenobu, 765 tally, 300-326 machine, 308, 316, 318 rules, 304, 305, 308, 315, 316, 325, 326 tame, 417, 422-438, 443-446 tameable, 425, 446 Tangrams, 895 Tantalizer, The Great, 892-894 Tapson, Frank, 920 tardy union, 312 Tarjan, Robert Endre, 224–226 tartan, 476 Tartan Theorem, 477 Tartans, Swirling, 476 tax exemption, 151-155 Taylor, Andrew, 875–876 Tchouka(illon), 761, 766 Tego, Theodore, 762 temperature, 124, 125, 131, 151-155, 164-173, 184, 185, 231 ε . 706 ambient, 164, 165, 188 critical, 167-171 policy, 124–125, 131–132 tendril, 564, 567, 568, 576, 577, 579 tendrilled crosses, 582 Tennis, 15 tentative tally, 306, 310 tepid component, 326, 327 tepid game, 308, 316, 327 tepid position, 306 tepid work, 316 terminal position, 2, 7, 28, 41, 416 ternary = base 3, 471, 472, 481ternary Gray code, 861

 \Diamond

tactical

diagram, 691

1000

Index

 \diamond

Ternups, 470-472 tesseract, magic, 891 tesseravore, 642 Test Remote Star, 246, 248, 249 Uppitiness, 245 Thatcher, J.W., 960 The More the Merrier, 533 Theorem At-least-one, 258 Decomposition, 565, 567 Don't-Break-It-Up, 213, 214, 216 Euler's, 571 Fundamental, of Zeroth Order Moribundity, 598Half-Tame, 435 Harmless Mutation, 562–564 Intermediate Value, 426, 438 Max-Flow, Miin-cut, 201, 205, 211 Mock Turtle, 464, 466 n-, 615 Noah's Ark, 423, 432–435, 443 non-arithmetic periodicity, 115 Number Avoidance, 147–149, 183 on simplifying games, 75-77 p-, 615 Quiet End, 618 Redwood Furniture, 212-214 Sidling, 340, 371–373 Simplest Form, 370 Star-Incentive, 259 Tartan, 477 Thirty-One, 723, 729 Twopins Decomposition, 500, 567 Uglification, 486 Zeckendorf's, 535 tetromino, 894 Theorem Beasley's Exit, 829, 830, 837, 838 Euler's 571, 584 Even Alteration, 511 Quiet End, 618, 629 Thirty-one, 723, 725, 729 theory, Green Hackenbush, 190–196, 201 Theory, Smith, 323, 395 thermal dissociation, 168

thermograph, 151-159, 161-165, 168-172, 174-176, 184, 185, 187, 358 compound, 164 extended, 161, 162 foundations, 155 four-stop, 158, 159 of **oof**, 358 thermographic thicket, 176 thermographs of star and up, 156 thermography, 151–174, 176–179, 188, 225, 262, 690 extended, 759 generalized, 759, 764 thermostat, 164, 185 THERMOSTRAT = Thermostatic Strategy, 163-167, 183-185, 187, 188, 224 thinking backwards, 384 forwards, 384 laterally, 384, 412 thinning a glider stream, 947 third cousin, 101, 103-105, 107, 109, 114, 116 Third One Lucky, 520 thirding, 620 thirteen's unlucky!, 420 Thirty-One Theorem, 723, 725, 729 Thistlethwaite, Morwen, 875, 876 Three B'ars Problems, 827 thunderbird, 937 Thorp, Edward, 765 Thompson, Ken, 632 Three B'ars Problem, 827 Three Men's Morris, 737 Three Up, 737 Three-Color Hackenbush, see also Hackenbush, Hotchpotch Three-Finger Morra, 15 three-quarters, 6, 17 timesstar, nim-product, 475 thumb-twiddling, 405, thunderbird, 937 Thursday, Maundy, 26 Tic Tac Toe = Noughts-and-Crosses, 114, 94, 731-736, 738, 740, 743, 749 tie \neq draw, 14 time bomb, 937 time, complete in exponential, 224

timer, 300,-303, 305-308, 310, 314-318, 322, 325 - 327tims, 475 tiniest, 706 tiniest value, 706 Tinsley, Marion, 763 tinted, 38, 39, 47-49, 51, 204, 206, 209, 210 tiny, 353, 357 tiny, 126, 127, 169, 170 -a-quarter, 132 -two, 126 -x, 125F&G, 706 tis, tisn, 342, 344, 374 Tit-Tat-Toe = Tic-Tac-Toe, 731toad, 930 Toads & Frogs, 12, 13, 40, 63–65, 67–70, 75, 76, 127-131, 133-137, 367, 375-377, 690, 863 toenail, 158 Toeplitz, Otto, 571, 584 toil, honest, 616 toll, 300-303, 305-306, 318, 310, 312, 315-321, 325hacking, 686 infinite, 312-320, 325-326 tombstone, 318 toolkit, 690 Top Entails, 396 top row traps, 697 Tortoises, see Hares, 862, 916 Tower of Hanoï, 861 trace, 712, 714-716, 720, 723, 728, 729 tracking, 65, 202, 205, 222 track = path, 202-209, 221, 222Trading Triangles, 254 traffic lights, 928, 929, 934, 935, 937 trailing, 402, 403 trains, 402, 937 transition phase, 168 two-ish region, 702 transitional region, 701 translation by nimbers, 259, 261 by numbers, 149, 153, 158

of four-stop games, 157-158 of switches, 123 traps, top row, 697 travestv, 286 Trawick, Charles, 936 Treblecross = $\cdot 007$, 94, 112 tree Australian, 22, 24, 214 binary, 22, 24 game, 40 green, 191-193 greenwood, 34 infinite, 332 redwood, 214-217 spanning, 216, 222, 224, 746 with extra twig, 214, 216 trev. 357 triamond, 895 triality traps, 714 triality triumphs, 526 Triamond, 896 Triangles, Trading, 254 triangular numbers, 254, 520 Tribulations, 520, 535 Tricky Six Puzzle, 867, 868, 876 trifurcated range, 705 trifurcated region, 701, 702 Trigg, Charles W., 919 Trim, 534 trimmed, 675 trimmed Delphinium, 686 Triplet Fives, 470 Triplets, 469 Triplicate Dawson's Kayles, 261 Triplicate Nim, 114, 116 tripling, 620 tromino, 895 Tromp, John, xvii, 673 truth, awful, 416 Tsai, Alice, 765 Tschantz, Steve, 52 Tsyan-Shizi = Wythoff's Game, 427tub, 929, 937 Tubergen, G. J. van, 384, 412 tuft, 502, 504, 506 tumblers, 936 Tuppins = Twopins, 500

1002

Index

Turing, Alan Mathison, 960 Turn-and-Eatcake, 236 Turning Corners, 473-475, 478 Turning Turtles, 461 Turnips, 470–472, 481–483 Tutte, William T., 538 Tweedledum and Tweedledee, 2, 3, 35 argument, 3, 35, 72 strategy, 349 twig, 192, 193, 196, 213, 216 twigs pale and pink, 344 redwood, 213, 214, 216 Twins, 469, 478 Acrostic, 473, 474, 483 Twisted Bynum, 236 Two-Ahead Rule, 199, 200, 241 two-dimensional games, 332, 333, 473-487, 632 Two-Dimensional Nim, 332–333 two-ish transition, 701, 702 two and two, 419 Twopins, 500, 501, 564, 566 vine, 546, 566-568 Two place, 713 Tyrrell, J.A., 924 U-turns, 754, 838 Uehara, R., 769 uggles, 483 uglification, 483-487 table, 485 Theorem, 486 ugly product, 483 Uiterwijk, Jos W. H. M., 143, 762, 765, 766 Ulehla, J., 607 Ullman, Daniel, 539 ultimate periodicity, 99, 101, 112, 289, 522, 529, 627, 752 unboundedly unbounded, 610 uncertainty, 230 under, 341, 353 underlying economy, 151 union, 299-326 misère partizan, 312 of variation sets, 605 tardy, 312 urgent, 312, 316 United Kingdom, 894

units, 824 fickle and firm, 424 universal machine, 957 unparted jungles, 210 unpredictability, 934 unproductivity, 825 unrestricted tallies, 314-318 unruly, 116 unsnappable vine, 567 unstable positions, 705 untinted nodes, 50, 204, 206, 210 **up**, \uparrow , 63–72, 150, 155, 229, 232, 234–236, 242, 247, 249-252, 255, 258, 261, 262 up-second, 235, 341 up-onth, 375 up, F&G, 706 upon, 341, 353, 708 upon* = delphinium, 341, 345, 353, 355, 708uppitiness = atomic weight, 200-211, 221, 222, 230-232, 234-242, 245-254, 258-262 uppity, equally, 242, 245 upset board, 317 upstart equality, 71 upsum, 336, 355, 357, 358, 367 urgent unions, 312, 316 Ussher, Archbishsop James, 905, 925 value 4-13, 17, 19-28, 33, 36-39, 41-45, 47-52,

57, 62-69, 71, 72, 75-78, 125, 690-710 complicated, 707, 708 early, 700 infinitesimal, 36, 169-171, 173, 229, 539, 690 initial, 698, 699 of Fox-and-Geese, 673 of the Initial F&G Positions, 688 mast, 152-154, 165, 166 mean, 149, 151-154, 165, 166, 188 startling, 38 stopping, 149, 150, 173 tiniest, 706 value versus average, 10 values Childish Hackenbush, 43, 52 Col, 37-39, 47-51 common, 110-112 Cram, 505, 506

 \diamond

Cutcake, 25 Domineering, 120-122, 138-142, 177 entailed, 397-404 exceptional, 90-92, 101, 108 excluded, 111 F&G, circled FOXTAC, 691 F&G, early, 689 F&G, initial, 698-701 Hackenbush, 4–6, 17, 19, 20, 23, 28 irregular, 90-92, 101, 108 loony, 558 loopy, 387, 389-391 Maundy Cake, 26, 220 nim-, 82-94, 96-102, 104-117, 191-196, 230. 387-408, 410, 411, 425, 426, 430, 434, 436, 442, 444, 446-450, 462-466, 469-488, 498, 506, 507, 514, 521, 529-533, 537, 557-560, 567-569, 579-583, 586-590, 597, 604-606, 636, 640, 751, 752 Nimstring, 556–560, 564–569, 581–583 non-loopy, 389 putative, 471, 472 rare, 110-112 redwood bed, 216 regular, 90, 91 Seating Boys-and-Girls, 132, 179 Seating Couples, 45 Ski-Jumps, 9-11, 19 small, 36, 229-262, Snort, 147, 161, 167-169, 177, 180-183 Streaking, 484, 487 Stripping, 484, 487 switch, 121-125 Toads-and-Frogs, 13, 127-131, 133-137 van den Herik, H. Jaap, 740, 765-767 van der Meulen, Maarten, 766 van Roode, Thea, xiii(3), 78, 519 Vandeventer, Joan, 925 Varga, Tamas, 875 vanishing reactions, 941–943, 947–948, 951 variations, 94, 310 variation set, 220, 605 item varieties, 362 victory, 163 vines, 564-570, 575-580, 582 virus, 936 voracity, 945 Vout, Colin, 38, 749

Wainwright, Robert T., 930, 960 Walden, W. E., 765 Walker, R.J., 924 Wallis John, 925 Walrus, 461-462 Walsh, J. L., 52 waltz .6, 96 wandering stones, 642-664 war, cold, 299, 300-302, 304, 306, 307 Ward, Steve, 939 warfare, jungle, 210 Wari, 761 warming, 187 Waterman, Lewis, 760 Watkins, Harold, 925 Weak squares, 713, 715, 717–721 weight, see atomic weight Weigi (See Go) welt, 506 Welter function, 506–508, 510–514 Welter's Game, 427, 506, 507, 510, 514, 515 Welter, C. P., 506, 540 Welton's Delphinium, 675, 693-695 Welton's region, 688, 693, 695, 697, 699–702, 704, 705 Welton, Jonathan, xiii, 673, 710 West, Julian, xiii, 144, 584 Whim, 534 Whinihan, Michael J., 540 Whist, 275 white heap, 532 White Knight, 56-59, 278 White, Farmer, 145-147 white stones = wandering stones, 642-664Whitgift, Archbishop, 905 whole numbers, 19 wholeness of Hackenbush Hotchpotch, 251 width, 164, 184 wiggly line, 560, 581 wild animals, 430-431 wild games, 430-431, 434-435, 437-438 Wilder, Thornton, 414 Willmott, S., 765 Wilson, David, xiii Wilson, Neil Y., 118 Wilson, Richard M., 867, 876, 925 win quickly!, 278

Windows, 478, 479, 481 winners and losers, 748 winning post, 395, 411 wire and string puzzles, 849-861 Wolfe, David, xvii(3), 18, 141, 143, 188, 224, 225, 227, 690, 710, 763-765 Wolves-and-Sheep, 665, 666 women beautiful and intriguing, 524 other, 44 wonders, numberless, 119 working out a horse, 28 world lost, 413-414 small, 229-262 World Games Review, 756 worthwhile move option, 213-216 worthy prelate, 905 wrestling match, Paterson's, 902 Wright House, 372 Wyt Queens, 59, 60, 74, 402, 427-428 Wythoff's Difference Rule, 74, 404 Wythoff's Game, 15, 60, 74, 427 Wythoff, W. A., 74, 79

Yamasaki, Yohei, 427, 430, 454, 540 Yanai, K., 766 Yedwab, Laura J., 188, 224, 227 Yeong-Nan, Yeh, 767 Yes!, 53, 112, 348 Yesha, Y. 224, 226 Yoshimura, J., 762 You-nit, 609 Zamkauskas, Walter, 756 Zeckendorf algorithm, 535, 537 Zeckendorf Theorem, 535 Zeilberger, Doron, 640 zero, 2, 7, 19, 28-35, 41 deriders of, 483 game, 7 place, 713 position, 2, 3 Zetters, T.G.L., 740, 768 Zieve, Michael, 118, 540 zigzag, 177 Zig-Zag, 632 zoo, Good Child's, 427 Zuzarte, Maria S.N., 607

1004